

flex - Fast Lexical Analyzer Generator

We can use **flex**<sup>a</sup> to automatically generate the lexical analyzer/scanner for the lexical atoms of a language.

<sup>a</sup>The original version was known as lex.



The input to the flex program (known as flex compiler) is a set of patterns or specification of the tokens of the source language. Actions corresponding to different matches can also be specified.





The flex software compiles the regular expression specification to a DFA and implements it as a C program with the action codes properly embedded in it.



The output of flex is a C file lex.yy.c with the function int yylex(void). This function, when called, returns a token. Attribute information of a token is available through the global variable yylval.



#### **flex** Specification

The definition has two parts: The portion within the pair of special parentheses %{ and %} is copied verbatim in lex.yy.c. Similarly the user code is also copied. The other part of the definition contains regular name definitions, start conditions and other options.

## **flex** Specification

The rules are specified as <pattern> <action> pairs. The patterns are extended regular expressions corresponding to different tokens. The corresponding actions are given as C code. An action is taken at the end of the match.

```
An Example: lex.l
/* A scanner for a toy language:
   ./lex_yacc/lex.l
*/
%{ // Copied verbatim
#include <string.h>
#include <stdlib.h>
#include "y.tab.h" // tokens are defined
int line_count = 0; yylType yylval;
%}
```

```
%option noyywrap
%x CMNT
DELIM ([ \t])
WHITESPACES ({DELIM}+)
POSI_DECI ([1-9])
DECI (0|{POSI_DECI})
          (0|{POSI_DECI}{DECI}*)
DECIMAL
LOWER
       ([a-z])
LETTER ({LOWER}|[:upper:]|_)
IDENTIFIER ({LETTER}({LETTER})*)
%%
```

```
"/*"
             {BEGIN CMNT;}
<CMNT>. {;}
<CMNT>\n {++line_count;}
<CMNT>"*/" {BEGIN INITIAL;}
n
              ++line_count;
              return (int)'\n';
             }
"\"".*"\""
yylval.string=(char *)malloc((yyleng+1)*(sizeof(dhar)));
 strncpy(yylval.string, yytext+1, yyleng-2);
 yylval.string[yyleng-2]='\0';
return STRNG;
             ł
```

{WHITESPACES]	} { ; }
if	<pre>{return IF;}</pre>
else	<pre>{return ELSE; }</pre>
while	<pre>{ return WHILE; }</pre>
for	{ return FOR; }
int	<pre>{ return INT; }</pre>
$\backslash$ (	<pre>{ return (int)'('; }</pre>
$\backslash$ )	<pre>{ return (int)')'; }</pre>
\{	<pre>{ return (int)'{'; }</pre>
\}	<pre>{ return (int)'}'; }</pre>
;	<pre>{ return (int)';'; }</pre>
,	<pre>{ return (int)','; }</pre>
=	<pre>{ return (int)'='; }</pre>
"<"	<pre>{ return (int)'&lt;'; }</pre>

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```
return BIN_OP;
              }
{IDENTIFIER} {
  yylval.string=(char *)malloc((yyleng+1)*(sizeof(char)));
  strncpy(yylval.string, yytext, yyleng + 1);
  return ID;
{DECIMAL}
               yylval.integer = atoi(yytext);
               return NUM;
             { ; }
%%
  The function yywrap() int yywrap(){return 1;} */
```

			y.tab.h
#ifndef	_Y_TAB_	_H	
#define	_Y_TAB_	_H	
#define	IF	300	
#define	ELSE	301	
#define	WHILE	302	
#define	FOR	303	
#define	INT	304	
#define	ID	306	
#define	NUM	307	
#define	STRNG	308	
<pre>#define</pre>	BIN_OP	310	

```
int yylex(void);
typedef union {
    char *string;
    int integer;
} yylType;
```

#endif

```
myLex.c++
#include <iostream>
using namespace std;
#include <stdlib.h>
#include "y.tab.h"
extern yylType yylval;
int main() // myLex.c++
{
     int s;
     while((s=yylex()))
     switch(s) {
      case '\n': cout << "\n";</pre>
```

		break;
case '(	)	cout << "<(> ";
		break;
case ')	':	cout << "<)> ";
		break;
<pre>case '{</pre>	':	cout << "<{> ";
		break;
<pre>case '}</pre>	':	cout << "<}> ";
		break;
case ';	:	cout << "<;> ";
		break;
case ',	:	cout << "<,> ";
		break;
case '=	:	cout << "<=> ";

```
break;
case '<' : cout << "<<> ";
           break;
case BIN_OP : cout << "<BIN_OP," << (char) yylval.int
              break;
case IF : cout << "<if>";
           break;
case ELSE : cout << "<else> ";
            break;
case WHILE : cout << "<while> ";
             break;
case FOR : cout << "<for> ";
           break;
case INT : cout << "<int> ";
```

break;

```
case ID : cout << "<ID," << yylval.string << "> ";
    free (yylval.string);
    break;
```

```
case NUM : cout << "<NUM," << yylval.integer << ">;
    break;
```

```
case STRNG : cout << "<STRNG," << yylval.string << ">
free (yylval.string) ;
break;
```

```
default: ;
```

```
return 0;
```

}

input /\* Scanner input \*/ int main() { int n, fact, i; scanf("%d", &n); fact=1; for(i=1; i<=n; ++i) fact = fact\*i;</pre> printf("%d! = %d\n", n, fact); return 0; } /\* End \*/

How to Compile?

