CS21201 Discrete Structures, Autumn 2021–2022

Second Test

Date: Oct 05, 2021	Time: 10:15am–11:30am	Maximum marks: 40

Prove that the following C function terminates for all non-negative integer inputs a, b, c. Here, the divisions by 2 are to be considered as divisions of int variables. (10)

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void wow ( int a, int b, int c )
{
    int r, s, t;
    while (1) {
        if ((a == b) || (b == c) || (c == a)) break;
        r = (a + b) / 2; s = (b + c) / 2; t = (c + a) / 2;
        a = r; b = s; c = t;
    }
}
```

Solution The loop in the function maintains the following two invariance properties.

- (1) a, b, c always remain non-negative.
- (2) $\max(a, b, c)$ decreases strictly from one iteration to the next.

The first invariance is obvious. In order to prove the second invariance, consider the situation a > b > c (the other situations can be handled analogously). In this case, $\max(a, b, c) = a$. After the loop body is executed once, the maximum becomes $\lfloor (a+b)/2 \rfloor \leq \lfloor (a+(a-1))/2 \rfloor = \lfloor a-\frac{1}{2} \rfloor = a-1 < a$.

Since $\max(a, b, c)$ is a non-negative integer by the first invariance, the second invariance implies that the loop cannot repeat forever.

- 2. 65 distinct integers are chosen in the range 1, 2, 3, ..., 2021. Prove that there must exist four of the chosen integers (call them a, b, c, d) such that a b + c d is a multiple of 2021. (10)
- Solution The total count of 2-subsets of the 65 chosen integers is $\binom{65}{2} = 2080 > 2021$. So we can find two distinct subsets $S = \{a, c\}$ and $T = \{b, d\}$ of the chosen integers such that $a + c \equiv b + d \pmod{2021}$, that is, $a b + c d \equiv 0 \pmod{2021}$. We need to show that $S \cap T = \emptyset$. Suppose not. Since S and T are distinct, we must have $|S \cap T| = 1$. Say, $a = b \pmod{c \neq d}$. The condition $a + c \equiv b + d \pmod{2021}$ implies that $c \equiv d \pmod{2021}$. But c and d are chosen in the range [1, 2021], so they must be equal, a contradiction.
- **3.** Let ρ and σ be two binary relations over the set \mathscr{A} . A *composite relation* $\rho \circ \sigma$ over \mathscr{A} is defined as

 $\rho \circ \sigma = \{(p,r) \mid \text{ there exists some } q \in \mathscr{A} \text{ such that } (p,q) \in \rho \text{ and } (q,r) \in \sigma \}.$

Prove the following assertions with precise formal justifications.

(a) If ρ and σ are equivalence relations, then $\rho \circ \sigma$ is an equivalence relation if and only if $\rho \circ \sigma = \sigma \circ \rho$. (6)

Solution $[\Rightarrow]$ Suppose that $(x,y) \in \rho \circ \sigma$ $(x,y \in \mathscr{A})$. Since $\rho \circ \sigma$ is an equivalence relation, we also have $(y,x) \in \rho \circ \sigma$ (symmetric property). This means that for some $\alpha \in \mathscr{A}$, we have $(y,\alpha) \in \rho$ and $(\alpha,x) \in \sigma$. Since ρ and σ are both equivalence relations, we further get $(\alpha, y) \in \rho$ and $(x, \alpha) \in \sigma$ (symmetric property) This means that $(x,y) \in \sigma \circ \rho$ (by definition). Therefore $\rho \circ \sigma \subseteq \sigma \circ \rho$. Similar arguments (in opposite direction) can be given to prove $\sigma \circ \rho \subseteq \rho \circ \sigma$, thereby establishing $\rho \circ \sigma = \sigma \circ \rho$.

[\Leftarrow] Since ρ and σ are both equivalence relations, $(x,x) \in \rho$ as well as $(x,x) \in \sigma$ (for all $x \in \mathscr{A}$). By the definition of composite relations, we immediately have $(x,x) \in \rho \circ \sigma$, proving that $\rho \circ \sigma$ is *reflexive*.

If $(x, y) \in \rho \circ \sigma$ $(x, y \in \mathscr{A})$, then for some $\alpha \in \mathscr{A}$, we have $(x, \alpha) \in \rho$ and $(\alpha, y) \in \sigma$. Since ρ and σ are both equivalence relations, we have $(\alpha, x) \in \rho$ and $(y, \alpha) \in \sigma$ (symmetric property). This means that $(y, x) \in \sigma \circ \rho$ (by definition). Finally, $\rho \circ \sigma = \sigma \circ \rho$ implies $(y, x) \in \rho \circ \sigma$. This proves that $\rho \circ \sigma$ is *symmetric*.

