# CS39002 Operating Systems Laboratory Spring 2025

# Lab Assignment: 3 Date of submission: 22–Jan –2025

# Performance Comparison of Some Scheduling Algorithms

In this assignment, you compare the performance of FCFS (First-Come-First-Serve) scheduling with that of RR (Round-Robin) scheduling, and also the effect of the time quantum q on the performance of RR scheduling. You deal with the case of a single CPU. FCFS scheduling can be considered as a special case of RR scheduling with  $q = \infty$ , so you only work on RR scheduling, and specify q as a parameter. You do not deal with real processes, but simulate the behavior of the CPU on processes with random behavior.

The details of the processes are stored in a text file *proc.txt*. The file starts with the number n of processes that you are going to deal with. Each of the remaining n lines in the file stores the following information about a process. The lines may appear in any sequence (not necessarily sorted by IDs or arrival times).

# ID ARRIVAL\_TIME $CPU_1 IO_1 CPU_2 IO_2 \ldots CPU_b -1$

Each process is given a unique integer ID (you may think of this as the PID of the process). This is followed by the time when the process arrives (a non-negative integer). This is followed by a sequence of alternate CPU- and IO-burst times (positive integers). The list ends with -1 as an IO-burst time indicating the end of the process. Assume that all times are in milliseconds. In reality, the burst times are not known beforehand, but you deal with the performance of the scheduling algorithms on the given processes, so assume that these burst times are exact. You are not required to implement SJF scheduling, so whether or not you can estimate burst times is not a concern to this assignment.

A process-generator program is supplied to you as *genproc.c.* Compile and run the code with n (number of processes) as an optional command-line argument. Without this argument, n = 100 is taken. This generates data for n processes with the following randomly chosen parameters.

Process type	Probability	CPU burst time	IO burst time	# CPU bursts	# IO burst
IO-bound	90%	1 – 15	50 - 200	4 - 10	3 – 9
CPU-bound	10%	100 - 300	50 - 200	3 – 7	2-6

Each process is assumed to start and end with CPU bursts, so the number of IO bursts is one less than the number of CPU bursts. You may use these statistics to determine your array sizes.

## Part 1: Read the input file

Define a structure to store the information about the processes. Each process requires storage for the ID, the arrival time, the number of bursts, and the burst times. In addition, you store additional information (like state of the process, which burst it is in, and so on) as needed by the simulation. It is your choice whatever you need to store in each process record. Read *proc.txt*, and populate an array of process-info structures. The queues to be used in the simulation will consist only of indices in this array (and nothing else). In essence, each entry in the process-info array *behaves* like the PCB of a process.

### Part 2: Data structure for the Ready Queue

For both FCFS and RR scheduling, the ready queue is implemented as a FIFO queue. Design a data structure to perform the init, front, enqueue, dequeue (and other operations) on a FIFO queue. Each element in the queue will be an index identifying the process in the process-info table. So it will be a queue of integers. Make your own queue implementation.

## Part 3: Handling the events [Discrete-event system simulation]

Your simulation neither creates any of the *n* processes nor runs the process on any CPU. On the contrary, your program assumes that the processes behave as stored in the process-info array, and schedules the processes on an imaginary CPU. In order to do so, you need to handle events *chronologically* starting at time 0. An *event* in the simulation is an incident that occurs at a particular time and that changes the configuration of the simulated system. In particular, you need to handle the following events.

- Arrival of a new process
- End of a process
- End of the use of CPU (after timeout or CPU-burst end)
- Rearrival of a process after IO completion

These events must be handled in a non-decreasing order of the times when they happen. A min-priority queue is used for that purpose, and is called the *event queue*. Like the ready queue, the elements of the queue are integers which are indices in the process-info table. At the beginning of the simulation, only the *n* arrival times of the processes are known. The event queue should be initialized by these *n* arrivals. Later, when the simulation proceeds, the exact times of the other types of events will be calculated, and inserted in the event queue. The first ordering parameter for the even queue is the times of occurrences of the events. For multiple events happening at the same time (like two processes complete IO and a new process arrives at the same time), some tie-breaking policies are to be used. In respect of the joining of a process to the ready queue, the following ordering is to be maintained:

Arrival (for the first time or after IO completion) < CPU timeout

Ties are possible even with this ordering. Give precedence to smaller IDs to break the ties. For example, suppose that a process with ID j arrives at the same time when a process with ID i completes an IO burst. If we have i < j, then the IO-completion event will appear earlier in the event queue than the new-arrival event.

