

Searching Elements in an Array: Linear and Binary Search

Searching

- Check if a given element (called **key**) occurs in the array.
 - Example: array of student records; rollno can be the key.
- Two methods to be discussed:
 - If the array elements are unsorted.
 - Linear search
 - If the array elements are sorted.
 - Binary search

Linear Search

Basic Concept

- **Basic idea:**
 - Start at the beginning of the array.
 - Inspect elements one by one to see if it matches the key.
- **Time complexity:**
 - A measure of how long an algorithm takes to run.
 - If there are n elements in the array:
 - **Best case:**
match found in first element (1 search operation)
 - **Worst case:**
no match found, or match found in the last element
(n search operations)
 - **Average case:**
 $(n + 1) / 2$ search operations

Contd.

```
/* The function returns the array index where the
   match is found. It returns -1 if there is no
   match.      */

int  linear_search (int a[], int size, int key)
{
    int pos = 0;
    while ((pos < size) && (a[pos] != key))
        pos++;
    if (pos < size)
        return pos;    /* Return the position of match */
    return -1;        /* No match found */
}
```

Contd.

```
int x[] = {12, -3, 78, 67, 6, 50, 19, 10};
```

- Trace the following calls :

`search (x, 8, 6) ;`

Returns 4



`search (x, 8, 5) ;`

Returns -1



Binary Search

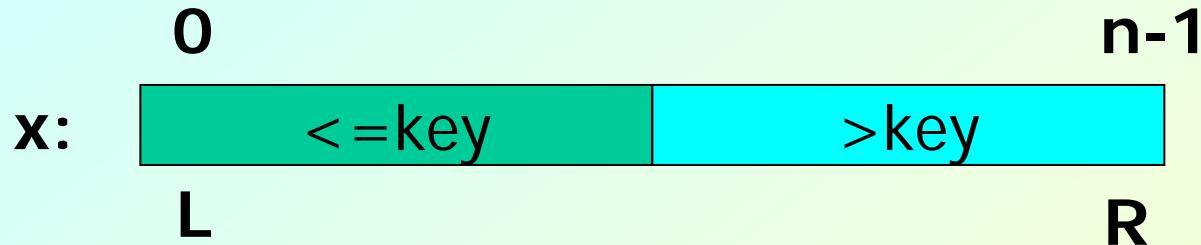
Basic Concept

- **Binary search works if the array is *sorted*.**
 - Look for the target in the middle.
 - If you don't find it, you can ignore half of the array, and repeat the process with the other half.
- **In every step, we reduce the number of elements to search in by half.**

The Basic Strategy

- What we want?

- Find split between values larger and smaller than **key**:



- Situation while searching:

- Initially L and R contains the indices of first and last elements.

- Look at the element at index $[(L+R)/2]$.

- Move L or R to the middle depending on the outcome of test.

Contd.

```
/* If key appears in x[0..size-1], return its location, pos  
such that x[pos]==key. If not found, return -1 */
```

```
int bin_search (int x[], int size, int key)  
{  
    int L, R, mid;  
    _____i  
    while ( _____ )  
    {  
        _____i  
    }  
    _____ i  
}
```

The basic search iteration

```
int bin_search (int x[], int size, int key)
{
    int L, R, mid;
    _____;
    while ( _____ )
    {
        mid = (L + R) / 2;
        if (x[mid] <= key)
            L = mid;
        else R = mid;
    }
    _____ ;
}
```

Loop termination

```
int bin_search (int x[], int size, int key)
{
    int L, R, mid;
    _____;
    while ( L+1 != R )
    {
        mid = (L + R) / 2;
        if (x[mid] <= key)
            L = mid;
        else R = mid;
    }
    _____ ;
}
```

Return result

```
int bin_search (int x[], int size, int key)
{
    int L, R, mid;
    _____;
    while ( L+1 != R )
    {
        mid = (L + R) / 2;
        if (x[mid] <= key)
            L = mid;
        else R = mid;
    }
    if (L >= 0 && x[L] == key) return L;
    else return -1;
}
```

