

SPARC Assembly Language Reference Manual

Sun Microsystems, Inc. 4150 Network Circle Santa Clara, CA 95054 U.S.A.

Part No: 816-1681-10 May 2002 Copyright 2002 Sun Microsystems, Inc. 4150 Network Circle, Santa Clara, CA 95054 U.S.A. All rights reserved.

This product or document is protected by copyright and distributed under licenses restricting its use, copying, distribution, and decompilation. No part of this product or document may be reproduced in any form by any means without prior written authorization of Sun and its licensors, if any. Third-party software, including font technology, is copyrighted and licensed from Sun suppliers.

Parts of the product may be derived from Berkeley BSD systems, licensed from the University of California. UNIX is a registered trademark in the U.S. and other countries, exclusively licensed through X/Open Company, Ltd.

Sun, Sun Microsystems, the Sun logo, docs.sun.com, AnswerBook, AnswerBook2, and Solaris are trademarks, registered trademarks, or service marks of Sun Microsystems, Inc. in the U.S. and other countries. All SPARC trademarks are used under license and are trademarks or registered trademarks of SPARC International, Inc. in the U.S. and other countries. Products bearing SPARC trademarks are based upon an architecture developed by Sun Microsystems. Inc.

The OPEN LOOK and Sun^{TM} Graphical User Interface was developed by Sun Microsystems, Inc. for its users and licensees. Sun acknowledges the pioneering efforts of Xerox in researching and developing the concept of visual or graphical user interfaces for the computer industry. Sun holds a non-exclusive license from Xerox to the Xerox Graphical User Interface, which license also covers Sun's licensees who implement OPEN LOOK GUIs and otherwise comply with Sun's written license agreements.

Federal Acquisitions: Commercial Software-Government Users Subject to Standard License Terms and Conditions.

DOCUMENTATION IS PROVIDED "AS IS" AND ALL EXPRESS OR IMPLIED CONDITIONS, REPRESENTATIONS AND WARRANTIES, INCLUDING ANY IMPLIED WARRANTY OF MERCHANTABILITY, FITNESS FOR A PARTICULAR PURPOSE OR NON-INFRINGEMENT, ARE DISCLAIMED, EXCEPT TO THE EXTENT THAT SUCH DISCLAIMERS ARE HELD TO BE LEGALLY INVALID.

Copyright 2002 Sun Microsystems, Inc. 4150 Network Circle, Santa Clara, CA 95054 U.S.A. Tous droits réservés.

Ce produit ou document est protégé par un copyright et distribué avec des licences qui en restreignent l'utilisation, la copie, la distribution, et la décompilation. Aucune partie de ce produit ou document ne peut être reproduite sous aucune forme, par quelque moyen que ce soit, sans l'autorisation préalable et écrite de Sun et de ses bailleurs de licence, s'il y en a. Le logiciel détenu par des tiers, et qui comprend la technologie relative aux polices de caractères, est protégé par un copyright et licencié par des fournisseurs de Sun.

Des parties de ce produit pourront être dérivées du système Berkeley BSD licenciés par l'Université de Californie. UNIX est une marque déposée aux Etats-Unis et dans d'autres pays et licenciée exclusivement par X/Open Company, Ltd.

Sun, Sun Microsystems, le logo Sun, docs.sun.com, AnswerBook, AnswerBook2, et Solaris sont des marques de fabrique ou des marques déposées, ou marques de service, de Sun Microsystems, Inc. aux Etats-Unis et dans d'autres pays. Toutes les marques SPARC sont utilisées sous licence et sont des marques de fabrique ou des marques déposées de SPARC International, Inc. aux États-Unis et dans d'autres pays. Les produits portant les marques SPARC sont basés sur une architecture développée par Sun Microsystems, Inc.

L'interface d'utilisation graphique OPEN LOOK et SunTM a été développée par Sun Microsystems, Inc. pour ses utilisateurs et licenciés. Sun reconnaît les efforts de pionniers de Xerox pour la recherche et le développement du concept des interfaces d'utilisation visuelle ou graphique pour l'industrie de l'informatique. Sun détient une licence non exclusive de Xerox sur l'interface d'utilisation graphique Xerox, cette licence couvrant également les licenciés de Sun qui mettent en place l'interface d'utilisation graphique OPEN LOOK et qui en outre se conforment aux licences écrites de Sun.