In order to handle the next event, extract the first element from the event queue. Set the current time to the time of the event. Depending on the type of the event, do the following.

Arrival of a process for the first time or after IO completion: Put the process at the back of the ready queue.

**CPU burst ends:** Check whether this was the last CPU burst of the process. If so, print that the process exits, and also the following performance measures pertaining to that process.

- Turnaround time
- Turnaround time as a percentage of running time (total CPU + IO burst times)
- Wait time (turnaround time running time)

Otherwise (that is, the process has more bursts to do), the next burst of the process will be an IO burst. Add the time of that IO burst to the current time. Insert the (re)arrival event of that process after IO completion, to the event queue.

**CPU timeout:** The current CPU burst is preempted, that is, more time is needed to complete the current CPU burst, so insert the process at the back of the ready queue.

After handling each event, check whether the CPU is free. If so, and if the ready queue is not empty, schedule the process at the front of the ready queue for the next time quantum q or for the rest of the next CPU burst time (whichever is smaller).

Eventually, all the n processes completes all the bursts, and the event queue becomes empty. Stop the simulation at that point.

Make your own array-based implementation of the event queue as a (binary) min-heap.

### **Part 4: Performance figures**

Print the per-process performance figures as explained in Part 3. In addition, print the following aggregate information after the simulation stops.

- Average wait time of a process
- Total turnaround time (simulation end time simulation start time (0))
- Total idle time of the CPU in the simulation interval (total turnaround time)
- Percentage utilization of the CPU during the simulation interval

#### Part 5: main() function

In your main() function, call the scheduler for  $q = \infty$  (a large integer like 10<sup>9</sup>) to simulate FCFS scheduling. Then make two other calls of the scheduler to simulate RR scheduling with q = 10 and q = 5.

#### Part 6: Verbose/Concise output

The default behavior of your code would be to print only the per-process and aggregate performance figures. Enable verbose output by a compile-time flag VERBOSE. In the verbose mode, print all individual events (arrival/rearrival, departure, CPU end, scheduling decision). The printing format will be specified at the end.

#### Makefile

You may use the following makefile.

```
compile: schedule.c
    gcc -Wall -o schedule schedule.c
run: compile
    ./schedule
vcompile: schedule.c
    gcc -Wall -o schedule -DVERBOSE schedule.c
vrun: vcompile
    ./schedule
db: genproc.c
    gcc -Wall -o genproc genproc.c
clean:
        -rm -f genproc schedule proc.txt
```

Submit a single file *schedule.c*.

# Sample Input

20																						
	1	Θ	11	115	8	144	10	111	3	109	12	76	5	170	9	-1						
	2	91	14	129	8	84	2	131	8	186	15	67	4	-1								
	3	493	13	138	1	60	1	134	9	64	9	176	5	50	6	144	5	198	14	-1		
	4	796	4	91	12	134	13	51	7	118	3	72	11	80	5	174	3	190	13	-1		
	5	978	170	90	227	72	165	179	207	98	113	141	140	- 1								
	61	1116	14	53	4	65	7	80	8	108	12	196	12	145	2	136	2	185	7	181	1	-1
	71	1347	15	69	5	86	6	160	4	139	11	134	10	-1								
	81	1683	12	185	5	100	10	147	6	162	3	-1										
	91	1774	14	166	8	75	11	75	8	69	15	52	2	128	10	-1						
1	0 2	2172	11	157	13	92	14	155	14	118	9	136	10	130	2	-1						
1	1 2	2249	12	164	3	105	15	119	15	98	14	102	7	57	12	184	15	89	7	-1		
1	2 2	2674	5	72	8	124	15	175	7	135	11	137	3	81	7	161	14	86	13	-1		
1		2884	5	106	2	97	10	186	8	-1												
1		2998	5	130	1	150	3	190	5	51	1	68	15	56	14	186	14	185	12	-1		
1		3222	15	100	11	85	6	92	15	79	2	124	9	70	8	124	2	53	1	92	14	-1
1		3587	3	156	7	150	7	147	11	127	3	77	11	102	14	-1						
1		3724	6	161	11	186	9	53	15	149	3	-1										
1		3761	14	96	10	133	8	75	8	136	1	173	5	64	13	188	8	107	15	-1		
1	93	3958	14	186	4	145	14	92	9	195	9	81	10	154	7	-1						
2	0 4	1335	280	116	221	142	147	70	175	81	134	69	190	111	237	-1						

