
Decidability

CS60001: Foundations of Computing Science



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Decidable Problems Concerning Regular Languages

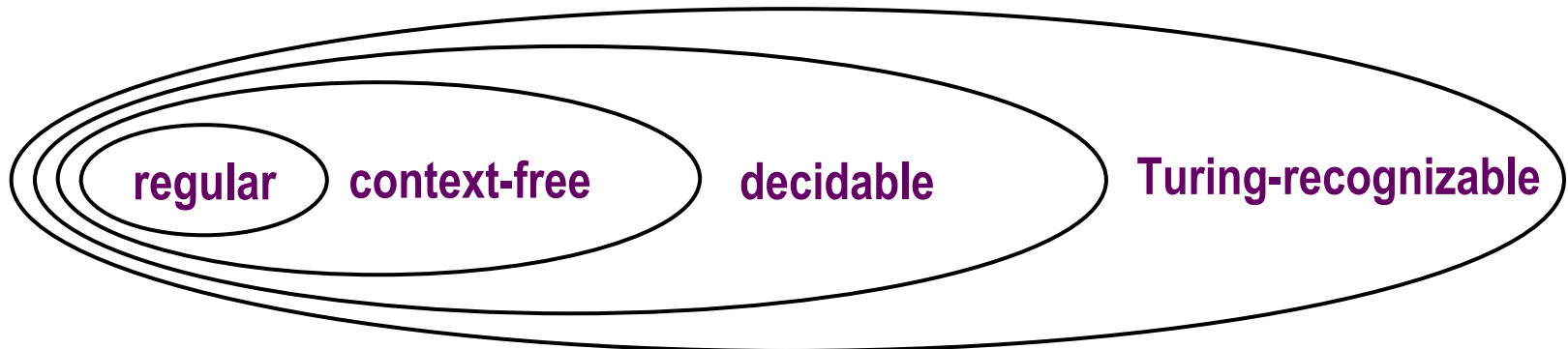
□ Theorems:

- Let $A_{\text{DFA}} = \{\langle B, w \rangle \mid B \text{ is a DFA that accepts input string } w\}$,
then A_{DFA} is a decidable language
- Let $A_{\text{NFA}} = \{\langle B, w \rangle \mid B \text{ is a NFA that accepts input string } w\}$,
then A_{NFA} is a decidable language
- Let $A_{\text{REX}} = \{\langle R, w \rangle \mid R \text{ is a regular expression that generates string } w\}$,
then A_{REX} is a decidable language
- Let $E_{\text{DFA}} = \{\langle A \rangle \mid A \text{ is a DFA and } L(A) = \Phi\}$,
then E_{DFA} is a decidable language (*emptiness testing*)
- Let $EQ_{\text{DFA}} = \{\langle A, B \rangle \mid A \text{ \& } B \text{ are DFAs and } L(A) = L(B)\}$,
then EQ_{DFA} is a decidable language

Decidable Problems Concerning Context-free Languages

□ Theorems:

- Let $A_{CFG} = \{\langle G, w \rangle \mid G \text{ is a CFG that generates string } w\}$, then A_{CFG} is a decidable language
- Let $E_{CFG} = \{\langle G \rangle \mid G \text{ is a CFG and } L(G) = \Phi\}$, then E_{CFG} is a decidable language (*emptiness testing*)
- Let $EQ_{CFG} = \{\langle G, H \rangle \mid G \text{ \& } H \text{ are CFGs and } L(G) = L(H)\}$, then EQ_{CFG} is a decidable language
- Every context-free language is decidable



The relationship among classes of languages

The Halting Problem

- Let $A_{TM} = \{ \langle M, w \rangle \mid M \text{ is a TM and } M \text{ accepts } w \}$
 - A_{TM} is *Turing-recognizable*
 - A_{TM} is *undecidable*
- The Diagonalization Method [Georg Cantor, 1873]
 - Definitions:
 - A function that is both one-to-one and onto is called a *correspondence*
 - A set is *countable* if either it is finite or it has the same size as \mathcal{N}
 - Example (Theorem): The set of real numbers (\mathcal{R}) is *uncountable*
 - Corollary: Some languages are not *Turing-recognizable*
- The Halting Problem (A_{TM}) is *undecidable*
- A Turing-unrecognizable Language
 - Theorem: A language is *decidable* if and only if it is *Turing-recognizable* and *co-Turing-recognizable*
 - Corollary: $\overline{A_{TM}}$ is *not* Turing-recognizable