Initialization

```
int bin_search (int x[], int size, int key)
{
    int L, R, mid;
    L = -1;    R = size;
    while ( L+1 != R )
    {
        mid = (L + R) / 2;
        if (x[mid] <= key)
            L = mid;
        else R = mid;
    }
    if (L >= 0 && x[L] == key) return L;
    else return -1;
}
```

Binary Search Examples

Sorted array

-17 -5 3 6 12 21 45 63 50

Trace :

	L= -1; R= 9;	x[4]=12;
	L= -1; R=4;	x[1]= -5;
binsearch (x, 9, 3);	L= 1; R=4;	x[2]=3;
	L=2; R=4;	x[3]=6;
binsearch (x, 9, 145);	L=2; R=3;	return L;
binsearch (x, 9, 45);		

We may modify the algorithm by checking equality with x[mid].

Is it worth the trouble ?

- Suppose that the array **x** has 1000 elements.
- Ordinary search
 - If **key** is a member of **x**, it would require 500 comparisons on the average.
- Binary search
 - after 1st compare, left with 500 elements.
 - after 2nd compare, left with 250 elements.
 - After at most 10 steps, you are done.

Time Complexity

- If there are n elements in the array.
 - Number of searches required:

$$\log_2 n$$

- For $n = 64$ (say).

- Initially, list size = 64.
- After first compare, list size = 32.
- After second compare, list size = 16.
- After third compare, list size = 8.
-
- After sixth compare, list size = 1.

$2^k = n$, where k is the number of steps.

$$\log_2 64 = 6$$

$$\log_2 1024 = 10$$

Sorting

The Basic Problem

- Given an array

$x[0], x[1], \dots, x[\text{size}-1]$

reorder entries so that

$x[0] \leq x[1] \leq \dots \leq x[\text{size}-1]$

- List is in non-decreasing order.
- We can also sort a list of elements in non-increasing order.

