

**CS21201 Discrete Structures**  
**Solutions to Practice Problems**

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**Propositional Logic**

1. Coding using Propositional Logic is as follows:

$$\begin{aligned} S1 & : \neg S \vee D \\ S2 & : \neg R \vee E \\ S3 & : (S \wedge \neg R) \vee (\neg S \wedge R) \\ S4 & : \neg S \vee E \\ S5 & : \neg R \vee D \\ G & : (\neg E \vee \neg D) \wedge (D \vee E) \end{aligned}$$

The goal should be easy enough to derive using the truth table method

2. **Solution I:** Assume that Ronaldo knows Mbappe means that they both know each other. Propositions used are as follows:

$$\begin{aligned} K & : \text{Mbappe knows Ronaldo (Ronaldo knows Mbappe)} \\ HL & : \text{Haaland likes the cookies} \\ MBP & : \text{Mbappe was on the pitch} \\ MEP & : \text{Messi was on the pitch} \end{aligned}$$

$$\text{Messi: } K \wedge HL$$

$$\text{Mbappe: } \neg K \wedge \neg MBP$$

$$\text{Haaland: } MBP \wedge MEP$$

Since one and only one of Haaland, Mbappe or Messi ate the cookies. We break the problem down into three cases depending on who ate the cookies:

- a. Messi: Both Mbappe and Haaland are telling the truth. But for that to be true,  $\neg MBP$  and  $MBP$  have to be true at the same time, **not possible**
- b. Haaland: Both Messi and Mbappe are telling the truth. But for that to be true,  $K$  and  $\neg K$  have to be true, **not possible**
- c. Mbappe: It can be seen that this is the only possible case. Since we can keep  $HL = 1$  and  $MEP = 1$ .

**Solution II:** Assume that Ronaldo knows Mbappe but Mbappe does not know Ronaldo. In that case, we design two predicates:

$$\text{KMR} : \text{Mbappe knows Ronaldo}$$

$$\text{KRM} : \text{Ronaldo knows Mbappe}$$

Messi:  $KRM \wedge HL$

Mbappe:  $\neg KMR \wedge \neg MBP$

Haaland:  $MBP \wedge MEP$

Since one and only one of Haaland, Mbappe or Messi ate the cookies. We break the problem down into three cases depending on who ate the cookies:

- a. Messi: Not possible, see **Solution I**
- b. Haaland: Both Messi and Mbappe are telling the truth. But for that to be true,  $KRM$  and  $\neg KMR$  have to be true (in addition to  $HL$  and  $\neg MBP$ ), **possible**
- c. Mbappe: Possible, see **Solution I**

Under this assumption, the identity of the thief cannot be ascertained.

3. Denote the  $NAND(x)$  function as  $N(x)$ . You can write  $\neg X$  as  $N(X, X)$ .
  - a.  $P \rightarrow Q \equiv \neg P \vee Q \equiv \neg(P \wedge \neg Q)$   
 $\equiv N(P, \neg Q)$  (Using De Morgan's Theorem)  
 $\equiv N(P, N(Q, Q))$
  - b.  $P \leftrightarrow Q \equiv (P \rightarrow Q) \wedge (Q \rightarrow P)$   
 $\equiv N(P, N(Q, Q)) \wedge N(Q, N(P, P))$

To simplify things a bit, choose:

$X = N(P, N(Q, Q))$  and  $Y = N(Q, N(P, P))$

$P \leftrightarrow Q \equiv X \wedge Y \equiv \neg(\neg X \vee \neg Y)$  (using De Morgan's Theorem)  
 $\equiv N(N(X, X), N(Y, Y))$

- c.  $P \oplus Q \equiv (P \vee Q) \wedge (\neg Q \vee \neg P)$   
 $X = (P \vee Q) \equiv N(\neg P, \neg Q) \equiv N(N(P, P), N(Q, Q))$   
 $Y = (\neg P \vee \neg Q) \equiv N(P, Q)$   
Repeat the solution to part (b).

4.
  - a. Notice that any formula constructed with  $\oplus$  will always be false when all propositions are false. However, consider  $\neg A$ . If  $A$  is false, then  $\neg A$  is true. Therefore, there are propositional logic formulae which can never be covered by  $\oplus$  and hence  $\{\oplus\}$  is not functionally complete.

- b. **This question is wrong.** It can be seen that any combination of  $\{\oplus, \wedge\}$  is always false whenever all propositions are false, since both  $A \oplus B$  and  $A \wedge B$  give false values when both A and B are false. Similar to part (a), consider  $\neg A$ , to continue.

5.

- a. It is sufficient to prove that  $(p \vee q) \wedge (\neg q \vee r) \wedge \neg(p \vee r)$  is a contradiction:

$$(p \vee q) \wedge (\neg q \vee r) \wedge \neg(p \vee r) \equiv (p \vee q) \wedge (\neg q \vee r) \wedge \neg p \wedge \neg r$$

By Disjunctive Syllogism,

$$(p \vee q) \wedge \neg p \rightarrow q \text{ and}$$

$$(\neg q \vee r) \wedge \neg r \rightarrow \neg q$$

Hence Proved

b.