# Sample Output (without the VERBOSE flag)

	FS Scheduling *	****				
739	: Process					, Wait time = 0
783	: Process	1 exits.	Turnaround	time =	783 (100%)	, Wait time = 0
1752	: Process	3 exits.	Turnaround	time =	1259 (123%)	, Wait time = 232
2363	: Process	4 exits.	Turnaround	time =	1567 (160%)	, Wait time = 586
2588	: Process	5 exits.	Turnaround	time =	1610 (100%)	, Wait time = 8
2598	: Process	7 exits.	Turnaround	time =	1251 (196%)	, Wait time = 612
2760	: Process	9 exits.	Turnaround	time =	986 (156%)	, Wait time = 353
2771	: Process	8 exits	Turnaround	time =	1088 (173%)	, Wait time = 458
2994	: Process					, Wait time = 660
3194	: Process					, Wait time = 161
3304	: Process					, Wait time = 6
3464	: Process					, Wait time = 197
3750	: Process	12 ovito	Turnaround	time -	1076 (102%)	, Wait time = 22
4097	: Process					, Wait time = 13
4134	: Process	15 exits.		time =	912 (101%)	, Wait time = 10
4329	: Process	1/ exits.	lurnaround	time =	605 (102%)	, Wait time = 12
4629	: Process					, Wait time = 227
5091	: Process					, Wait time = 276
5248	: Process					, Wait time = 370
6308	: Process	20 exits.	Turnaround	time =	1973 (100%)	, Wait time = 0
Average	wait time = 21	L0.15				
Total t	urnaround time	= 6308				
CPU idl	e time = 2752					
CPU uti	lization = 56.3	37%				
**** RR	Scheduling wit	th a = 10 ***	**			
739	: Process			time =	648 (100%)	, Wait time = 0
783	: Process					, Wait time = 0
1548	: Process	3 exits	Turnaround	time =	1055 (103%)	, Wait time = 28
1838	: Process					, Wait time = 61
2040			Turnaround	time =	603 (100%)	Watt time = 54
2336	: Process			tune =		, Wait time = 54
	: Process					, Wait time = 23
2426	: Process					, Wait time = 92
2461	: Process	9 exits.	Turnaround	time =	687 (109%)	, Wait time = 54
2918	: Process					, Wait time = 338
3077	: Process					, Wait time = 44
3317	: Process	13 exits.	Turnaround	time =	433 (105%)	, Wait time = 19
3336	: Process	11 exits.	Turnaround	time =	1087 (107%)	, Wait time = 69
3765	: Process	12 exits.	Turnaround	time =	1091 (104%)	, Wait time = 37
4132	: Process	14 exits.	Turnaround	time =	1134 (104%)	, Wait time = 48
4146	: Process	15 exits.	Turnaround	time =	924 (102%)	, Wait time = 22
4358	: Process					, Wait time = 41
4441	: Process					, Wait time = 39
4879	: Process	18 exits.	Turnaround	time =	1118 (106%)	, Wait time = 64
4906	: Process	19 exits.	Turnaround	time =	948 (103%)	, Wait time = 28
6383	: Process					, Wait time = 75
0000		Lo chico.		c crite	2010 (2010)	, note cane is
Average	wait time = 50	5 80				
	urnaround time					
	e time = 2827	- 0305				
		710/				
CPU ULL	lization = 55.7	1%				
++++ 00	Cabadultan	a e asso	+			
	Scheduling wit				CAO (4000)	under et al.
739	: Process					, Wait time = 0
783	: Process					, Wait time = 0
1543	: Process	3 exits.	Iurnaround	time =	1050 (102%)	, Wait time = 23
1828	: Process	4 exits.	Turnaround	time =	1032 (105%)	, Wait time = 51
2037	: Process	7 exits.	Turnaround	time =	690 (108%)	, Wait time = 51
2334	: Process	8 exits.	Turnaround	time =	651 (103%)	, Wait time = 21
2419	: Process	6 exits.	Turnaround	time =	1303 (107%)	. Wait time = 85
2444	: Process	9 exits.	Turnaround	time =	670 (106%)	, Wait time = 37
2916	Process	5 exits	Turnaround	time =	1938 (121%)	Wait time = 336