Example

- **Original list:**

10, 30, 20, 80, 70, 10, 60, 40, 70

- **Sorted in non-decreasing order:**

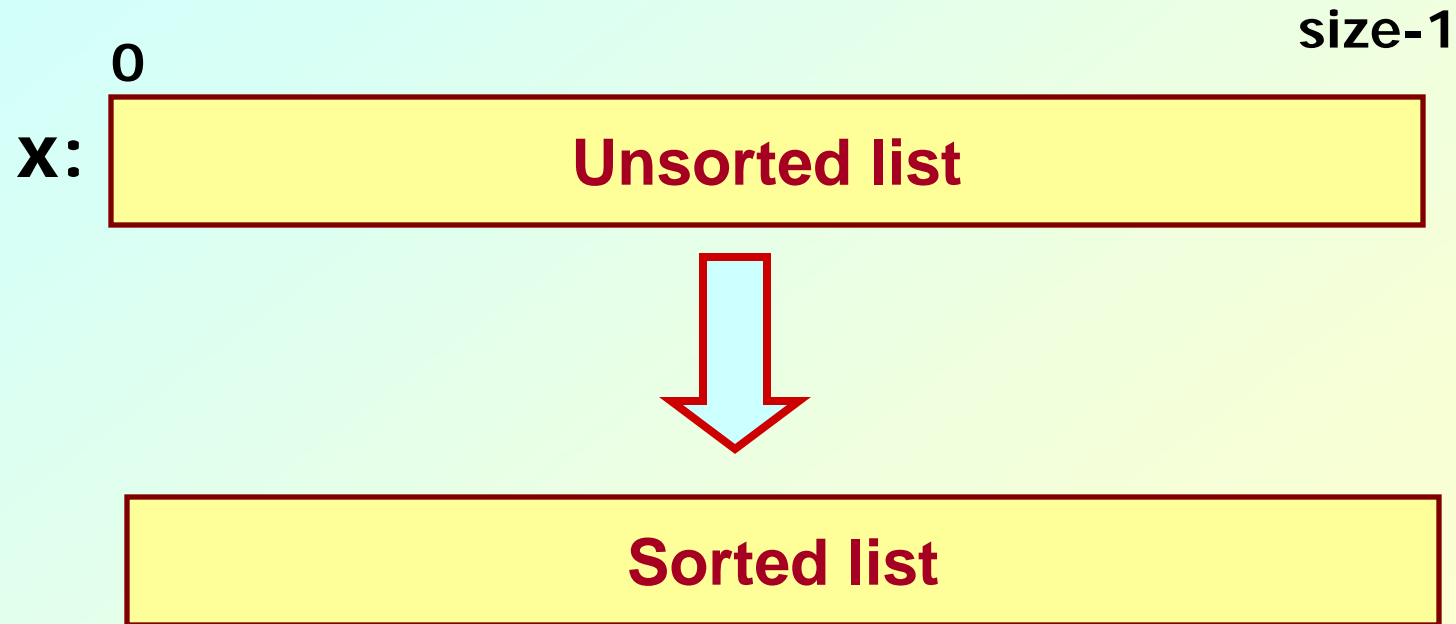
10, 10, 20, 30, 40, 60, 70, 70, 80

- **Sorted in non-increasing order:**

80, 70, 70, 60, 40, 30, 20, 10, 10

Sorting Problem

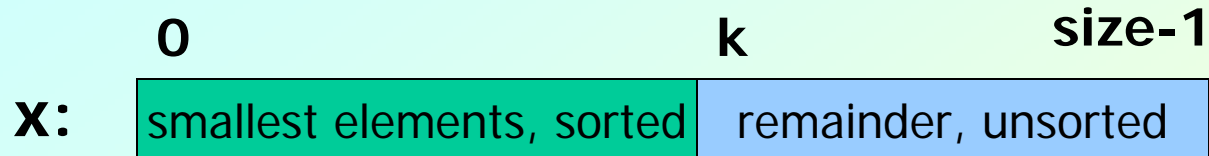
- What we want ?
 - Data sorted in order



Selection Sort

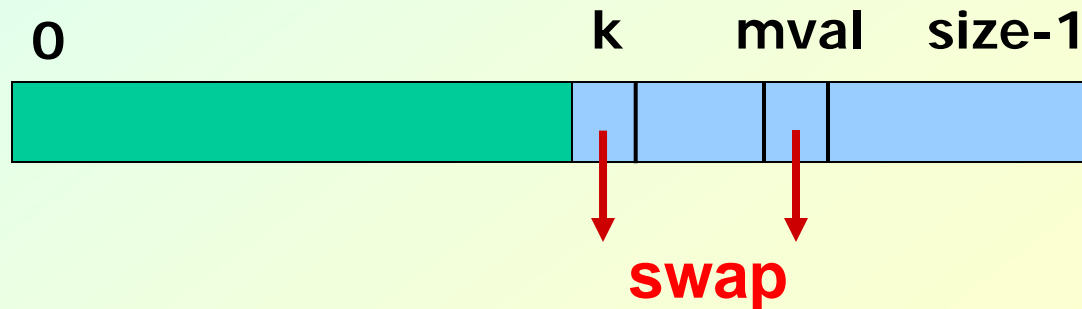
How it works?

- **General situation :**



- **Step :**

- Find smallest element, **mval**, in $x[k..size-1]$
- Swap smallest element with $x[k]$, then increase k .



Subproblem

```
/* Yield index of smallest element in x[k..size-1];*/  
  
int min_loc (int x[], int k, int size)  
{  
    int j, pos;  
  
    pos = k;  
    for (j=k+1; j<size; j++)  
        if (x[j] < x[pos])  
            pos = j;  
    return pos;  
}
```


The main sorting function

```
/* Sort x[0..size-1] in non-decreasing order */  
  
int sel_sort (int x[], int size)  
{ int k, m;  
  
  for (k=0; k<size-1; k++)  
  {  
    m = min_loc (x, k, size);  
    temp = a[k];  
    a[k] = a[m];  
    a[m] = temp; } }  
}
```

SWAP

```

int main()
{
    int x[ ]={-45,89,-65,87,0,3,-23,19,56,21,76,-50};
    int i;
    for(i=0;i<12;i++)
        printf("%d ",x[i]);
    printf("\n");
    sel_sort(x,12);
    for(i=0;i<12;i++)
        printf("%d ",x[i]);
    printf("\n");
}

```

```

-45 89 -65 87 0 3 -23 19 56 21 76 -50
-65 -50 -45 -23 0 3 19 21 56 76 87 89