```
A Simple Makefile
```

```
objfiles = myLex.o lex.yy.o
```

```
a.out : $(objfiles)
    c++ $(objfiles)
```

```
myLex.o : myLex.c++ y.tab.h
    c++ -c -Wall myLex.c++
```

```
lex.yy.o : lex.yy.c
c++ -c lex.yy.c
```

```
lex.yy.c : lex.l y.tab.h
```

```
flex lex.l
clean :
  rm a.out lex.yy.c $(objfiles)
```





There is a problem with the rule for comment.
It will give wrong result for the following
comment (input4)
/\* There is a string "within \*/ this
 comment" and it will not work \*/

Try with the rule in lex4.1



When the function yylex() is called, it reads data from the input file (by default it is stdin) and returns a token. At the end-of-file (EOF) the function returns zero (0).

#### Input/Output File Names

The name of the input file to the scanner can be specified through the global file pointer (FILE \*) yyin i.e. yyin = fopen("flex.spec", "r");. Similarly, the output file of the scanner<sup>a</sup> can also be supplied through the file pointer yyout i.e.

```
<sup>a</sup>Note that it will affect the output of ECHO and not change the output of printf() etc.
```

## Pattern and Action

A pattern must start from the first column of the line and the action must start from the same line. Action corresponding to a pattern may be empty. The special parenthesis %{ and %} of the definition part should also start from the first column.

#### Matching Rules

If more than one pattern matches, then the longest match is taken. If both matches are of same length, then the earlier rule will take precedence. In our example if, else, while, for, int are specified above the identifier. If we change the order, the keywords will not be identified separately.

### Matching Rules

In our example '<=', is treated as two symbols '<' and '='. But if there is a rule for '<=', a single token can be generated.

# Lexeme

If there is a match with the input text, the matched lexeme is pointed by the char pointer yytext (global) and its length is available in the variable yyleng. The yytext can be defined to be an array by using %array in the definition section. The array size will be determined by YYLMAX.



The command ECHO copies the text of yytext to the output stream of the scanner<sup>a</sup>. The default rule is to echo any unmatched character to the output stream. (lex1.1)

<sup>a</sup>By default it is **stdout** and it can be changed by the file pointer **yyout**.

## yymore(), yyless(int)

- yymore(): the current lexeme remains in the yytext. The next lexme is appended to it.
- yyless(n): the scanner puts back all but the first n characters in the input buffer. Next time they will be rescanned. Both
  yytext and yyleng (n now) are properly modified. (lex1.1)

# input(), unput(c)

input(): reads the next character from the
input stream.
unput(c): puts the character c back in the
input stream<sup>a</sup>.

<sup>a</sup>flex manual: An important potential problem when using unput() is that if you are using %pointer (the default), a call to unput() destroys the contents of yytext, starting with its rightmost character and devouring one character to the left with each call. If you need the value of yytext preserved after a call to unput() (as in the above example), you must either first copy it elsewhere, or build your scanner using %array instead.
There is a mechanism to put a conditional guard for a rule. Let the condition for a rule be <C>. It is written as <C>P {action} The scanner will try to match the input with the pattern P, only if its start condition is C. In the example we have three rules guarded by the start condition <CMNT> used to remove the comment of a C program without any action.

A start condition is activated by a BEGIN action. In our example the start condition begins (BEGIN CMNT) after the scanner sees the starting of a comment (/\*).

The state of the scanner when no other start condition is active is called INITIAL. All rules without any start condition or with the start condition <INITIAL> are active at this state. The command BEGIN(0) or BEGIN(INITIAL) brings the scanner from any other state to this state.

In our example the scanner comes back to **INITIAL** after consuming the comment.

- Start conditions can be defined in the definition part of the specification using %s or %x. The start condition specified by %x is called an exclusive start condition. When such a start is active, only the rules guarded by it are activated. On the other hand %s specifies a an inclusive start condition.
- In our example, only three rules are active after the scanner sees the beginning of a comment (/\*).

```
The scope of start conditions may be nested
and the start conditions can be stacked. Three
routines are of interest:
void yy_push_state(int new_state),
void yy_pop_state() and
int yy_top_state().
```

# yywrap()

Once the scanner receives the EOF indication from the YY\_INPUT macro (zero), the scanner calls the function yywrap(). If there is no other input file, the function returns true and the function yylex() returns zero (0). But if there is another input file, yywrap() opens it, sets the file pointer yyin to it and returns false. The scanner starts consuming input from the new file.

%option noyywrap

This option in the definition makes scanner behave as if yywrap() has returned true. No yywrap() function is required.



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#### Architecture of Flex Scanner

The flex compiler includes a fixed program in its scanner. This program simulates the DFA. It reads the input and simulates the state transition using the state transition table constructed from the flex specification. Other parts of the scanner are as follows:

#### Architecture of Flex Scanner

- The transition table, start state and the final states this comes from the construction of the DFA.
- Declarations and functions given in the definition and user code of the specification file - these are copied verbatim.

#### Architecture of Flex Scanner

Actions specified in the rules corresponding to different patterns are put in such a way that the simulator can initiate them when a pattern is matched (in the corresponding final state).



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