

## CS39002 OPERATING SYSTEMS LABORATORY

### SPRING 2025

LAB ASSIGNMENT: 10  
DATE: 02-APRIL -2025

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### Demand Paging with Page Replacement

In this assignment, we continue with the embedded device of LA9. The specifications of this system are reproduced below. The input is to be read from *search.txt* which is organized in the same format as in LA9.

Total memory = 64 MB  
Page/Frame size = 4 KB  
OS usage = 16 MB (higher side)  
User space = 48 MB  
Number of user frames = 12288 (numbered 0, 1, 2, ..., 12287)  
Number of processes =  $n$  (in the range [50, 500], PID's are numbered 0, 1, 2, ...,  $n - 1$ )  
Number of binary searches per process =  $m$  (in the range [10, 100])  
Virtual memory size = 2048 pages (this is also the page-table size of each process)  
Number of essential pages = 10 (page numbers 0, 1, 2, ..., 9)  
Number of pages for storing the array  $A$  = 2038 (page numbers 10, 11, 12, ..., 2047)  
Size of each page-table entry = 16 bits (unsigned short int)  
14-bit frame addresses (bits 0 – 13), Bit 15 is the valid bit, Bit 14 to be used as the reference bit

In LA9, processes are swapped out when memory is full. In this assignment, we keep all the processes in memory. When memory is *full*, we replace valid pages by new pages. Approximate LRU policy is used for page replacement.

#### Global frame allocation

Maintain a list of free frames (FFLIST). Each entry in the FFLIST consists of (1) a frame number, (2) the last owner (PID) of the frame (–1 if none), and (3) the page number of the last owner stored in the frame (–1 if no last owner). In LA9, the FFLIST consisted only of free-frame numbers, and was implemented as a FIFO queue. Here, you need to make passes through the FFLIST, so it is technically no longer a queue. Use an array of triples (frame number, last owner, last owner's page number) as the FFLIST. Also maintain a counter NFF (number of free frames). Make your own implementation.

To start with, add all the 12288 user frames to the FFLIST. When you start the  $n$  processes, allocate 10 essential frames to each from the FFLIST. Subsequently, free frames are allocated on demand. However, we will never allow NFF (the number of free frames) to fall below NFFMIN = 1000. So long as we have  $NFF > NFFMIN$ , any process requesting a frame is granted a free frame. You do not need to check how many frames are already owned by the process. In our application, processes carry out binary searches in a round-robin fashion, so this frame-allocation strategy will ensure more or less equal distribution of frames to the processes. Although the sizes of  $A$  for different processes vary within a factor of 2, a proportional frame-preallocation scheme based on the sizes of  $A$  may be more appropriate, but you do not need to do this.

When  $NFF = NFFMIN$ , and a process requests a frame, a free frame in the FFLIST is no longer allocated straightaway to the process. On the contrary, a free frame in the FFLIST is swapped with a frame currently allocated to the process. This situation is described in detail in the next section.

When a process exits (after completing its  $m$  searches), all the frames held by that process at that time are returned to the FFLIST. These frames join the FFLIST with no last owner, because the exiting process will never try to use/reclaim these (or other) frames. Free frames released by an exiting process will be allocated to the existing processes on demand.

### Local page replacement

A process  $i$ , in each iteration of its binary-search loop, makes a memory access of  $A[M]$ . Assume that this element is stored in logical page  $p$  of the process. If the page table of process  $i$  indicates the valid bit to be 1 for the  $p$ -th page, that page is already residing in a frame, so the process proceeds to the next iteration of the binary-search loop. In the other case, that is, the valid bit is 0, this memory access encounters a page fault. If  $NFF > NFFMIN$ , the process gets a free frame from the FFLIST, and continues. If  $NFF = NFFMIN$ , then a page replacement is necessary.

