## Modeling of Finite State Machines

Debdeep Mukhopadhyay

# Definition

- 5 Tuple:  $(Q, \Sigma, \delta, q_0, F)$
- Q: Finite set of states
- Σ: Finite set of alphabets
- δ: Transition function
  - $-Q\chi \Sigma \rightarrow Q$
- q<sub>0</sub> is the start state
- F is a set of accept states. They are also called final states.

#### Some Examples



What does this FSM do?

It accepts the empty string or any string that ends with 0

These set of strings which takes the FSM to its accepting states are often called **language** of the automaton.



• Accepts strings that starts and ends with the same bits.



- FSM accepts if the running sum of the input strings is a multiple of 3.
- RESET symbol resets the running sum to 0.

# Designing FSMs

- Its an art.
- Pretend to be an FSM and imagine the strings are coming one by one.
- Remember that there are finite states.
- So, you cannot store the entire string, but only crucial information.
- Also, you do not know when the string ends, so you should always be ready with an answer.

## Example

- Design a FSM which accepts 0,1 strings which has an odd number of 1's.
- You require to remember whether there are odd 1's so far or even 1's so far.



## Example

- Design a FSM that accepts strings that contain 001 as substrings.
- There are 4 possibilities
  - No string
  - seen a 0
  - seen a 00
  - seen a 001

#### Answer



- Note that their may be cases where design of FSMS are not possible.
- Like design an FSM for strings which has the same number of 0's and 1's.

#### How to model such FSMs?



#### Simple Model of FSM

#### Mealy Machine/Moore Machine





#### Modeling FSMs using Verilog

#### Issues

- State Encoding
  - sequential
  - gray
  - Johnson
  - one-hot

# **Encoding Formats**

No	Sequential	Gray	Johnson	One-hot
0	000	000	0000	0000001
1	001	001	0001	00000010
2	010	011	0011	00000100
3	011	010	0111	00001000
4	100	110	1111	00010000
5	101	111	1110	00100000
6	110	101	1100	01000000
7	111	100	1000	10000000

## Comments on the coding styles

- Binary: Good for arithmetic operations. But may have more transitions, leading to more power consumptions. Also prone to error during the state transitions.
- Gray: Good as they reduce the transitions, and hence consume less dynamic power. Also, can be handy in detecting state transition errors.

# **Coding Styles**

- Johnson: Also there is one bit change, and can be useful in detecting errors during transitions. More bits are required, increases linearly with the number of states. There are unused states, so we require either explicit asynchronous reset or recovery from illegal states (even more hardware!)
- One-hot: yet another low power coding style, requires more no of bits. Useful for describing bus protocols.

#### Good and Bad FSM



#### **FSM State Diagram**

```
always@(posedge Clock)
begin
parameter ST_Read=0,ST_Write=1,ST_Delay=3;
 integer state;
 case(state)
    ST_Read:
     begin
        Read=1;
        Write=0;
        State=ST Write;
     end
```

```
ST Write:
  begin
    Read=0;
    Write=1;
    if(SlowRam) State=ST Delay;
    else State=ST Read;
  end
```

ST Delay: begin Read=0; Write=0; State=ST\_Read; end endcase end

# Why Bad?

- No reset. There are unused states in the FSM.
- Read and Write output assignments also infer an extra flip-flop.
- No default, latch is inferred.
- There is feedback logic.

# Good verilog

always @(posedge Clock) begin if(Reset) CurrentState=ST Read; else CurrentState=NextState; end

# Good verilog

```
always@(CurrentState or SlowRAM)
begin
 case(CurrentState)
   ST Read:
     begin
       Read=1; Write=0;
       NextState=ST Write;
     end
```

## Good Verilog

ST\_Write: begin Read=0; Write=1; if(SlowRAM) NextState=ST\_Delay; else NextState=ST\_Read; end

# Good Verilog

```
ST Delay:
  begin
    Read=0; Write=0; NextState=ST Read;
  end
default:
  begin
    Read=0; Write=0; NextState=ST Read;
  end
 endcase
end
```

#### One Bad and four good FSMs



always @(posedge Clock or posedge Reset) begin

if(Reset) begin

```
Y=1;
STATE=ST0;
end
```

```
else
case(STATE)
ST0: begin Y=1; STATE=ST1; end
ST1: begin Y=2;
if(Control) STATE=ST2;
else STATE=ST3;
ST2: begin Y=3; STATE=ST3; end
ST3: begin Y=4; STATE=ST0; end
endcase
end
```

Output Y is assigned under synchronous always block so extra three latches inferred.

#### Good FSMs

- Separate CS, NS and OL
- Combined CS and NS. Separate OL
- Combined NS and OL. Separate CS

# Next State (NS)

```
always @(control or currentstate)
begin
  NextState=ST0;
  case(currentstate)
   ST0: begin
          NextState=ST1;
        end
   ST1: begin ...
   ST3:
          NextState=ST0;
endcase
end
```

# Current State (CS)

always @(posedge Clk or posedge reset) begin if(Reset) currentstate=ST0; else currentstate=Nextstate; end

# Output Logic (OL)

always @(Currentstate) begin case(Currentstate) ST0: Y=1; ST1: Y=2; ST2: Y=3; ST3: Y=4; end

#### CS+NS

```
always@(posedge Clock or posedge reset)
begin
 if(Reset)
       State=ST0;
  else
     case(STATE)
                                        default not required
         ST0: State=ST1;
                                       as it is in edge triggered
         ST1: if(Control) ...
                                         always statement
         ST2: ...
         ST3: STATE=ST0;
      endcase
  end
```

#### CS+NS



### NS+OL

```
always @(Control or Currentstate)
begin
  case(Currentstate)
    ST0: begin
           Y=1;
           NextState=ST1;
          end
    ST1: ...
    ST2: ...
    ST3: ...
    default: ...
   endcase
 end
```

# NS+OL

always @(posedge clock or posedge reset) begin if(reset) Currentstate=ST0; else Currentstate=NextState; end