
Tutorial 6: One-way Functions

Submission Guidelines All problems must be solved in class today. Searching on the internet for solutions is strictly discouraged.

1. Show that any IND-EAV-secure private-key encryption scheme implies the existence of a one-way function.

A: Refer to Katz-Lindell book (section 6.7) for the proof.

2. Suppose $f : \{0, 1\}^n \rightarrow \{0, 1\}^n$ is a one-way function. Prove that the function $g : \{0, 1\}^{2n} \rightarrow \{0, 1\}^{2n}$ defined as $g(x_1, x_2) = (f(x_1), x_2)$ for $x_1, x_2 \in \{0, 1\}^n$ is also one-way.

A: Let \mathcal{A}_g be an inverting adversary for g . We construct an adversary \mathcal{A}_f that inverts f .

Description of \mathcal{A}_f :

- \mathcal{A}_f received $f(x)$ for some $x \xleftarrow{U} \{0, 1\}^n$.
- It then picks $x_2 \xleftarrow{U} \{0, 1\}^n$ and provides $g(x, x_2) = (f(x), x_2)$ to \mathcal{A}_g .
- \mathcal{A}_g sends some $x', x'_2 \in \{0, 1\}^n$ and halts.
- \mathcal{A}_f just relays x' to its challenger and terminates.

Clearly,

$$\Pr[\mathcal{A}_f \text{ wins}] = \Pr[f(x') = f(x)] = \Pr[g(x', x'_2) = g(x, x_2)] = \Pr[\mathcal{A}_g \text{ wins}]$$

thus implying that if f is one-way, then so is g .

3. Show (formally) that if a one-to-one function has a hard-core predicate, then it is one-way. Where exactly do you need the one-to-one property?

A: Let $f : X \rightarrow Y$ be a 1-1 function and let hc be a hard-core predicate for f . Let \mathcal{A} be an adversary inverting f . We show how to build a prediction adversary \mathcal{B} for hc .

Description of \mathcal{B} :

- Receives from its challenger $y = f(x)$ for some $x \xleftarrow{U} X$.
- Provides y to \mathcal{A} and receives $x' \in X$ in return.
- Returns $\text{hc}(x')$ and halts.

If \mathcal{A} wins, then $f(x') = f(x)$. f is 1-1 implies that $x = x'$ and hence $\text{hc}(x) = \text{hc}(x')$. In this case, \mathcal{B} wins. We therefore can conclude that hc is (ϵ, t) -hardcore predicate for f implies that f is $(\epsilon, t - O(1))$ one-way.