```

Example

x: 3 12 -5 6 142 21 -17 45

x: -17 12 -5 6 142 21 3 45

x: -17 -5 12 6 142 21 3 45

x: -17 -5 3 6 142 21 12 45

x: -17 -5 3 6 142 21 12 45

x: -17 -5 3 6 12 21 142 45

x: -17 -5 3 6 12 21 142 45

x: -17 -5 3 6 12 21 45 142

Analysis

- How many steps are needed to sort n items ?
 - Total number of steps *proportional* to n^2 .
 - No. of comparisons?

$$(n-1)+(n-2)+\dots+1 = n(n-1)/2$$



Of the order of n^2

- Worst Case? Best Case? Average Case?

Insertion Sort

Insertion Sort

```
void insert_sort (int list[], int size)
{
    int i,j,item;

    for (i=1; i<size; i++)
        {
            item = list[i] ;
            for (j=i-1; (j>=0)&& (list[j] > i); j--)
                list[j+1] = list[j];
            list[j+1] = item ;
        }
}
```

```

int main()
{
    int x[ ]={-45,89,-65,87,0,3,-23,19,56,21,76,-50};
    int i;
    for(i=0;i<12;i++)
        printf("%d ",x[i]);
    printf("\n");
    insert_sort(x,12);
    for(i=0;i<12;i++)
        printf("%d ",x[i]);
    printf("\n");
}

```

```

-45 89 -65 87 0 3 -23 19 56 21 76 -50
-65 -50 -45 -23 0 3 19 21 56 76 87 89

```


Time Complexity

- **Number of comparisons and shifting:**

- **Worst case?**

$$1 + 2 + 3 + \dots + (n-1) = n(n-1)/2$$

- **Best case?**

$$1 + 1 + \dots \text{ to } (n-1) \text{ terms} = (n-1)$$

Bubble Sort

How it works?

- The sorting process proceeds in several passes.
 - In every pass we go on comparing neighboring pairs, and swap them if out of order.
 - In every pass, the largest of the elements under considering will *bubble* to the top (i.e., the right).

10	5	17	11	-3	12
5	10	17	11	-3	12
5	10	17	11	-3	12
5	10	11	17	-3	12
5	10	11	-3	17	12
5	10	11	-3	12	17

Largest



- **How the passes proceed?**
 - In pass 1, we consider index 0 to $n-1$.
 - In pass 2, we consider index 0 to $n-2$.
 - In pass 3, we consider index 0 to $n-3$.
 -
 -
 - In pass $n-1$, we consider index 0 to 1.

Bubble Sort

```
void swap(int *x, int *y)
{
    int tmp = *x;
    *x = *y;
    *y = tmp;
}
```

```
void bubble_sort
    (int x[], int n)
{
    int i,j;

    for (i=n-1; i>0; i--)
        for (j=0; j<i; j++)
            if (x[j] > x[j+1])
                swap(&x[j],&x[j+1]);
}
```

```

int main()
{
    int x[ ]={-45,89,-65,87,0,3,-23,19,56,21,76,-50};
    int i;
    for(i=0;i<12;i++)
        printf("%d ",x[i]);
    printf("\n");
    bubble_sort(x,12);
    for(i=0;i<12;i++)
        printf("%d ",x[i]);
    printf("\n");
}

```

```

-45 89 -65 87 0 3 -23 19 56 21 76 -50
-65 -50 -45 -23 0 3 19 21 56 76 87 89

```

Time Complexity

- **Number of comparisons :**

- **Worst case?**

$$1 + 2 + 3 + \dots + (n-1) = n(n-1)/2$$

- **Best case?**

Same

- **How do you make best case with $(n-1)$ comparisons only?**
 - **By maintaining a variable `flag`, to check if there has been any swaps in a given pass.**
 - **If not, the array is already sorted.**


```
void bubble_sort
    (int x[], int n)
{
    int i,j;
    int flag = 0;

    for (i=n-1; i>0; i--)
    {
        for (j=0; j<i; j++)
            if (x[j] > x[j+1])
            {
                swap(&x[j],&x[j+1]);
                flag = 1;
            }
        if (flag == 0) return;
    }
}
```

