

# Advanced graph theory: Tutorial 3: CS60047

## Autumn 2024

Time: 70 minutes; Maximum marks: 100

August 23, 2024

1. The easier part of Kuratowski's theorem is to show that the presence of homeomorphs of  $K_5$  or  $K_{3,3}$  as subgraphs would make a graph non-planar. Prove this by (i) showing that  $K_5$  and  $K_{3,3}$  are non-planar, and (ii) the presence of a homeomorph of a non-planar graph causes non-planarity.

[A graph  $G$  is a homeomorph of another graph  $H$  if  $G$  can be obtained by repeatedly adding degree-2 vertices  $w$  by deleting edge  $\{u, v\}$ , and adding edges  $\{u, w\}$  and  $\{w, v\}$ . Note that  $H$  is planar if and only if its homeomorph  $G$  is planar.]

[This amounts to showing the necessary condition that homeomorphs of none of the two Kuratowski's graphs can appear as subgraphs in a planar graph.]

[The tougher part of Kuratowski's theorem is to show that a graph is planar if it does not have subgraphs homeomorphic to the any of the two Kuratowski graphs.]

2. Show that every 2-connected 3-regular graph has a 1-factor.
3. Use Hall's theorem to show that  $\beta(G) \leq \alpha'(G)$  in a bipartite graph  $G(X \cup Y, E)$ .

(10 marks)

[Show that for any minimum cardinality vertex cover  $C$  in  $G$ , we can demonstrate a matching  $M$  of the same size. The minimum vertex cover  $C$  can span both the two partites of the bipartite graph  $G$ . Such a partition of the vertex cover  $C$  into two parts, say  $C \cap X$  and  $C \cap Y$ , would create two disjoint sub-bipartite graphs  $G(C \cap X, Y \setminus C)$  and  $G(C \cap Y, X \setminus C)$ , where  $N(S)$  denotes the neighbourhood set of vertices of  $S$  in  $G$ . Apply Hall's theorem to both these sub-bipartite graphs.]

4. Minimum  $s - t$  separating vertex cuts are called  $s - t$  *clots*. The cardinality  $\kappa(s, t)$  of an  $s - t$  *clot* is called the  $s - t$  *clot number*.

Show that the vertex connectivity  $\kappa(G)$  of a graph  $G$  is the minimum of all non-adjacent pair clot numbers in  $G$ .

A *thread* between  $s$  and  $t$  of a graph  $G$  is a set of paths between  $s$  and  $t$  which have pairwise no vertices common except for  $s$  and  $t$ . The number of paths in a thread is called the *ply number* of the thread. The maximum ply number of a thread between  $s$  and  $t$  is called the  $s - t$  *ply number* and written as  $p(s, t)$ . The *ply number*  $p(G)$  of  $G$  is the minimum of  $p(s, t)$  over all non-adjacent pairs  $s$  and  $t$  in  $G$ .

Explain the vertex version of the theorem of Menger in terms of the *ply* and *clot* numbers,  $p(s, t)$  and  $\kappa(s, t)$ , respectively. [5 marks]

A graph  $G$  is called  $k$ -*connected* if  $k \leq \kappa(G)$ . Show that a graph is  $k$ -connected iff any two of its vertices are joined by a thread of ply number  $k$ , as per your above statement of Menger's theorem in terms of ply and clot numbers [3 marks].

5. Show that  $K_5$  and the Petersen graph are not planar graphs.

## 6. An easy application for bipartite graphs

Use Hall's theorem show that a bipartite graph  $G(V, E)$  has a perfect matching if and only if  $|N(S)| \geq |S|$  for every  $S \subseteq V$ , where  $V$  is the set of all the vertices of the graph  $G(V, E)$ .

7. Every  $k$ -regular,  $(k-1)$ -connected graph with even order has a 1-factor.
8. If  $G$  is a  $k$ -regular,  $(k-1)$ -connected graph of even order and  $G'$  is obtained from  $G$  by removing any  $(k-1)$  edges then  $G'$  has a 1-factor.