# INTRODUCTION TO RECURSIVE FORMULATIONS FOR ALGORITHM DESIGN: I



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## **Overview**

- □ <u>Algorithms</u> and Programs
- Pseudo-Code
- Algorithms + Data Structures = Programs
- Initial Solutions + Analysis + Solution Refinement + Data Structures = Final Algorithm
- Use of <u>Recursive Definitions</u> as Initial Solutions
- Recurrence Equations for Proofs and Analysis
- Solution Refinement through Recursion Transformation and Traversal

Space

Data Structures for saving past computation for future use Time 2 complexity

#### Sample Problems:

- 1. Finding the Largest
- 2. Largest and Smallest
- 3. Largest and Second Largest
- 4. Fibonacci Numbers
- 5. Searching for an element in an ordered / unordered List
- 6. Sorting
- 7. Pattern Matching
- 8. Permutations and Combinations
- 9. Layout and Routing

10. Shortest Paths

#### First Problem: Largest of a Set of Numbers

Sequential Comparison L= ZZ1, Z2, ..., Rn Z x is an integer max (L  $z = z_1, z_2, \dots, z_n z_1$ z = 1 return  $(z_1)$ - 54 y) return (x1) T(n) = -T(n-1) + 1, n > 1

25,8,3,1,2,6,12212 ₹8,3,1,2,6,123 12 3 \$ 3, 1, 2, 6, 12 3, 2 3 25, 2, 6, 12 3, 12 6 companyions 52,6,123

### **Finding Largest: Recursive Formulation**

maz2(L) {x<sub>1</sub>, x<sub>2</sub>,..., x<sub>n</sub>} |L| = 1 return (x<sub>1</sub>) lit Linto 2 non-empty sets L1, L2 = marca(L)mand potwin (y,) etse return (y2) boof by INDUCTION Correct ners T(n) = T(K) + T(n-K) n= 0, if n=1 = (n - 1)



#### Largest: Analysis

**Proof of Correctness** By induction Base, condition n=1 Inductive condition correctly for all n<no we prove inductively it is true for no+1 True for all of 2 1

# **Complexity Analysis** T(n) = T(k) + T(n-k) + 1for n>1 =0, $\eta = 1$ T(x) = x - 1T(n) = (K-1) + (n-K-1)- 71-4

#### **Comparison Tournament**



### 2<sup>nd</sup> Problem: Largest and Smallest

Sequential Comparison maxmin (L) let L= 2x1, 22, ..., Xnz  $f_{1}(L) = 1$  return  $(\langle z_1, z_1 \rangle)$ 1 = L- ZZZ  $\langle y_{1}, y_{2} \rangle = maximin(L')$  $\Psi(x_1, y_1) \text{ then } m_1 = x_1$ else  $m_1 = y_1$  $3f(x_1 \leq y_2)$  then  $m_2 = x_1$ else Ma= Ha, return (<m, m2>) T(n) = 0, M = 1= T(m-1) + 2, m>1

25,8,3,1,2,6,12 3 (12,1) 5 28,3,1,2,6,12% 8 3, 1, 2, 6, 123 人」人(2,1) え」,2,6,123 1 52,6,124 2 \$ 3,6,12 3, (12,12)

### Largest and Smallest: Recursive Definition

maxmin 2 (L let L= {x1, x2, ..., xn}  $\mathfrak{g}[L] = 1$ , return  $(\langle \mathfrak{z}_{1}, \mathfrak{z}_{1} \rangle)$ |L| = 2, if (2/722)return ({x1, 22}) else return ({a2, x1>) split Linto 2 non-empty sets 1  $\langle y_1, y_2 \rangle = maximun a. (4)$  $\langle z_1, z_2 \rangle = maximin a($ Lg -  $(y_1 > Z_1)$   $m_1 = y_1$  else  $m_1 = Z_1$  $g_{f} (\dot{y}_{2} \leq z_{2}) m_{f} = \dot{y}_{2} else m_{f} = z_{2}$ return (<m, ma>)



### Largest & Smallest: Choice of Split



N × (n-2) T(n) = 0, y = 1= 1, y = 2= T(K) + T(n-K) + 2, if 372If we choose \$1=1 T(1) + T(n-1) + 2= 2TV - 3 (base case of) 5(2) = 1choose &= 2, = T(2) + T(n-2) + 2= T(2) + T(2) + T(n-4) + 2 + 231/2-

#### **Tournament Based Analysis & Design**



optimal spl (2) = even then split into even, even parts 121 = ord then split into even, odd 6 elements シース 1→ 3,3 → 4,2

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Thank you

Any Questions?