- i. **This question has been changed,** consider the corrected problem as uploaded on the webpage:

$$p \rightarrow s \equiv \neg p \vee s \quad (1)$$

$$p \vee (q \wedge r) \quad (2)$$

$$s \vee (q \wedge r) \quad (3) \text{ Resolution Rule (2, 3)}$$

$$(s \vee q) \quad (4) \text{ Distribution (3)}$$

$$(s \vee r) \quad (5) \text{ Distribution (3)}$$

Hence proved

Note that in the original problem, the goal could not be derived from the statements.

- ii.  $(\neg q \vee r) \quad (1)$   
 $(p \vee q \vee t) \quad (2)$   
 $(p \vee r \vee t) \quad (3) \text{ Resolution Rule (2, 3)}$   
 $(\neg p \vee s) \quad (4)$   
 $(s \vee r \vee t) \quad (5) \text{ Resolution Rule (3, 4)}$   
 $\neg t \vee (s \wedge r) \quad (6)$   
 $(s \vee r) \vee (s \wedge r) \quad (7) \text{ Resolution Rule (5, 6)}$   
 $(s \vee r \vee s) \wedge (s \vee r \vee r) \quad (8) \text{ Distribution (7)}$   
 $(s \vee r) \quad (9) \text{ Simplification}$

6.

a. Trivial, consider an alternate form of  $A \rightarrow B$  as  $\neg A \vee B$ .

$Q = \neg(S_1 \wedge S_2 \wedge \dots \wedge S_n) \vee G$  is a tautology.

Therefore,  $\neg Q$  is a contradiction.

$$\neg Q = \neg(\neg(S_1 \wedge S_2 \wedge \dots \wedge S_n) \vee G) = (S_1 \wedge S_2 \wedge \dots \wedge S_n) \wedge \neg G$$

b. S1:  $(p \rightarrow q) \rightarrow q \equiv \neg(\neg p \vee q) \vee q \equiv (p \wedge \neg q) \vee q \equiv p \vee q$

S2:  $(p \rightarrow p) \rightarrow r \equiv \neg(\neg p \vee p) \vee r \equiv r$

Hence the goal directly follows from S2.

7. Consider the propositions for the problem:

WF : Water altar contains the mysteries of fire

EF : Earth altar contains the mysteries of fire

EW : Earth altar holds the secrets of water

AW : Air altar holds the secrets of water

FE : Fire altar holds the energy of earth

AE : Air altar holds the energy of earth

AF : Air altar holds the mysteries of fire

AA : Air altar holds the strength of air

EA : Earth altar holds the strength of air

FW : Fire altar holds the secrets of water

FA : Fire altar holds the strength of air

FE : Fire altar holds the energy of earth

a.  $\neg(WF \wedge EF) \equiv \neg WF \vee \neg EF$

b.  $\neg(\neg EW \wedge \neg AW) \equiv EW \vee AW$

c.  $\neg(FE \wedge (\neg AW \wedge \neg AE \wedge \neg AF \wedge \neg AA))$

$$\neg FE \vee \neg(\neg AW \wedge \neg AE \wedge \neg AF \wedge \neg AA) \equiv \neg FE \vee AW \vee AE \vee AF \vee AA$$

d.  $\neg(\neg EA \wedge (\neg FW \wedge \neg FE \wedge \neg FF \wedge \neg FA))$

$$EA \vee FW \vee FE \vee FF \vee FA$$

From the second statement, it is clear that either the earth altar or the air altar holds the power of water.

**Earth altar holds the power of water**

Consider  $EW = 1, FA = 1, AE = 1, WF = 1$

**Air altar holds the power of water**

Consider  $AW = 1, FA = 1, AE = 1, WF = 1$

Therefore, both the Earth and Air altars can hold the power of water.

8.  $(r \rightarrow \neg q)$  (1)  
 $(p \rightarrow q)$  (2)  
 $(q \rightarrow \neg r)$  (3) Contrapositive (1)  
 $(p \rightarrow \neg r)$  (4) Transitive Implication  
 $(r \rightarrow \neg p)$  (5) Contrapositive (4)

It is clear from the above derivation that  $\rightarrow$  holds. However,  $\leftarrow$  does not hold.  
 Choose the values of p, q and r as  $r = T$ ,  $p = F$  and  $q = T$ .

