## INDIAN INSTITUTE OF TECHNOLOGY, KHARAGPUR Computer Science and Engineering

## Switching Circuits and Logic Design (CS21002)

Class Test – I (Spring)

Name:		Roll number:
Date: Wed, Feb 05, 2020	Marks: 90	Time: 6:15-7:15pm (AN)
Place: CSE-107, CSE-108, CSE-119	). CSE-120. CSE-302	Students: 72+73=145

## Answer ALL the questions. Do rough work on available blank pages/sheets.

Q1: Run the *double dabble* (also called *add-3 and shift*) algorithm to convert the binary number 110100101 to BCD showing each step clearly. The operations should be either L Sft for left shift or Add 3. The entries for B1, B2 and B3 should be their values after application of the indicated operation.

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Operation	B2	B1	В0	421
Start	0000	0000	0000	110100101
L Sft	0000	0000	0001	110100101
L Sft	0000	0000	0011	110100101
L Sft	0000	0000	0110	110100101
Add 3	0000	0000	<u>1001</u>	110100101
L Sft	0000	0001	0011	110100101
T. 00	0000	0010	0110	110100101
L Sft	0000	0010	0110	110100101
A 112	0000	0010	1001	110100101
Add 3	0000	0010	<u>1001</u>	110100101
1 00	0000	0101	0010	110100101
L Sft	0000	0101	0010	110100101
A 11 2	0000	1000	0010	110100101
Add 3	0000	1000	0010	110100101
I CG	0001	0000	0101	110100101
L Sft	0001	0000	0101	110100101
A dd 2	0001	0000	1000	110100101
Add 3	0001	0000	<u>1000</u>	110100101
L Sft	0010	0001	0000	110100101
L SII				110100101
L Sft	0100	0010	0001	110100101
Finish	4	2	1	110100101
1 1111311	1		1	

Q2: Represent 29 and -17 in 8-bit signed 2's complement representation; then perform the following operations on your representations, showing each step clearly:

(a) 29+(-17)

Item	Binary representation
17	0001 0001
-17 (2's complement)	1110 1111
29	0001 1101
+ (-17)	1110_1111
Result12	0000 1100

(b) (-17)+(-29)

Item	Binary representation
29	0001 1101
-29 (2's complement)	1110 0011
-17	1110 1111
+ (-29)	1110_0011
Result46	1101 0010

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Q3: Bits 1111011 corresponding to  $\langle p_1, p_2, d_1, p_3, d_2, d_3, d_4 \rangle$  for single bit ECC Hamming code is received; here  $d_1, d_2, d_3$  and  $d_4$  are the data bits while  $p_1, p_2$  and  $p_3$  are the parity bits. Assume that at most one of the bits that are received may be corrupted (i.e.  $\bar{b}$  received instead of b).

Association of parity bits to the data bits may be done according to the table below.

Bits positions	7	6	5	4	3	2	1
(starting at 1)							
Binary	111	110	101	100	011	010	001
Received	$d_4$	$d_3$	$d_2$	$p_3$	$d_1$	$p_2$	$p_1$
data/parity bit							
Association	$p_3, p_2, p_1$	$p_3, p_2$	$p_3, p_1$	$p_3$	$p_2, p_1$	$p_2$	$p_1$
Recomputed bit	$d_4^{ m r}$	$d_3^{ m r}$	$d_2^{ m r}$	$p_3^{ m r}$	$d_1^{ m r}$	$p_2^{ m r}$	$p_1^{ m r}$
name							

Let  $e_i$  be the Boolean value that is 1 if and only if the parity condition involving bit  $p_i$  is violated. Fill up the following table to present the Boolean expression (using  $\cdot, +, \oplus, \leftrightarrow, '$ ) for  $e_i$  in terms of the received bits.

Parity violation indicator	Boolean exp	Truth value	
$ e_1 $	$e_1 =$	$p_1 \oplus d_1 \oplus d_2 \oplus d_4$	1
$e_2$	$e_2 = _{\_\_}$	$p_2 \oplus d_1 \oplus d_3 \oplus d_4$	0
$e_3$	$e_3 = $	$p_3 \oplus d_2 \oplus d_3 \oplus d_4$	1

Fill up the following table to present the Boolean expression (using  $\cdot, +, \oplus, \leftrightarrow, '$ ) to capture that an error has occured in the indicated bit, as error flags (EF); your expressions should be in terms of  $e_1, e_2, e_3$ .

Bit	Value	EF	Boolean expression to compute error flag	EF value
$p_1$	1	$f_1$	$f_1 = e_1 \cdot e_2' \cdot e_3'$	0
$p_2$	1	$f_2$	$f_2 = $ $e_1' \cdot e_2 \cdot e_3'$	0
$p_3$	1	$f_3$	$f_3 = e_1' \cdot e_2' \cdot e_3$	0
$d_4$	1	$g_4$	$g_4 = e_1 \cdot e_2 \cdot e_3$	
$d_3$	1	$g_3$	$g_3 = e_1' \cdot e_2 \cdot e_3$	_0_
$d_2$	0	$g_2$	$g_2 = e_1 \cdot e_2' \cdot e_3$	_1_
$d_1$	1	$g_1$	$g_1 = e_1 \cdot e_2 \cdot e_3'$	0

Finally, fill up the following table to present the Boolean expression (using  $\cdot$ , +,  $\oplus$ ,  $\leftrightarrow$ ,') in terms of received bits and the computed flag bits that will allow the correct bits to be recovered.

Recomputed bit	Boolean expression to recompute the bit	Recomputed value
$p_1^{ m r}$	$p_1^{ m r}= \qquad \qquad p_1\oplus f_1$	_1_
$p_2^{ m r}$	$p_2^{ ext{r}} = p_2 \oplus f_2$	1
$p_3^{ m r}$	$p_3^{ ext{r}} =  p_3 \oplus f_3$	1
$d_4^{ m r}$	$d_4^{ ext{r}} = \underline{\qquad \qquad d_4 \oplus g_4}$	1
$d_3^{ m r}$	$d_3^{ ext{r}} = \underline{\qquad \qquad \qquad } d_3 \oplus g_3$	1
$d_2^{ m r}$	$d_2^{ extsf{r}} = \qquad \qquad$	
$d_1^{ m r}$	$d_1^{ ext{r}} = \underline{\qquad \qquad d_1 \oplus g_1}$	1