2419: ProcessG exits. Turnaround time = 1033 (107%), Wait time = 852444: Process9 exits. Turnaround time = 670 (106%), Wait time = 852444: Process9 exits. Turnaround time = 670 (106%), Wait time = 3363065: Process10 exits. Turnaround time = 1938 (121%), Wait time = 3323077: Process10 exits. Turnaround time = 423 (102%), Wait time = 323174: Process11 exits. Turnaround time = 423 (102%), Wait time = 323184: Process11 exits. Turnaround time = 1065 (105%), Wait time = 473780: Process12 exits. Turnaround time = 1066 (105%), Wait time = 524092: Process14 exits. Turnaround time = 1004 (101%), Wait time = 84133: Process15 exits. Turnaround time = 607 (102%), Wait time = 194331: Process16 exits. Turnaround time = 607 (102%), Wait time = 274866: Process18 exits. Turnaround time = 842 (103%), Wait time = 514908: Process19 exits. Turnaround time = 950 (103%), Wait time = 306380: Process20 exits. Turnaround time = 2045 (104%), Wait time = 72

Average wait time = 48.25 Total turnaround time = 6380 CPU idle time = 2824 CPU utilization = 55.74%

## Sample Output (with the VERBOSE flag)

\*\*\*\* RR Scheduling with q = 10 \*\*\*\* cheduling with q = 10 \*\*\*\* : Starting : Process 1 joins ready queue upon arrival : Process 1 is scheduled to run for time 10 : Process 1 joins ready queue after timeout : Process 1 is scheduled to run for time 1 : CPU goes idle : Process 2 joins ready queue upon arrival : Process 2 joins ready queue after timeout : Process 2 is scheduled to run for time 10 : Process 2 is scheduled to run for time 4 : CPU goes idle : Process 1 joins ready queue after IO compl : Process 1 joins ready queue after IO completion : Process 1 is scheduled to run for time 8 : CPU goes idle CPU goes idle
Process 2 joins ready queue after IO completion
Process 2 is scheduled to run for time 8
CPU goes idle
Process 1 joins ready queue after IO completion
Process 1 is scheduled to run for time 10
CPU goes idle
Process 2 joins ready queue after IO completion
Process 2 is scheduled to run for time 2
CPU goes idle
Process 1 ions ready queue after IO completion 328 : CPU goes tole : Process 1 joins ready queue after IO completion : Process 1 is scheduled to run for time 3 : CPU goes idle : Process 2 joins ready queue after IO completion : Process 2 is scheduled to run for time 8 459 CPU goes idle CPU goes idle Process 3 joins ready queue upon arrival Process 3 is scheduled to run for time 10 Process 3 is scheduled to run for time 3 503 : Process 1 joins ready queue after IO completion : Process 1 joins ready queue after IO completion : Process 1 joins ready queue after timeout : Process 1 joins ready queue after timeout : Process 1 is scheduled to run for time 2 521 : CPU goes idle : Process 1 joins ready queue after IO completion : Process 1 is scheduled to run for time 5 644 : CPU goes idle : CPU goes idle : Process 3 joins ready queue after IO completion : Process 3 is scheduled to run for time 1 : CPU goes idle : Process 2 joins ready queue after IO completion : Process 2 is scheduled to run for time 10 : Process 2 is scheduled to run for time 5 : CPU goes idle : Process 3 joins ready queue after IO completion 668 : Process 3 joins ready queue after IO completion Process 3 is scheduled to run for time 1 : CPU goes idle : Process 2 joins ready queue after IO completion : Process 2 is scheduled to run for time 4 : Process 2 exits. Turnaround time = 648 (100%), Wait time = 0 . cru goes tale : Process 1 joins ready queue after IO completion : Process 1 is scheduled to run for time 9 : Process 1 exits. Turnaround time = 783 (100%), Wait time = 0 : CPU goes idle : CPU goes idle 

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: Process 20 is scheduled to run for time 10 : Process 20 joins ready queue after timeout : Process 20 is scheduled to run for time 10 : Process 20 is scheduled to run for time 10 : Process 20 joins ready queue after timeout : Process 20 is scheduled to run for time 10 : Process 20 is scheduled to run for time 10 : Process 20 joins ready queue after timeout 6336 6356 Process 20 joins ready queue after timeout Process 20 is scheduled to run for time 7 Process 20 exits. Turnaround time = 2048 (104%), Wait time = 75 : CPU goes idle Average wait time = 56.80

Total turnaround time = 6383 CPU idle time = 2827 CPU utilization = 55.71%

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