- 9.
- a.  $(Smoke \rightarrow Fire) \rightarrow (\neg Smoke \rightarrow \neg Fire)$  (1)  
 $\equiv (Smoke \rightarrow Fire) \rightarrow (Fire \rightarrow Smoke)$  (2) Contrapositive (1)  
**Valid:** No ( $Smoke = F, Fire = T$ )  
**Satisfiable:** Yes ( $Smoke = T, Fire = T$ )
- b.  $((Smoke \wedge Heat) \rightarrow Fire) \leftrightarrow ((Smoke \rightarrow Fire) \vee (Heat \rightarrow Fire))$   
 $((Smoke \wedge Heat) \rightarrow Fire) \equiv \neg Smoke \vee \neg Heat \vee Fire$   
 $\equiv \neg Smoke \vee Fire \vee \neg Heat \vee Fire \equiv (Smoke \rightarrow Fire) \vee (Heat \rightarrow Fire)$   
**Valid:** Yes  
**Satisfiable:** Yes
- c.  $Big \vee Dumb \vee (Big \rightarrow Dumb)$   
 $\equiv Big \vee Dumb \vee (\neg Big \vee Dumb) \equiv Dumb$   
**Valid:** No ( $Dumb = F$ )  
**Satisfiable:** Yes ( $Dumb = T$ )
- d.  $(Big \wedge Dumb) \vee \neg Dumb \equiv (Big \vee \neg Dumb) \wedge (Dumb \vee \neg Dumb)$   
 $(Big \vee \neg Dumb)$   
**Valid:** No ( $Big = F, Dumb = T$ )  
**Satisfiable:** Yes (all other truth values)
- e.  $(Smoke \rightarrow Fire) \rightarrow ((Smoke \wedge Heat) \rightarrow Fire)$   
 $Heat = T \Rightarrow (Smoke \rightarrow Fire) \rightarrow (Smoke \rightarrow Fire) = T$   
 $Heat = F \Rightarrow (Smoke \rightarrow Fire) \rightarrow T = T$   
**Valid:** Yes  
**Satisfiable:** Yes

**CS21201 Discrete Structures**  
**Solutions to Tutorial Problems**

**Propositional Logic**

1.

Tutorial - 2  
Propositional Logic

**Question - 1**

(a) Propositions used are as follows:

p: Swapna wrote the paper  
t: Swapna typed the answers  
h: Swapna has a camera  
c: Swapna completed in time

(b)

$$\begin{array}{l} S1: p \vee t \\ S2: (p \vee h) \rightarrow \neg c \\ S3: \neg c \\ \hline G: p \end{array}$$

(c)  $p = F, c = F, h = F, t = T$  S1, S2 and S3 are true but G is false hence  $S1 \wedge S2 \wedge S3 \rightarrow G$  is not a Tautology.

2.

*Solution* The first deduction goes as follows.

$\frac{\neg(s \vee u)}{\therefore \neg s, \neg u}$	$\frac{\neg u}{t \rightarrow u} \therefore \neg t$	$\frac{r \rightarrow (s \vee t) \quad \neg s, \neg t}{\therefore \neg r}$	$\frac{(\neg p \vee q) \rightarrow r \quad \neg r}{\therefore p \wedge \neg q}$	$\frac{q \leftrightarrow v \quad \neg q}{\therefore \neg v}$
$\frac{(v \wedge \neg w) \vee (\neg v \wedge w) \rightarrow \neg p \quad p}{\therefore (\neg v \vee w) \wedge (v \vee \neg w)}$	$\frac{(\neg v \vee w) \wedge (v \vee \neg w) \quad \neg v}{\therefore \neg w}$			

3.

a.

*Solution* We may use the following propositions.

- rob : The murder was done for robbery.
- pol : The murder was a political assassination.
- wom : The murder was for a woman.
- tak : Something was taken from the murderer's place.
- imm : The assassin left immediately after work done.
- trc : The assassin left tracks all over the room.

b.

- Solution*
- (1)  $\neg \text{rob} \rightarrow \text{pol} \vee \text{wom}$
  - (2)  $\text{rob} \rightarrow \text{tak}$
  - (3)  $\neg \text{tak}$
  - (4)  $\text{pol} \rightarrow \text{imm}$
  - (5)  $\text{trc}$
  - (6)  $\text{trc} \rightarrow \neg \text{imm}$

c.

*Solution*

$$\begin{array}{l} \text{rob} \rightarrow \text{tak} \\ \neg \text{tak} \\ \hline \therefore \neg \text{rob} \quad (\text{Modus Tollens}) \\ \\ \neg \text{rob} \\ \neg \text{rob} \rightarrow \text{pol} \vee \text{wom} \\ \hline \therefore \text{pol} \vee \text{wom} \quad (\text{Modus Ponens}) \\ \\ \text{trc} \\ \text{trc} \rightarrow \neg \text{imm} \\ \hline \therefore \neg \text{imm} \quad (\text{Modus Ponens}) \\ \\ \neg \text{imm} \\ \text{pol} \rightarrow \text{imm} \\ \hline \therefore \neg \text{pol} \quad (\text{Modus Tollens}) \\ \\ \text{pol} \vee \text{wom} \\ \neg \text{pol} \\ \hline \therefore \text{wom} \quad (\text{Disjunctive Syllogism}) \end{array}$$

d. The murder was done for a woman