Some Efficient Sorting Algorithms

- Two of the most popular sorting algorithms are based on **divide-and-conquer** approach.
 - Quick sort
 - Merge sort
- Basic concept (divide-and-conquer method):

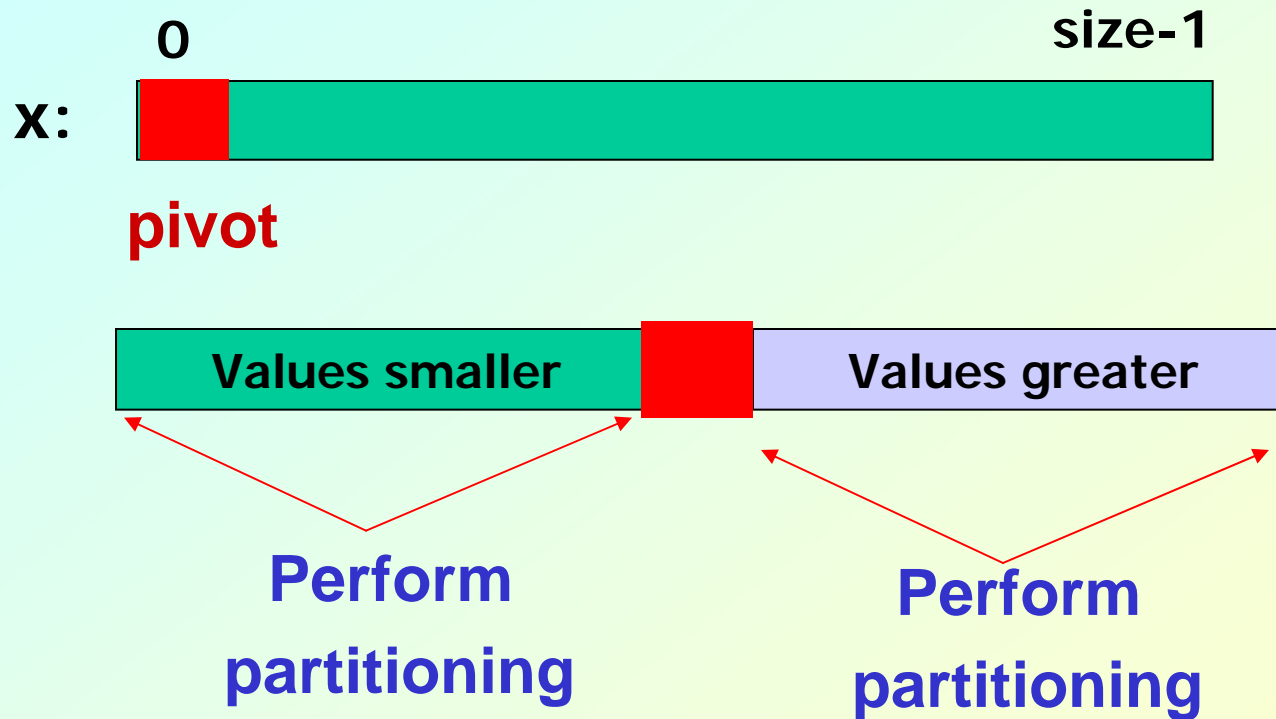
```
sort (list)
{
    if the list has length greater than 1
    {
        Partition the list into lowlist and highlist;
        sort (lowlist);
        sort (highlist);
        combine (lowlist, highlist);
    }
}
```

Quick Sort

How it works?

- At every step, we select a *pivot element* in the list (usually the first element).
 - We put the pivot element in the *final position* of the sorted list.
 - All the elements less than or equal to the pivot element are to the left.
 - All the elements greater than the pivot element are to the right.