We use the approximate LRU replacement algorithm. The 14-th bit in each page-table entry is used as the **reference bit**. Moreover, each page-table entry should additionally contain a 16-bit counter for storing the history of accesses made in the recent past. The 10 essential pages always have 0xffff as this history, and will never be replaced. A page  $q$  with valid bit 1 and with the minimum history will be selected from the page table for replacement. When a new page is loaded (after a page fault or otherwise), its history is set to 0xffff (it is the most recently used page at that moment, and must not be replaced in the next iteration of the binary-search loop).

Any access of  $A[M]$  (with or without page fault) will set the reference bit of the corresponding page. After the end of the binary-search loop, the history of each valid page is updated by a 1-bit right shift followed by inserting the reference bit at the MSB (15-th) position. The reference bits of the valid pages are then cleared.

Once a page  $q$  (a valid page of minimum history) is located for replacement (to accommodate page  $p$  that caused the page fault), the following steps should be carried out in the given order. Recall again that page replacement is done only when  $NFF = NFFMIN$ .

1. First check in the FFLIST whether there exists a free frame  $f$  with last owner  $i$  and page number  $p$ . If that is the case, page  $p$  of process  $i$  has been replaced in the recent past. Use this frame  $f$  to store the new page  $p$ . No reloading of the page from the hard disk is needed in this case. In your program, you actually do not load pages, but follow this strategy to simulate a real-life situation correctly.
2. If Attempt 1 fails, check in the FFLIST whether there is a free frame with no owner. If such a frame  $f$  is found, allocate  $f$  to page  $p$ .
3. If both Attempts 1 and 2 fail, try to locate in the FFLIST a free frame whose last owner was  $i$  (of course, the page number stored will be different from  $p$ , otherwise Attempt 1 would have succeeded). If such a frame  $f$  is found, allocate  $f$  to page  $p$ . If there are multiple free frames with  $i$  as the last owner, pick any one of them.
4. If all the above three attempts fail, pick a random free frame  $f$  from FFLIST for allocating to page  $p$ .

Once a free frame  $f$  is obtained to accommodate page  $p$  of process  $i$ , you need to free the frame  $g$  storing the victim page  $q$  of process  $i$ . Add  $(g, i, q)$  to FFLIST. Also update the page table of process  $i$  to reflect these changes.

When a process exits, add, to the FFLIST, all the frames currently held by that process, after clearing the ownership information in all these frames.

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Use a compile-time flag `VERBOSE` for the type of output. In the default (non-verbose) mode, print for each process (and also the aggregate) the number of page accesses, the number and percentage of page faults, the number and percentage of page replacements, and the number (and percentage) of page replacements handled by each of Attempts 1 – 4. In the verbose mode, supply more detailed paging information. See the sample provided.

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**Submit a single C/C++ file `LRU.c(pp)`.**

## Sample Output

We use the same *search.txt* as in LA9, with  $n = 128$  process, each making  $m = 64$  binary searches. The non-verbose output (statistics only) is given below. An archive (storing other files) will be provided to you. The percentages of faults and replacements are with respect to the number of memory accesses, whereas the percentages in attempt statistics are with respect to the number of page replacements.