CETTE PUBLICATION EST FOURNIE "EN L'ETAT" ET AUCUNE GARANTIE, EXPRESSE OU IMPLICITE, N'EST ACCORDEE, Y COMPRIS DES GARANTIES CONCERNANT LA VALEUR MARCHANDE, L'APTITUDE DE LA PUBLICATION A REPONDRE A UNE UTILISATION PARTICULIERE, OU LE FAIT QU'ELLE NE SOIT PAS CONTREFAISANTE DE PRODUIT DE TIERS. CE DENI DE GARANTIE NE S'APPLIQUERAIT PAS, DANS LA MESURE OU IL SERAIT TENU JURIDIQUEMENT NUL ET NON AVENU.





Contents

SPARC Assembler for SunOS 5.x

Preface 7

```
1.1 Introduction
1.2 Operating Environment 11
1.3 SPARC Assembler for SunOS 4.1 Versus SunOS 5.x
    1.3.1 Labeling Format
    1.3.2 Object File Format
                              12
    1.3.3 Pseudo-Operations
    1.3.4 Command Line Options
                                   12
Assembler Syntax
2.1 Syntax Notation
                      13
2.2 Assembler File Syntax
                            14
    2.2.1 Lines Syntax
    2.2.2 Statement Syntax
                            14
2.3 Lexical Features
    2.3.1 Case Distinction
                            14
    2.3.2 Comments
    2.3.3 Labels
    2.3.4 Numbers
                     15
    2.3.5 Strings
                  15
    2.3.6 Symbol Names
                           16
    2.3.7 Special Symbols - Registers
                                      16
    2.3.8 Operators and Expressions
                                      18
```

11

	2.3.9 SPA	ARC V9 Op	erator	s and Ex	pressions	19
2.4	Assemble	er Error Me	essages	20		
_	. 11	17. 1.				

Executable and Linking Format 24

- 3.1 ELF Header
- 3.2 Sections
 - 3.2.1 Section Header 26
 - 3.2.2 Predefined User Sections
 - 3.2.3 Predefined Non-User Sections 31
- 32 3.3 Locations
- 3.4 Addresses 32
- 3.5 Relocation Tables 33
- 3.6 Symbol Tables 33
- 3.7 String Tables 34
- 3.8 Assembler Directives
 - 3.8.1 Section Control Directives 35
 - 3.8.2 Symbol Attribute Directives 35
 - 3.8.3 Assignment Directive
 - 3.8.4 Data Generating Directives 36

Converting Files to the New Format 37

- 4.1 Introduction 37
- 4.2 Conversion Instructions 37
- 4.3 Examples 38

Instruction-Set Mapping 39

- 5.1 Table Notation
- 5.2 Integer Instructions
- 5.3 Floating-Point Instruction 48
- 5.4 Coprocessor Instructions 50
- 5.5 Synthetic Instructions
- 5.6 V8/V9 Natural Pseudo Instructions 53

Pseudo-Operations 55

A.1 Alphabetized Listing with Descriptions 55

В	Examples of Pseudo-Operations 63				
	B.1 Example 1 63				
	B.2 Example 2 64				
	B.3 Example 3 64				
	B.4 Example 4 65				
	B.5 Example 5 65				
C	Using the Assembler Command Line 67				
	C.1 Assembler Command Line 67				
	C.2 Assembler Command Line Options 68				
	C.3 Disassembling Object Code 71				
D	An Example Language Program 73				
E	SPARC-V9 Instruction Set 79				
	E.1 SPARC-V9 Changes 79				
	E.1.1 Registers 79				
	E.1.2 Alternate Space Access 81				
	E.1.3 Byte Order 81				
	E.2 SPARC-V9 Instruction Set Changes 81				
	E.2.1 Extended Instruction Definitions to Support the 64-bit Model 82				
	E.2.2 Added Instructions to Support 64 bits 82				
	E.2.3 Added Instructions to Support High-Performance System Implementation 83				
	E.2.4 Deleted Instructions 83				
	E.2.5 Miscellaneous Instruction Changes 84				
	E.3 SPARC-V9 Instruction Set Mapping 84				
	E.4 SPARC-V9 Floating-Point Instruction Set Mapping 93				
	E.5 SPARC-V9 Synthetic Instruction-Set Mapping 95				
	E.6 UltraSPARC and VIS Instruction Set Extensions 97				
	E.6.1 Graphics Data Formats 97				
	E.6.2 Eight-bit Format 97				
	E.6.3 Fixed Data Formats 98				
	E.6.4 SHUTDOWN Instruction 98				
	E.6.5 Graphics Status Register (GSR) 98				
	E.6.6 Graphics Instructions 99				
	E.6.7 Memory Access Instructions 103				