Partitioning



Example

26	<u>33</u>	35	29	19	12	<u>22</u>
	22	<u>35</u>	29	19	<u>12</u>	33
	22	12	<u>29</u>	<u>19</u>	35	33
	22	12	<u>19</u>	<u>29</u>	35	33
<u>19</u>	<u>22</u>	<u>12</u>	26	<u>29</u>	<u>35</u>	<u>33</u>

The partitioning process

↙ ↘

Recursively carry out the partitioning

```
void print (int x[], int low, int high)
{
    int i;

    for(i=low; i<=high; i++)
        printf(" %d", x[i]);
    printf("\n");
}
```

```
void swap (int *a, int *b)
{
    int tmp=*a;
    *a=*b;
    *b=tmp;
}
```



```
void partition (int x[], int low, int high)
{
    int i = low+1, j = high;
    int pivot = x[low];
    if (low >= high) return;
    while (i<j) {
        while ((x[i]<pivot) && (i<high)) i++;
        while ((x[j]>=pivot) && (j>low)) j--;
        if (i<j) swap (&x[i], &x[j]);
    }
    if (j==high) {
        swap (&x[j], &x[low]);
        partition (x, low, high-1);
    }
    else
        if (i==low+1)
            partition (x, low+1, high);
        else {
            swap (&x[j], &x[low]);
            partition (x, low, j-1);
            partition (x, j+1, high);
        }
}
```

```
int main (int argc, char *argv[])
{
    int *x;
    int i=0;
    int num;

    num = argc - 1;
    x = (int *) malloc(num * sizeof(int));

    for (i=0; i<num; i++)
        x[i] = atoi(argv[i+1]);

    partition(x,0,num-1);

    printf("Sorted list: ");
    print (x,0,num-1);
}
```

Trace of Partitioning

```
./a.out 45 -56 78 90 -3 -6 123 0 -3 45 69 68
```

```
45 -56 78 90 -3 -6 123 0 -3 45 69 68
```

```
-6 -56 -3 0 -3 45 123 90 78 45 69 68
```

```
-56 -6 -3 0 -3 68 90 78 45 69 123
```

```
-3 0 -3 45 68 78 90 69
```

```
-3 0 69 78 90
```

```
Sorted list: -56 -6 -3 -3 0 45 45 68 69 78 90 123
```