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```
gcc -Wall -o runsearch LRU.c
./runsearch
+++ Page access summary
```

PID	Accesses	Faults	Replacements	Attempts
0	1339	409 (30.55%)	335 (25.02%)	6 + 5 + 324 + 0 (1.79% + 1.49% + 96.72% + 0.00%)
1	1326	356 (26.85%)	266 (20.06%)	1 + 4 + 261 + 0 (0.38% + 1.50% + 98.12% + 0.00%)
2	1322	379 (28.67%)	291 (22.01%)	9 + 7 + 275 + 0 (3.09% + 2.41% + 94.50% + 0.00%)
3	1282	336 (26.21%)	266 (20.75%)	8 + 9 + 249 + 0 (3.01% + 3.38% + 93.61% + 0.00%)
4	1280	367 (28.67%)	287 (22.42%)	13 + 11 + 263 + 0 (4.53% + 3.83% + 91.64% + 0.00%)
5	1290	331 (25.66%)	252 (19.53%)	4 + 4 + 244 + 0 (1.59% + 1.59% + 96.83% + 0.00%)
6	1288	373 (28.96%)	293 (22.75%)	4 + 14 + 275 + 0 (1.37% + 4.78% + 93.86% + 0.00%)
7	1341	408 (30.43%)	314 (23.42%)	11 + 15 + 288 + 0 (3.50% + 4.78% + 91.72% + 0.00%)
8	1328	387 (29.14%)	298 (22.44%)	3 + 13 + 282 + 0 (1.01% + 4.36% + 94.63% + 0.00%)
9	1302	381 (29.26%)	293 (22.50%)	0 + 9 + 284 + 0 (0.00% + 3.07% + 96.93% + 0.00%)
10	1334	403 (30.21%)	316 (23.69%)	12 + 8 + 296 + 0 (3.80% + 2.53% + 93.67% + 0.00%)
11	1339	402 (30.02%)	319 (23.82%)	11 + 12 + 296 + 0 (3.45% + 3.76% + 92.79% + 0.00%)
12	1342	417 (31.07%)	322 (23.99%)	15 + 15 + 292 + 0 (4.66% + 4.66% + 90.68% + 0.00%)
13	1334	390 (29.24%)	304 (22.79%)	6 + 13 + 285 + 0 (1.97% + 4.28% + 93.75% + 0.00%)
14	1310	371 (28.32%)	291 (22.21%)	2 + 5 + 284 + 0 (0.69% + 1.72% + 97.59% + 0.00%)
15	1320	381 (28.86%)	300 (22.73%)	10 + 16 + 274 + 0 (3.33% + 5.33% + 91.33% + 0.00%)
16	1337	385 (28.80%)	292 (21.84%)	4 + 12 + 276 + 0 (1.37% + 4.11% + 94.52% + 0.00%)
17	1334	393 (29.46%)	313 (23.46%)	6 + 15 + 292 + 0 (1.92% + 4.79% + 93.29% + 0.00%)
18	1295	362 (27.95%)	284 (21.93%)	8 + 12 + 264 + 0 (2.82% + 4.23% + 92.96% + 0.00%)
19	1334	384 (28.79%)	298 (22.34%)	12 + 12 + 274 + 0 (4.03% + 4.03% + 91.95% + 0.00%)
20	1335	409 (30.64%)	318 (23.82%)	7 + 12 + 299 + 0 (2.20% + 3.77% + 94.03% + 0.00%)
21	1304	360 (27.61%)	270 (20.71%)	10 + 7 + 253 + 0 (3.70% + 2.59% + 93.70% + 0.00%)
22	1291	367 (28.43%)	285 (22.08%)	8 + 11 + 266 + 0 (2.81% + 3.86% + 93.33% + 0.00%)
23	1326	409 (30.84%)	325 (24.51%)	8 + 12 + 305 + 0 (2.46% + 3.69% + 93.85% + 0.00%)
24	1318	395 (29.97%)	311 (23.60%)	1 + 16 + 294 + 0 (0.32% + 5.14% + 94.53% + 0.00%)
