
Tutorial 1: Perfect Secrecy

Submission Guidelines All problems must be solved in class today. Searching on the internet for solutions is strictly discouraged.

- The shift cipher (also called Caesar cipher) works as follows. The English alphabet is represented by numbers from 0 to 25 i.e., $\{A, B, \dots, Z\}$ are mapped to $\{0, 1, \dots, 25\}$ in the same order.

Define $\mathcal{K} = \{0, 1, 2, \dots, 25\}$, $\mathcal{M} = \mathcal{C} = \{0, 1, 2, \dots, 25\}^*$

$\text{Gen}()$: $k \xleftarrow{\text{U}} \mathcal{K}$

$\text{Enc}(k, \mathbf{m} = m_1 m_2 \dots m_n)$: Set $c_i \leftarrow m_i + k \pmod{26}$. Ciphertext is given by $\mathbf{c} = c_1 c_2 \dots c_n$.

$\text{Dec}(k, \mathbf{c} = c_1 c_2 \dots c_n)$: Recover message components as $m_i \leftarrow c_i - k \pmod{26}$.

- Is it perfectly secret?
 - Can you modify the description to make it perfectly secret?
- Prove or refute: Every encryption scheme for which the size of the key space \mathcal{K} equals the size of the message space \mathcal{M} and for which the key is chosen uniformly from \mathcal{K} , is perfectly secret.
 - Let $\varepsilon < 1$ be a constant. An encryption scheme $\mathcal{E} = (\text{Gen}, \text{Enc}, \text{Dec})$ over $(\mathcal{K}, \mathcal{M}, \mathcal{C})$ is called ε -almost perfectly secret if for any distribution over \mathcal{M} , any $m \in \mathcal{M}$ and any $c \in \mathcal{C}$,

$$|\Pr[M = m | C = c] - \Pr[M = m]| < \varepsilon.$$

Show that if \mathcal{E} is an ε -almost perfectly secret encryption scheme, then $|\mathcal{K}| \geq (1 - \varepsilon)|\mathcal{M}|$.

- Modify the one-time pad as follows. Let $\mathcal{K} \subseteq \{0, 1\}^\ell$ and the key generation be such that first a key \tilde{k} is uniformly chosen from $\{0, 1\}^{\ell/2}$ and then the key is defined by $k = \tilde{k} \parallel \tilde{k}$ where \parallel denotes concatenation. (In addition, define $\mathcal{M} = \mathcal{C} = \{0, 1\}^\ell$.) Is this scheme perfectly secure? Justify your answer.