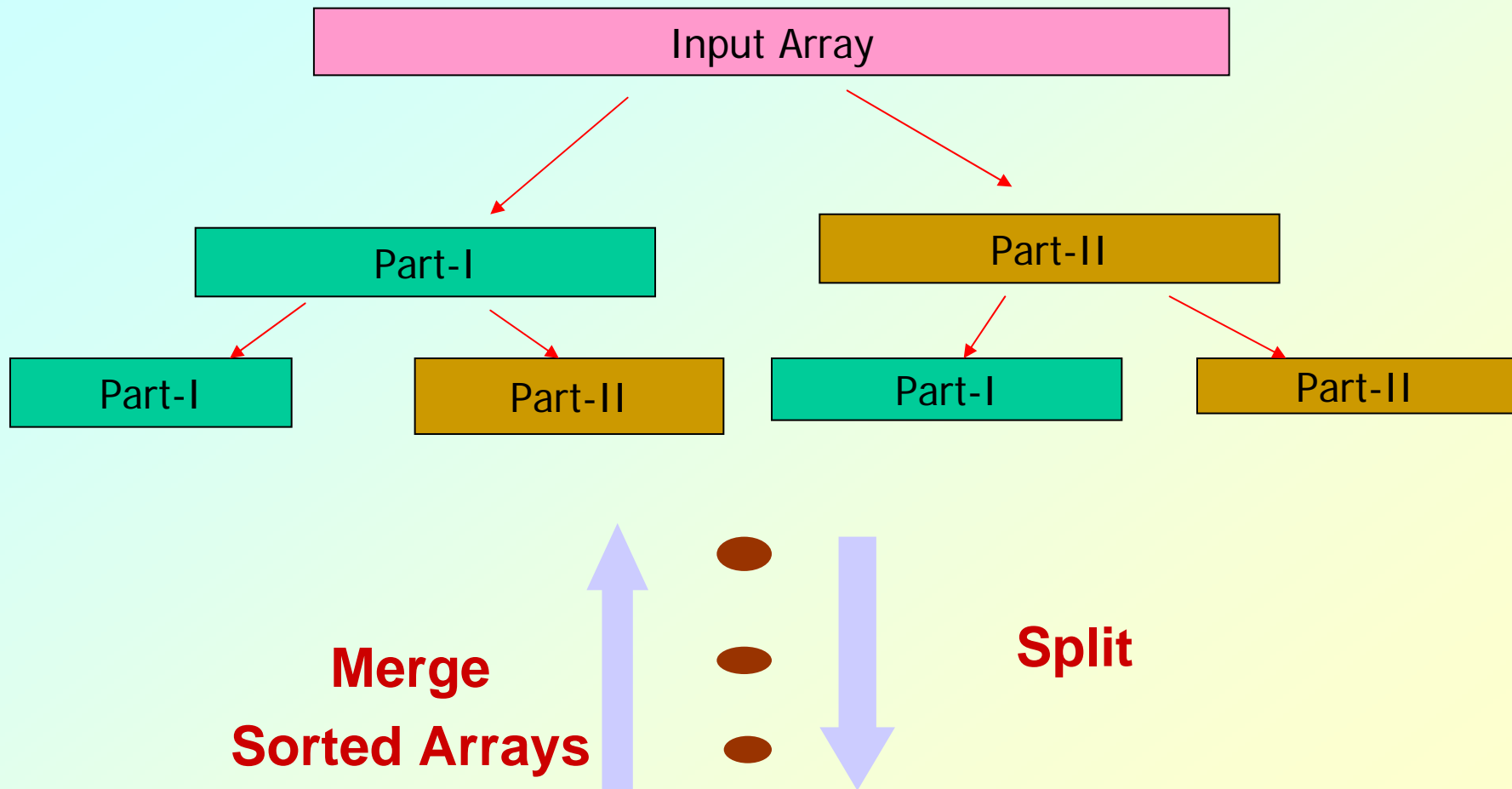
Time Complexity

- **Worst case:**
 n^2 ==> list is already sorted
- **Average case:**
 $n \log_2 n$
- **Statistically, quick sort has been found to be one of the fastest algorithms.**

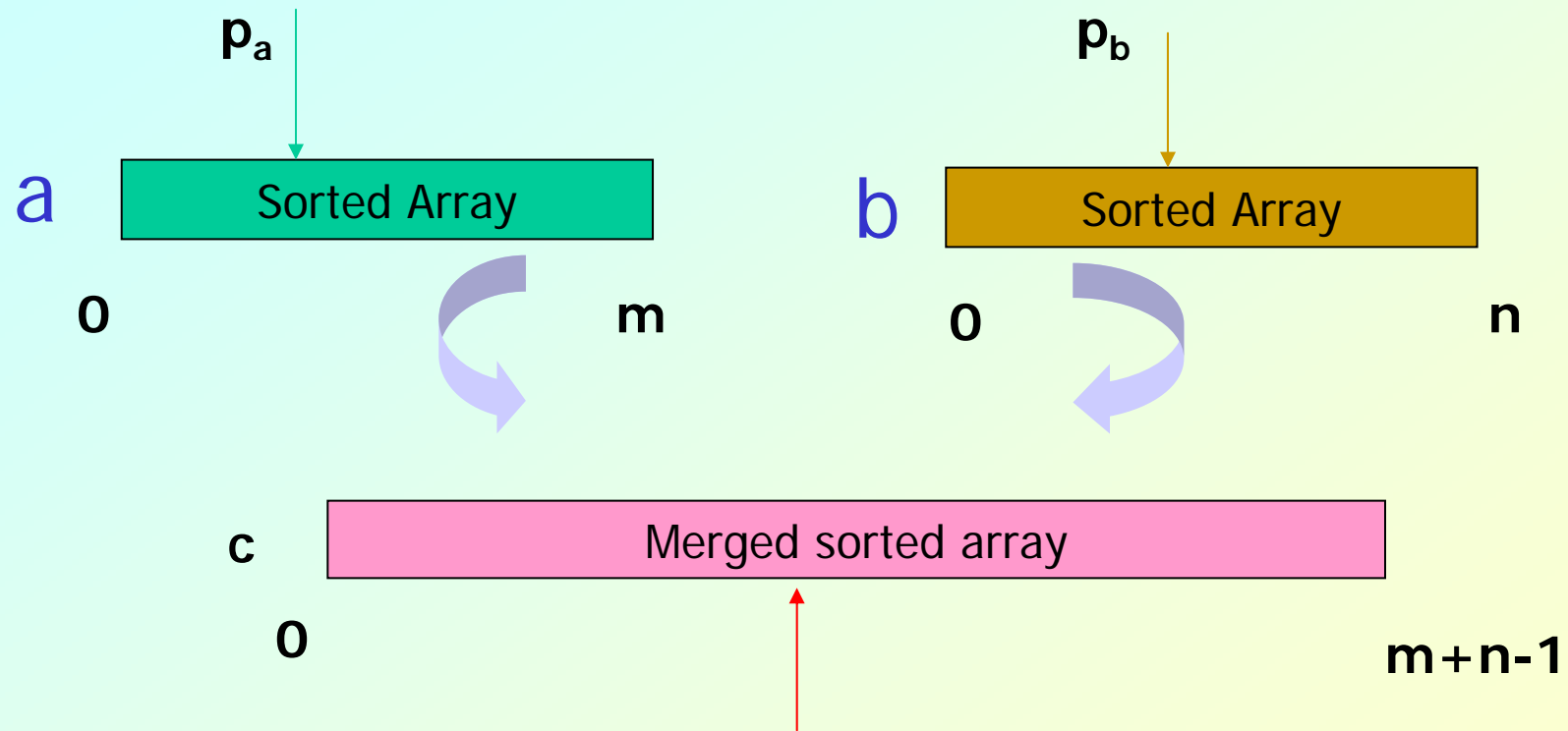
- **Corollary of quick sort:**
 - **Given a list of numbers stored in an array, determine how many numbers are smaller than a given number p ?**
 - **Given a list of integers (negative and non-negative), reorder the list so that all negative numbers precede the non-negative ones.**