25	1338	411 (30.72%)	323 (24.14%)	6 + 12 + 305 + 0 (1.86% + 3.72% + 94.43% + 0.00%)
26	1279	339 (26.51%)	262 (20.48%)	4 + 7 + 251 + 0 (1.53% + 2.67% + 95.80% + 0.00%)
27	1331	397 (29.83%)	306 (22.99%)	5 + 7 + 294 + 0 (1.63% + 2.29% + 96.08% + 0.00%)
28	1312	371 (28.28%)	287 (21.88%)	15 + 13 + 259 + 0 (5.23% + 4.53% + 90.24% + 0.00%)
29	1296	355 (27.39%)	273 (21.06%)	10 + 8 + 255 + 0 (3.66% + 2.93% + 93.41% + 0.00%)
30	1318	384 (29.14%)	304 (23.07%)	11 + 11 + 282 + 0 (3.62% + 3.62% + 92.76% + 0.00%)
31	1337	394 (29.47%)	305 (22.81%)	3 + 10 + 292 + 0 (0.98% + 3.28% + 95.74% + 0.00%)
32	1301	361 (27.75%)	275 (21.14%)	9 + 12 + 254 + 0 (3.27% + 4.36% + 92.36% + 0.00%)
33	1299	350 (26.94%)	266 (20.48%)	11 + 8 + 247 + 0 (4.14% + 3.01% + 92.86% + 0.00%)
34	1282	349 (27.22%)	273 (21.29%)	8 + 10 + 255 + 0 (2.93% + 3.66% + 93.41% + 0.00%)
35	1325	365 (27.55%)	280 (21.13%)	1 + 11 + 268 + 0 (0.36% + 3.93% + 95.71% + 0.00%)
36	1339	422 (31.52%)	341 (25.47%)	10 + 13 + 318 + 0 (2.93% + 3.81% + 93.26% + 0.00%)
37	1337	420 (31.41%)	330 (24.68%)	5 + 11 + 314 + 0 (1.52% + 3.33% + 95.15% + 0.00%)
38	1305	367 (28.12%)	287 (21.99%)	11 + 9 + 267 + 0 (3.83% + 3.14% + 93.03% + 0.00%)
39	1338	376 (28.10%)	290 (21.67%)	5 + 16 + 269 + 0 (1.72% + 5.52% + 92.76% + 0.00%)
40	1306	391 (29.94%)	308 (23.58%)	8 + 12 + 288 + 0 (2.60% + 3.90% + 93.51% + 0.00%)
41	1341	390 (29.08%)	307 (22.89%)	5 + 11 + 291 + 0 (1.63% + 3.58% + 94.79% + 0.00%)
42	1321	388 (29.37%)	303 (22.94%)	14 + 13 + 276 + 0 (4.62% + 4.29% + 91.09% + 0.00%)
43	1342	390 (29.06%)	304 (22.65%)	7 + 7 + 290 + 0 (2.30% + 2.30% + 95.39% + 0.00%)
44	1294	381 (29.44%)	304 (23.49%)	6 + 13 + 285 + 0 (1.97% + 4.28% + 93.75% + 0.00%)
45	1335	394 (29.51%)	310 (23.22%)	0 + 8 + 302 + 0 (0.00% + 2.58% + 97.42% + 0.00%)
46	1343	422 (31.42%)	336 (25.02%)	7 + 11 + 318 + 0 (2.08% + 3.27% + 94.64% + 0.00%)
47	1288	366 (28.42%)	280 (21.74%)	9 + 12 + 259 + 0 (3.21% + 4.29% + 92.50% + 0.00%)
48	1316	355 (26.98%)	272 (20.67%)	8 + 12 + 252 + 0 (2.94% + 4.41% + 92.65% + 0.00%)
49	1334	398 (29.84%)	311 (23.31%)	9 + 14 + 288 + 0 (2.89% + 4.50% + 92.60% + 0.00%)
50	1311	379 (28.91%)	289 (22.04%)	9 + 10 + 270 + 0 (3.11% + 3.46% + 93.43% + 0.00%)
51	1314	368 (28.01%)	281 (21.39%)	8 + 13 + 260 + 0 (2.85% + 4.63% + 92.53% + 0.00%)
