

CS60025 Algorithmic Game Theory

Autumn 2026

Quiz 2

17 September, 3:00 PM – 3:20 PM

Marks :10 × 1.5=15

Name: _____

Roll No.: _____

Please write your answers in the space provided. Do not mark on the questions.

Abbreviations: Dominant Strategy Equilibrium (DSE), Pure Strategy Nash Equilibrium (PSNE),
Mixed Strategy Nash Equilibrium (MSNE), Correlated Equilibrium (CE), Coarse Correlated
Equilibrium (CCE)

1. A finite game is a potential game if and only if for every two strategy profiles $\mathbf{s}^1, \mathbf{s}^2$ that differ in two players' choices (say, i and j), the following holds.
 - A. $\pi_j(\mathbf{s}^2) - \pi_j(s_i^2, \mathbf{s}_{-i}^1) = \pi_i(\mathbf{s}^2) - \pi_i(s_j^2, \mathbf{s}_{-j}^1)$
 - B. $\pi_i(s_i^2, \mathbf{s}_{-i}^1) - \pi_i(\mathbf{s}^1) = \pi_i(\mathbf{s}^2) - \pi_i(s_j^2, \mathbf{s}_{-j}^1)$
 - C. $\pi_i(s_i^2, \mathbf{s}_{-i}^1) - \pi_i(\mathbf{s}^1) + \pi_j(\mathbf{s}^2) - \pi_j(s_i^2, \mathbf{s}_{-i}^1) = \pi_j(s_j^2, \mathbf{s}_{-j}^1) - \pi_j(\mathbf{s}^1) + \pi_i(\mathbf{s}^2) - \pi_i(s_j^2, \mathbf{s}_{-j}^1)$
 - D. $\pi_i(s_i^2, \mathbf{s}_{-i}^1) - \pi_i(\mathbf{s}^1) + \pi_i(\mathbf{s}^2) - \pi_i(s_i^2, \mathbf{s}_{-i}^1) = \pi_j(s_j^2, \mathbf{s}_{-j}^1) - \pi_j(\mathbf{s}^1) + \pi_j(\mathbf{s}^2) - \pi_j(s_j^2, \mathbf{s}_{-j}^1)$
2. Which of the following is true?
 - A. Every CE is an PSNE.
 - B. Every MSNE is a CCE.**
 - C. Every CE is a DSE
 - D. Every CCE is an MSNE.
3. Pick the incorrect option. Computing PSNE for symmetric congestion games is
 - A. in NP
 - B. PLS-Complete
 - C. in TFNP
 - D. in P**
4. Which of the following is not true for the problem Π of computing MSNE in a finite bimatrix game (A, B) , where players 1, 2 have n, m choices respectively ($A \in \mathbb{R}^{n \times m}, B \in \mathbb{R}^{m \times n}$)?
 - A. $\Pi \in \text{PPAD-Complete}$
 - B. $\Pi \in \text{FNP-Complete}$**
 - C. $\Pi \in \text{P}$ when $B = -A^T$
 - D. $\Pi \in \text{FNP}$
5. Best response dynamics on a finite potential game
 - A. always takes exponential time to converge to a PSNE.
 - B. always converges to a PSNE.**

- C. may not reach a PSNE in finite time.
D. reaches a PSNE in time polynomial in the game parameters
6. When every player runs no external regret dynamics, the time averaged product distribution of the distributions output by the players is a
- A. CE
B. Approximate CCE
C. Approximate CE
D. CCE
7. Consider a network congestion game on two parallel edges e_1, e_2 shared by 3 players. The edge costs are $c_1(k) = 3k$ and $c_2(k) = k^2 + 1$, where k is the number of players using edge e . If 2 players choose e_1 and 1 player chooses e_2 , what is potential function value $\Phi(s) = \sum_e \sum_{i=1}^{n_e(s)} c_e(i)$?
- A. 15
B. 11
C. 13
D. 9
8. Consider the following 2×2 game:

	C	D
C	(5, 5)	(1, 6)
D	(6, 1)	(0, 0)

A mediator samples an outcome from the joint distribution $\sigma(C, C) = \sigma(C, D) = \sigma(D, C) = \frac{1}{3}$ and $\sigma(D, D) = 0$. Which of the following statements regarding their conditional expected payoffs is correct?