Merge Sort

Merge Sort



Merging two sorted arrays



Move and copy elements pointed by p_a if its value is smaller than the element pointed by p_b in $(m+n-1)$ operations; otherwise, copy elements pointed by p_b .

Example

- **Initial array A contains 14 elements:**
 - 66, 33, 40, 22, 55, 88, 60, 11, 80, 20, 50, 44, 77, 30
- **Pass 1 :: Merge each pair of elements**
 - (33, 66) (22, 40) (55, 88) (11, 60) (20, 80) (44, 50) (30, 70)
- **Pass 2 :: Merge each pair of pairs**
 - (22, 33, 40, 66) (11, 55, 60, 88) (20, 44, 50, 80) (30, 77)
- **Pass 3 :: Merge each pair of sorted quadruplets**
 - (11, 22, 33, 40, 55, 60, 66, 88) (20, 30, 44, 50, 77, 80)
- **Pass 4 :: Merge the two sorted subarrays to get the final list**
 - (11, 20, 22, 30, 33, 40, 44, 50, 55, 60, 66, 77, 80, 88)

```
void merge_sort (int *a, int n)
{
    int i, j, k, m;
    int *b, *c;

    if (n>1) {
        k = n/2;    m = n-k;
        b = (int *) calloc(k,sizeof(int));
        c = (int *) calloc(m,sizeof(int));
        for (i=0; i<k; i++)
            b[i]=a[i];
        for (j=k; j<n; j++)
            c[j-1]=a[j];

        merge_sort (b, k);
        merge_sort (c, m);
        merge (b, c, a, k, m);
        free(b); free(c);
    }
}
```

```

void merge (int *a, int *b, int *c, int m, int n)
{
    int i, j, k, p;
    i = j = k = 0;
    do {
        if (a[i] < b[j]) {
            c[k]=a[i]; i=i+1;
        }
        else {
            c[k]=b[j]; j=j+1;
        }
        k++;
    } while ((i<m) && (j<n));
    if (i == m) {
        for (p=j; p<n; p++) { c[k]=b[p]; k++; }
    }
    else {
        for (p=i; p<m; p++) { c[k]=a[p]; k++; }
    }
}

```

```
main()
{
    int i, num;
    int a[ ] = {-56,23,43,-5,-3,0,123,-35,87,56,75,80};

    for (i=0;i<12;i++)
        printf ("%d ",a[i]);
    printf ("\n");

    merge_sort (a, 12);

    printf ("\nSorted list:");
    for (i=0;i<12;i++)
        printf (" %d", a[i]);
    printf ("\n");
}
```