52	1335	417 (31.24%)	342 (25.62%)	7 + 12 + 323 + 0 (2.05% + 3.51% + 94.44% + 0.00%)
53	1318	358 (27.16%)	276 (20.94%)	2 + 6 + 268 + 0 (0.72% + 2.17% + 97.10% + 0.00%)
54	1326	391 (29.49%)	311 (23.45%)	5 + 12 + 294 + 0 (1.61% + 3.86% + 94.53% + 0.00%)
55	1336	373 (27.92%)	286 (21.41%)	8 + 10 + 268 + 0 (2.80% + 3.50% + 93.71% + 0.00%)
56	1317	378 (28.70%)	296 (22.48%)	6 + 10 + 280 + 0 (2.03% + 3.38% + 94.59% + 0.00%)
57	1334	402 (30.13%)	308 (23.09%)	0 + 5 + 303 + 0 (0.00% + 1.62% + 98.38% + 0.00%)
58	1335	396 (29.66%)	309 (23.15%)	6 + 7 + 296 + 0 (1.94% + 2.27% + 95.79% + 0.00%)
59	1321	374 (28.31%)	287 (21.73%)	4 + 7 + 276 + 0 (1.39% + 2.44% + 96.17% + 0.00%)
60	1338	397 (29.67%)	316 (23.62%)	1 + 3 + 312 + 0 (0.32% + 0.95% + 98.73% + 0.00%)
61	1330	391 (29.40%)	308 (23.16%)	7 + 6 + 295 + 0 (2.27% + 1.95% + 95.78% + 0.00%)
62	1295	374 (28.88%)	301 (23.24%)	0 + 5 + 296 + 0 (0.00% + 1.66% + 98.34% + 0.00%)
63	1312	375 (28.58%)	290 (22.10%)	1 + 3 + 286 + 0 (0.34% + 1.03% + 98.62% + 0.00%)
64	1339	419 (31.29%)	330 (24.65%)	3 + 7 + 320 + 0 (0.91% + 2.12% + 96.97% + 0.00%)
65	1314	381 (29.00%)	295 (22.45%)	9 + 6 + 280 + 0 (3.05% + 2.03% + 94.92% + 0.00%)
66	1325	404 (30.49%)	314 (23.70%)	2 + 6 + 306 + 0 (0.64% + 1.91% + 97.45% + 0.00%)
67	1307	361 (27.62%)	282 (21.58%)	0 + 0 + 281 + 1 (0.00% + 0.00% + 99.65% + 0.35%)
68	1311	352 (26.85%)	272 (20.75%)	6 + 6 + 260 + 0 (2.21% + 2.21% + 95.59% + 0.00%)
69	1333	414 (31.06%)	323 (24.23%)	6 + 8 + 309 + 0 (1.86% + 2.48% + 95.67% + 0.00%)
70	1310	389 (29.69%)	305 (23.28%)	0 + 4 + 301 + 0 (0.00% + 1.31% + 98.69% + 0.00%)
71	1285	344 (26.77%)	274 (21.32%)	5 + 7 + 262 + 0 (1.82% + 2.55% + 95.62% + 0.00%)
72	1333	403 (30.23%)	315 (23.63%)	3 + 5 + 307 + 0 (0.95% + 1.59% + 97.46% + 0.00%)
73	1280	315 (24.61%)	233 (18.20%)	1 + 0 + 231 + 1 (0.43% + 0.00% + 99.14% + 0.43%)

74	1315	367	(27.91%)	282	(21.44%)	3 +	6 + 273 +	0	(1.06% + 2.13% + 96.81% + 0.00%)
75	1321	368	(27.86%)	278	(21.04%)	0 +	0 + 277 +	1	(0.00% + 0.00% + 99.64% + 0.36%)
76	1305	361	(27.66%)	294	(22.53%)	2 +	7 + 285 +	0	(0.68% + 2.38% + 96.94% + 0.00%)
77	1326	409	(30.84%)	320	(24.13%)	2 +	5 + 313 +	0	(0.62% + 1.56% + 97.81% + 0.00%)
78	1286	345	(26.83%)	265	(20.61%)	4 +	7 + 254 +	0	(1.51% + 2.64% + 95.85% + 0.00%)
