# Tutorial 2: CS21003 Algorithms I 

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1. Let $A$ be a $n \times n$ two-dimensional array, in which all the rows and all the columns are sorted in ascending order from smaller to larger indices. Given a key $x$, your task is to find out whether $x$ is an element of $A$. Develop an efficient algorithm for the same and analyse the complexity.
2. Let $X=\left\{x_{1}, x_{2}, \ldots, x_{n}\right\}$ be an array of $n$ positive integers and $a$ is an integer. Propose an efficient algorithm to determine whether there are two elements in $X$ whose sum is exactly a. Derive the time-complexity of the proposed algorithm.
3. Analyze the complexity of the algorithm:
```
func(n){
    k=0, y=2
    while (kin){
                y=y*y,k++
    }
    return y
}
```

4. Let $A$ be an array of $n$ integers. In this exercise, we will show that the expected number of comparisons made by the randomized quick sort algorithm to sort $A$ is $\mathcal{O}(n \log n)$.
(a) Write a pseudo-code for randomized quick sort.
(b) Write a pseudo-code for the partition algorithm where the pivot element is involved in every comparisons.
(c) For every $1 \leqslant i<j \leqslant n$, define an indicator random variable $Z_{i, j}$ for the event that $\mathcal{A}[i]$ and $A[j]$ are compared with each other by the algorithm.
(d) Compute $\mathbb{E}\left[Z_{i, j}\right]$
(e) Define a random variable $Z$ to be the number of pairs of elements compared by the randomized quick sort algorithm to sort $A$. Write $Z$ in terms of $Z_{i, j}, 1 \leqslant i<j \leqslant n$
(f) Compute $\mathbb{E}[Z]$