Let $x, y, z \in \mathscr{A}$. Suppose that $(x, y) \in \rho \circ \sigma$ and $(y, z) \in \rho \circ \sigma$. Since $(x, y) \in \rho \circ \sigma$, there exists $\alpha \in \mathscr{A}$ such that $(x, \alpha) \in \rho$ and $(\alpha, y) \in \sigma$. Since $(y, z) \in \rho \circ \sigma$, there exists $\beta \in \mathscr{A}$ such that $(y, \beta) \in \rho$ and $(\beta, z) \in \sigma$. But

then, since $(\alpha, y) \in \sigma$ and $(y, \beta) \in \rho$, we have $(\alpha, \beta) \in \sigma \circ \rho$ (by definition). It is given that $\sigma \circ \rho = \rho \circ \sigma$, so $(\alpha, \beta) \in \rho \circ \sigma$, that is, there exist $\delta \in \mathscr{A}$, such that $(\alpha, \delta) \in \rho$, and $(\delta, \beta) \in \sigma$. Since ρ is transitive, and (x, α) and (α, δ) are in ρ , we have $(x, \delta) \in \rho$. Moreover, since σ is transitive, and (β, β) and (β, z) are in σ , we have $(\delta, z) \in \sigma$. By definition, we then have $(x, z) \in \rho \circ \sigma$, that is, $\rho \circ \sigma$ is *transitive*.

(b) The *inverse* of a relation τ over \mathscr{A} is defined as $\tau^{-1} = \{(q, p) \mid (p, q) \in \tau\}$ $(p, q \in \mathscr{A})$. Prove that $(\rho \circ \sigma)^{-1} = (\sigma^{-1} \circ \rho^{-1}).$ (4)

Solution Let (y,x) ∈ (ρ ∘ σ)⁻¹ for x, y ∈ A. By definition, (x,y) ∈ (ρ ∘ σ), that is, for some α ∈ A, we have (x, α) ∈ ρ and (α, y) ∈ σ. This also implies that (α,x) ∈ ρ⁻¹ and (y, α) ∈ σ⁻¹. Therefore (y,x) ∈ σ⁻¹ ∘ ρ⁻¹, concluding that (ρ ∘ σ)⁻¹ ⊆ (σ⁻¹ ∘ ρ⁻¹).
On the other hand, let (y,x) ∈ σ⁻¹ ∘ ρ⁻¹ for x, y ∈ A. Then, for some α ∈ A, we have (y,α) ∈ σ⁻¹ and (α,x) ∈ ρ⁻¹ (by definition). This also implies that (α,y) ∈ σ and (x,α) ∈ ρ. Since (x,y) ∈ ρ ∘ σ, we have (y,x) ∈ (ρ ∘ σ)⁻¹, concluding that (σ⁻¹ ∘ ρ⁻¹) ⊆ (ρ ∘ σ)⁻¹.
Together, we have proved that (ρ ∘ σ)⁻¹ = (σ⁻¹ ∘ ρ⁻¹).

4. Let $\mathscr{P}(S)$ denote the power set of *S*. For a function $f: X \to Y$, define two functions $g: \mathscr{P}(A) \to \mathscr{P}(B)$ and $h: \mathscr{P}(B) \to \mathscr{P}(A)$ as

$$g(A) = \{b \mid \exists a \in A, f(a) = b\}, \text{ and}$$

$$h(B) = \{a \mid f(a) \in B\}$$

for all $A \subseteq X$ and $B \subseteq Y$. Prove the following assertions with precise formal justifications.

(a) f is injective if and only if h(g(A)) = A for all $A \subseteq X$.

Solution [If] To show that if h(g(A)) = A for all $A \subseteq X$, then f is injective.

Let $f(x_1) = f(x_2)$ for some $x_1, x_2 \in X$. Then, $x_1 \in h(g(\{x_1\}))$. But $h(g(\{x_1\})) = h(\{f(x_1)\}) = h(\{f(x_2)\}) = h(g(\{x_2\})) = \{x_2\}$ (by taking $A = \{x_2\}$ in the hypothesis). It follows that $x_1 \in \{x_2\}$, that is, $x_1 = x_2$.

[Only if] To show that if *f* is injective, then h(g(A)) = A for all $A \subseteq X$.

 $[A \subseteq h(g(A))] \quad a \in A \Rightarrow f(a) \in g(A) \Rightarrow a \in h(g(A)).$

 $[h(g(A)) \subseteq A]$ $a \in h(g(A)) \Rightarrow f(a) \in g(A) \Rightarrow \exists x \in A, f(x) = f(a) \Rightarrow x = a \text{ (since } f \text{ is injective)} \Rightarrow a \in A.$

(b) f is surjective if and only if g(h(B)) = B for all $B \subseteq Y$.

Solution [If] To show that if g(h(B)) = B for all $B \subseteq Y$, then f is surjective.

Take any $b \in Y$, and $B = \{b\}$. By hypothesis, $g(h(B)) = B = \{b\}$. This implies that there exists $a \in h(B)$ such that f(a) = b. Since $h(B) \subseteq X$, it follows that f is surjective.

[Only if] To show that if *f* is surjective, then g(h(B)) = B for all $B \subseteq Y$.

 $[g(h(B)) \subseteq B]$ Let $b \in g(h(B))$. By the definition of g, there exists $a \in h(B)$ such that f(a) = b. But then by the definition of h, we have $f(a) \in B$, that is, $b \in B$.

 $[B \subseteq g(h(B))]$ Let $b \in B$. Since f is surjective, we have b = f(a) for some $a \in X$. By the definition of h, we then have $a \in h(B)$. By the definition of g, we have $f(a) \in g(h(B))$, that is, $b \in g(h(B))$.

(5)

(5)