79	1331	390	(29.30%)	306	(22.99%)	1 +	7 + 298 +	0	(0.33% + 2.29% + 97.39% + 0.00%)
80	1313	372	(28.33%)	284	(21.63%)	4 +	6 + 274 +	0	(1.41% + 2.11% + 96.48% + 0.00%)
81	1331	377	(28.32%)	286	(21.49%)	11 +	7 + 268 +	0	(3.85% + 2.45% + 93.71% + 0.00%)
82	1315	379	(28.82%)	293	(22.28%)	6 +	6 + 281 +	0	(2.05% + 2.05% + 95.90% + 0.00%)
83	1310	375	(28.63%)	290	(22.14%)	3 +	3 + 284 +	0	(1.03% + 1.03% + 97.93% + 0.00%)
84	1315	369	(28.06%)	286	(21.75%)	0 +	3 + 283 +	0	(0.00% + 1.05% + 98.95% + 0.00%)
85	1319	353	(26.76%)	265	(20.09%)	2 +	9 + 254 +	0	(0.75% + 3.40% + 95.85% + 0.00%)
86	1340	407	(30.37%)	316	(23.58%)	2 +	5 + 309 +	0	(0.63% + 1.58% + 97.78% + 0.00%)
87	1297	351	(27.06%)	263	(20.28%)	2 +	3 + 258 +	0	(0.76% + 1.14% + 98.10% + 0.00%)
88	1289	373	(28.94%)	291	(22.58%)	0 +	0 + 290 +	1	(0.00% + 0.00% + 99.66% + 0.34%)
89	1309	384	(29.34%)	305	(23.30%)	4 +	7 + 294 +	0	(1.31% + 2.30% + 96.39% + 0.00%)
90	1308	361	(27.60%)	281	(21.48%)	0 +	5 + 276 +	0	(0.00% + 1.78% + 98.22% + 0.00%)
91	1320	385	(29.17%)	302	(22.88%)	5 +	5 + 292 +	0	(1.66% + 1.66% + 96.69% + 0.00%)
92	1280	339	(26.48%)	259	(20.23%)	0 +	5 + 254 +	0	(0.00% + 1.93% + 98.07% + 0.00%)
93	1299	369	(28.41%)	298	(22.94%)	3 +	7 + 288 +	0	(1.01% + 2.35% + 96.64% + 0.00%)
94	1315	382	(29.05%)	299	(22.74%)	1 +	4 + 294 +	0	(0.33% + 1.34% + 98.33% + 0.00%)
95	1325	396	(29.89%)	313	(23.62%)	3 +	7 + 303 +	0	(0.96% + 2.24% + 96.81% + 0.00%)
96	1334	409	(30.66%)	316	(23.69%)	1 +	3 + 312 +	0	(0.32% + 0.95% + 98.73% + 0.00%)
97	1312	376	(28.66%)	294	(22.41%)	4 +	7 + 283 +	0	(1.36% + 2.38% + 96.26% + 0.00%)
98	1341	392	(29.23%)	304	(22.67%)	1 +	7 + 296 +	0	(0.33% + 2.30% + 97.37% + 0.00%)
99	1337	402	(30.07%)	316	(23.64%)	2 +	6 + 308 +	0	(0.63% + 1.90% + 97.47% + 0.00%)
100	1328	385	(28.99%)	317	(23.87%)	7 +	8 + 302 +	0	(2.21% + 2.52% + 95.27% + 0.00%)
101	1336	412	(30.84%)	330	(24.70%)	5 +	4 + 321 +	0	(1.52% + 1.21% + 97.27% + 0.00%)
102	1305	345	(26.44%)	269	(20.61%)	7 +	9 + 253 +	0	(2.60% + 3.35% + 94.05% + 0.00%)
103	1280	353	(27.58%)	275	(21.48%)	2 +	7 + 266 +	0	(0.73% + 2.55% + 96.73% + 0.00%)
104	1337	388	(29.02%)	298	(22.29%)	1 +	6 + 291 +	0	(0.34% + 2.01% + 97.65% + 0.00%)
105	1339	397	(29.65%)	310	(23.15%)	3 +	7 + 300 +	0	(0.97% + 2.26% + 96.77% + 0.00%)
