

Practice Problems: Fast Fourier Transform and Dynamic Programming

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Submit the solutions of the questions marked (★) in PDF format generated using Latex by **April 11, 2026**.

1. In the inverse DFT problem, we are given evaluation of a polynomial of degree at most n at every n -th root of unity. The goal is to find the coefficient representation of the polynomial. Design a $\mathcal{O}(n \log n)$ time algorithm for the problem.
2. (★) Design a $\mathcal{O}(n \log n)$ time algorithm to multiply two n bit integers.
3. [CLRS] Consider two sets A and B , each having n integers in the range from 0 to $10n$. The Cartesian sum of A and B is defined by $C = \{x + y : x \in A, y \in B\}$. The integers in C lie in the range from 0 to $20n$. Show how, in $\mathcal{O}(n \log n)$ time, to find the elements of C and the number of times each element of C is realized as a sum of elements in A and B .
4. A set of arcs \mathcal{F} in a directed graph $\mathcal{G} = (\mathcal{V}, \mathcal{A})$ is called a directed feedback arc set of \mathcal{G} if $\mathcal{G}[\mathcal{V}, \mathcal{A} \setminus \mathcal{F}]$ is acyclic. Show that we can compute a minimum cardinality directed arc set of any directed graph in $\mathcal{O}^*(2^n)$ time, where n is the number of vertices of \mathcal{G} .
5. (★) The cut-width of a vertex ordering π of a graph $\mathcal{G} = (\mathcal{V}, \mathcal{E})$ is

$$\max_{v \in \mathcal{V}} |\{w, x\} \in \mathcal{E} : \pi(w) \leq \pi(v) < \pi(x)\}|.$$

The cut-width of a graph is the minimum cut-width taken over all orderings of its vertices. Show that we can compute the cut-width of a graph in time $\mathcal{O}^*(2^n)$ time, where n is the number of vertices of \mathcal{G} .

6. The domatic number of a graph \mathcal{G} is the minimum integer k such that $\mathcal{V}[\mathcal{G}]$ can be partitioned into k sets $\mathcal{V}_1, \dots, \mathcal{V}_k$ such that each \mathcal{V}_i is a dominating set of \mathcal{G} . Show how we can compute the domatic number of any graph in $\mathcal{O}^*(3^n)$ time, where n is the number of vertices of \mathcal{G} .
7. In the EXACT SAT problem, we are given a CNF formula. The task is to check if there exists an assignment of its variables so that every clause has exactly one literal set to TRUE. Show that there is a $\mathcal{O}^*(2^m)$ time algorithm for EXACT SAT, where m is the number of clauses in the formula.