

This exam contains 1 page and 4 problems.

You should *not* use your books, notes, or any calculator. Be *precise* in your answers. Intuitive justifications may not carry any marks, when you are asked to prove. All the *sub-parts* of a problem should be answered at *one place* only. On multiple attempts, *cross* any attempt that you do not want to be graded for.

There are no clarifications. In case of doubt, you can take a valid assumption, state that properly and continue.

1. (4 points) Let  $L_1$  and  $L_2$  be two infinite languages, defined over the alphabet  $\{a, b\}$ , satisfying  $L_1 \cap L_2 = \phi$  and  $L_1L_2 = L_2L_1$ . If such a language pair exists, give an example. If not, you must prove it.
2. (5 points) Construct an NFA to accept the regular expression  $b(((ba)^* + bbb)^* + a)^*b$ , such that the number of states are as minimum as possible. You should not use  $\epsilon$ -transitions. [Hint: the required states are  $\leq 5$ . You start losing marks as the number of states in your NFA increase beyond the required number.]
3. (5 points) Show that regular languages are closed under doubling. If language  $L$  is regular, then so also is the language  $L_2$  defined as

$$L_2 = \{\text{two } x \mid x \in L\}$$

where string doubling (two) is defined inductively as

$$\text{two } \epsilon = \epsilon$$

$$\text{two } ax = aa \cdot (\text{two } x)$$

4. (6 points) Use the pumping lemma to prove that the following language is non-regular.

$$L = \{b^p ab^q \mid p, q \geq 0, |p - 2q| = r^2, r > 0\}$$