

Programming & Data Structure

CS 11002

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Searching



Largest Element

(1)

Write a recursive C function that takes an `int` array and the number of elements present in the array as parameters, and returns the largest element of the array.

Iterative Function

```
int largestElement(int a[], int n) {  
    int i, largest = a[0];  
    for(i=1; i<n; ++i)  
        if(a[i] > largest) largest = a[i];  
    return largest;}  

```

Largest Element

(2)

Inductive Def $\Rightarrow$ Recursive Function

$$l(a[0 \cdots n-1]) = \begin{cases} a[0] & \text{if } n = 1, \\ \max(a[0], l(a[1 \cdots n-1])) & \text{if } n > 1. \end{cases}$$

```
int largestElement(int a[], int n) {  
    int temp;  
    if(n == 1) return a[0];  
    temp = largestElement(a+1, n-1);  
    return MAX(a[0], temp);  
} // largestFR.c
```

Reversal

(1)

Consider the problem of reversing a string or the digits of a number.

Ex: $\text{reverse}(274) = 472$; $\text{reverse}(\text{palin}) = \text{nilap}$.

Iterative Function

```
int reverse(int n){  
    int rev = 0;  
    while(n){  
        rev=10*rev+n%10; n /= 10;}  
    return rev;  
} // reverseIntF.c
```

Sequential / Linear Search

(1)

Write a function that takes n data stored in an array of type `int` and an integer as the *key*. The function sequentially searches the array for the *key*. If the *key* is present among the data in the array, then it returns the corresponding array *index*; otherwise it returns -1 .

Example

`a` =

8	3	11	5	7
---	---	----	---	---

key = 5: The function should return 3.

key = 6: The function should return -1 .

Sequential / Linear Search

(2)

Iterative C Function

```
#define NOTFOUND -1
int seqSearch(int a[], int n, int key){
    int i;
    for(i=0; i<n; ++i)
        if(a[i] == key) break;
    if(i == n) return NOTFOUND;
    return i;
} // seqSearchF.c
```

Sequential / Linear Search

(3)

Runtime

Best case: $\text{key} = a[0]$ and the **for loop** is executed only once.

Worst case: When $\text{key} = a[n - 1]$ or there is no match. The number of times the loop is executed is **n**, the number of data.

Sequential / Linear Search

(4)

Recursive C Function

```
#define NOTFOUND -1
int seqSearch(int a[],int n,int key){
    int temp;
    if(a[0] == key) return 0;
    if(n == 1) return NOTFOUND;
    temp = seqSearch(a+1, n-1, key);
    if(temp == NOTFOUND) return NOTFOUND;
    return temp + 1;
} // seqSearchFR.c
```

Sequential / Linear Search

(5)

Runtime

Best case: $\text{key} = a[0]$ and there is only one call.

Worst case: When $\text{key} = a[n - 1]$ or there is no match. The number of times the function is called is n , the number of data.

Space Usage: The number of stack frames (activation records) used by the recursive function is proportional to n in the worst case. On the contrary, the iterative function uses constant amount of extra space.

Binary Search

(1)

Search for a *key* can be made more efficient if the elements stored in the array are *sorted* in *ascending order* (or descending order).

Binary Search

(2)

key = 21

0	1	2	3	4	5	6	7	8	9
2	6	7	9	10	13	21	23	24	28

2	6	7	9	10	13	21	23	24	28
---	---	---	---	----	----	----	----	----	----

2	6	7	9	10	13	21	23	24	28
---	---	---	---	----	----	----	----	----	----

2	6	7	9	10	13	21	23	24	28
---	---	---	---	----	----	----	----	----	----

lo	hi	mid
----	----	-----

0 9 4

4 9 6

4 6 5

6 = 6



end of while loop

Binary Search

(3)

Iterative C Function

```
int binarySearch(int a[], int lo,
                  int hi, int key) {
    while(lo != hi) {
        int mid = (lo+hi)/2;
        if(key <= a[mid]) hi = mid;
        else lo = mid+1; }
    if(key == a[low]) return low;
    return -1;
} // binarySearchF.c
```

Binary Search

(4)

Runtime

The **while loop** is executed at most $\log_2 n + 1 = O(\log n)$ times.

For a large value of n , this gives a definite advantage over *sequential search* that executes the loop n times on a '*bad*' data set.

But in case of binary search, sorted data is required.

Binary Search

(5)

Inductive Definition

$$\text{bS}(a, l, h, k) = \begin{cases} l & \text{if } l = h \text{ \& } a[l] = k, \\ -1 & \text{if } l = h \text{ \& } a[l] \neq k, \\ \text{bS}(a, l, m, k) & \text{if } l < h \text{ \& } k \leq a[m] \\ \text{bS}(a, m + 1, h, k) & \text{if } l < h \text{ \& } k > a[m] \end{cases}$$

where l is the *low index*, h is the *high index*, k is the *key* to search, and $m = \frac{l+h}{2}$.

Binary Search

(6)

Recursive C Function

```
int binarySearch(int a[], int l, int h, int key){  
    if(l == h) {  
        if(a[l] == key) return l;  
        else return -1;}  
    if(l < h) {  
        int m = (l+h)/2;  
        if(a[m] >= key)  
            return binarySearch(data, l, m, key);  
        else return binarySearch(data, m+1, h, key);}  
} // binarySearchR.c
```

Binary Search

(7)

Runtime

For n elements, let T_n be the runtime of recursive binary search. Then

$$T_n = \begin{cases} c_0 & \text{if } n = 1 \\ T_{n/2} + c_1 & \text{if } n > 1, \end{cases}$$

where c_0 and c_1 are two constants.

Using the *iterative method*, we have

$$T_{n/2} = T_{n/2^2} + c_1, T_{n/2^2} = T_{n/2^3} + c_1, \dots, T_1 = c_0.$$

Hence, $T_n = c_0 + c_1 \log_2 n = O(\log_2 n)$.