- A. $\mathbb{E}_{s \sim \sigma}[u_1(s) \mid s_1 = C] = 3$ **and** $\mathbb{E}_{s \sim \sigma}[u_1(D, s_2) \mid s_1 = C] = 3$
B. $\mathbb{E}_{s \sim \sigma}[u_1(s) \mid s_1 = C] = 3$ and $\mathbb{E}_{s \sim \sigma}[u_1(D, s_2) \mid s_1 = C] = 3.5$
C. $\mathbb{E}_{s \sim \sigma}[u_1(s) \mid s_1 = C] = 2.5$ and $\mathbb{E}_{s \sim \sigma}[u_1(D, s_2) \mid s_1 = C] = 1.5$
D. $\mathbb{E}_{s \sim \sigma}[u_1(s) \mid s_1 = C] = 5$ and $\mathbb{E}_{s \sim \sigma}[u_1(D, s_2) \mid s_1 = C] = 6$
9. Consider the following two-player coordination game with strategy sets $S_1 = S_2 = \{B, S\}$:

	B	S
B	(2, 1)	(0, 0)
S	(0, 0)	(1, 2)

Which of the following distributions, over profiles $(B, B), (B, S), (S, B), (S, S)$, is **NOT** a CE?

- A. $\left(\frac{1}{6}, \frac{1}{3}, \frac{1}{3}, \frac{1}{6}\right)$
B. $\left(\frac{1}{2}, 0, 0, \frac{1}{2}\right)$
C. $\left(\frac{2}{9}, \frac{4}{9}, \frac{1}{9}, \frac{2}{9}\right)$

D. $\left(\frac{1}{3}, \frac{1}{6}, \frac{1}{6}, \frac{1}{3}\right)$

10. Consider an n -player cost minimisation game where $C_i(\mathbf{s}) \geq 0$ denotes the cost function for player $i \in [n]$ under strategy profile $\mathbf{s}$. At each iteration $t \in [T]$, each player i chooses a mixed strategy $p_i^t \in \Delta(S_i)$, forming the product distribution $\sigma_t = \prod_{j=1}^n p_j^t$. How is player i 's time-averaged external regret $R_i(T)$ defined, and what does no-external-regret dynamics guarantee regarding the empirical time-averaged distribution $\sigma = \frac{1}{T} \sum_{t=1}^T \sigma_t$?

- A. $R_i(T) = \frac{1}{T} \left(\max_{s_i^* \in S_i} \sum_{t=1}^T \mathbb{E}_{\mathbf{s}^t \sim \sigma_t} [C_i(\mathbf{s}^t) - C_i(s_i^*, \mathbf{s}_{-i}^t)] \right)$, and σ converges to a PSNE.
- B. $R_i(T) = \frac{1}{T} \left(\min_{s_i^* \in S_i} \sum_{t=1}^T \mathbb{E}_{\mathbf{s}_{-i}^t \sim \sigma_{-i}^t} [C_i(s_i^*, \mathbf{s}_{-i}^t)] - \sum_{t=1}^T \mathbb{E}_{\mathbf{s}^t \sim \sigma_t} [C_i(\mathbf{s}^t)] \right)$, and σ converges to an approximate CCE.
- C. $R_i(T) = \frac{1}{T} \left(\sum_{t=1}^T \mathbb{E}_{\mathbf{s}^t \sim \sigma_t} [C_i(\mathbf{s}^t)] - \min_{s_i^* \in S_i} \sum_{t=1}^T \mathbb{E}_{\mathbf{s}_{-i}^t \sim \sigma_{-i}^t} [C_i(s_i^*, \mathbf{s}_{-i}^t)] \right)$, and σ converges to an exact CE.
- D. $R_i(T) = \frac{1}{T} \left(\sum_{t=1}^T \mathbb{E}_{\mathbf{s}^t \sim \sigma_t} [C_i(\mathbf{s}^t)] - \min_{s_i^* \in S_i} \sum_{t=1}^T \mathbb{E}_{\mathbf{s}_{-i}^t \sim \sigma_{-i}^t} [C_i(s_i^*, \mathbf{s}_{-i}^t)] \right)$, and σ converges to an approximate CCE.