106	1289	370	(28.70%)	295	(22.89%)	3 +	5 + 287 +	0	(1.02% + 1.69% + 97.29% + 0.00%)
107	1285	374	(29.11%)	300	(23.35%)	5 +	7 + 288 +	0	(1.67% + 2.33% + 96.00% + 0.00%)
108	1339	405	(30.25%)	319	(23.82%)	6 +	7 + 306 +	0	(1.88% + 2.19% + 95.92% + 0.00%)
109	1314	402	(30.59%)	319	(24.28%)	2 +	5 + 312 +	0	(0.63% + 1.57% + 97.81% + 0.00%)
110	1338	406	(30.34%)	316	(23.62%)	1 +	3 + 312 +	0	(0.32% + 0.95% + 98.73% + 0.00%)
111	1282	364	(28.39%)	283	(22.07%)	2 +	4 + 277 +	0	(0.71% + 1.41% + 97.88% + 0.00%)
112	1339	408	(30.47%)	310	(23.15%)	1 +	6 + 303 +	0	(0.32% + 1.94% + 97.74% + 0.00%)
113	1335	401	(30.04%)	312	(23.37%)	3 +	4 + 305 +	0	(0.96% + 1.28% + 97.76% + 0.00%)
114	1337	406	(30.37%)	319	(23.86%)	3 +	6 + 310 +	0	(0.94% + 1.88% + 97.18% + 0.00%)
115	1303	377	(28.93%)	303	(23.25%)	10 +	9 + 284 +	0	(3.30% + 2.97% + 93.73% + 0.00%)
116	1311	372	(28.38%)	283	(21.59%)	3 +	5 + 275 +	0	(1.06% + 1.77% + 97.17% + 0.00%)
117	1322	372	(28.14%)	281	(21.26%)	1 +	5 + 275 +	0	(0.36% + 1.78% + 97.86% + 0.00%)
118	1317	347	(26.35%)	260	(19.74%)	4 +	6 + 250 +	0	(1.54% + 2.31% + 96.15% + 0.00%)
119	1316	372	(28.27%)	288	(21.88%)	2 +	8 + 278 +	0	(0.69% + 2.78% + 96.53% + 0.00%)
120	1314	371	(28.23%)	284	(21.61%)	2 +	3 + 279 +	0	(0.70% + 1.06% + 98.24% + 0.00%)
121	1300	370	(28.46%)	287	(22.08%)	6 +	5 + 276 +	0	(2.09% + 1.74% + 96.17% + 0.00%)
122	1285	352	(27.39%)	274	(21.32%)	7 +	5 + 262 +	0	(2.55% + 1.82% + 95.62% + 0.00%)
123	1296	368	(28.40%)	281	(21.68%)	8 +	7 + 266 +	0	(2.85% + 2.49% + 94.66% + 0.00%)
124	1321	357	(27.02%)	276	(20.89%)	8 +	7 + 261 +	0	(2.90% + 2.54% + 94.57% + 0.00%)
125	1326	404	(30.47%)	313	(23.60%)	2 +	6 + 305 +	0	(0.64% + 1.92% + 97.44% + 0.00%)
126	1318	389	(29.51%)	310	(23.52%)	1 +	4 + 305 +	0	(0.32% + 1.29% + 98.39% + 0.00%)
127	1316	387	(29.41%)	302	(22.95%)	6 +	7 + 289 +	0	(1.99% + 2.32% + 95.70% + 0.00%)
Total	168588	48691	(28.88%)	37943	(22.51%)	642 +	1000 + 36297 +	4	(1.69% + 2.64% + 95.66% + 0.01%)

You may use the following makefile.

```

run: LRU.c
    gcc -Wall -o runsearch LRU.c
    ./runsearch

vrunk: LRU.c
    gcc -Wall -DVERBOSE -o runsearch LRU.c
    ./runsearch

db: gensearch.c
    gcc -Wall -o gensearch gensearch.c
    ./gensearch

clean:
    -rm -f runsearch gensearch

```