Preface

The SunOS assembler that runs on the SPARC operating environment, referred to as the "SunOS SPARC" in this manual, translates source files that are in assembly language format into object files in linking format.

In the program development process, the assembler is a tool to use in producing program modules intended to exploit features of the SPARC architecture in ways that cannot be easily done using high level languages and their compilers.

Whether assembly language is chosen for the development of program modules depends on the extent to which and the ease with which the language allows the programmer to control the architectural features of the processor.

The assembly language described in this manual offers full direct access to the SPARC instruction set. The assembler may also be used in connection with SunOS 5.x macro preprocessors to achieve full macro-assembler capability. Furthermore, the assembler responds to directives that allow the programmer direct control over the contents of the relocatable object file.

This document describes the language in which the source files must be written. The nature of the machine mnemonics governs the way in which the program's executable portion is written. This document includes descriptions of the pseudo operations that allow control over the object file. This facilitates the development of programs that are easy to understand and maintain.

Before You Read This Book

You should also become familiar with the following:

- Manual pages: as(1), ld(1), cpp(1), elf(3f), dis(1), a.out(1)
- *SPARC Architecture Manual* (Version 8 and Version 9)

■ System V Application Binary Interface: SPARCTM Processor Supplement

How This Book is Organized

This book is organized as follows:

Chapter 1 discusses features of the SunOS 5.x SPARC Assembler.

Chapter 2 describes the syntax of the SPARC assembler that takes assembly programs and produces relocatable object files for processing by the link editor.

Chapter 3 describes the relocatable ELF files that hold code and data suitable for linking with other object files.

Chapter 4 describes how to convert existing SunOS 4.1 SPARC assembly files to the SunOS 5.x assembly file format.

Chapter 5 describes the relationship between hardware instructions of the SPARC architecture and the assembly language instruction set.

Appendix A lists the pseudo-operations supported by the SPARC assembler.

Appendix B shows some examples of ways to use various pseudo-operations.

Appendix C describes the available assembler command-line options.

Appendix D describes an example C language program with comments to show correspondence between the assembly code and the C code.

Appendix E describes the SPARC-V9 instruction set and the changes due to the SPARC-V9 implementation.

Accessing Sun Documentation Online

The docs.sun.comSM Web site enables you to access Sun technical documentation online. You can browse the docs.sun.com archive or search for a specific book title or subject. The URL is http://docs.sun.com.

Typographic Conventions

The following table describes the typographic changes used in this book.

TABLE P-1 Typographic Conventions

Typeface or Symbol	Meaning	Example
AaBbCc123	The names of commands, files, and directories; on-screen computer output	Edit your .login file. Use ls -a to list all files. machine_name% you have mail.
AaBbCc123	What you type, contrasted with on-screen computer output	machine_name% su Password:
AaBbCc123	Command-line placeholder: replace with a real name or value	To delete a file, type rm filename.
AaBbCc123	Book titles, new words, or terms, or words to be emphasized.	Read Chapter 6 in <i>User's Guide</i> . These are called <i>class</i> options. You must be <i>root</i> to do this.

Shell Prompts in Command Examples

The following table shows the default system prompt and superuser prompt for the C shell, Bourne shell, and Korn shell.

TABLE P-2 Shell Prompts

Shell	Prompt
C shell prompt	machine_name%
C shell superuser prompt	machine_name#
Bourne shell and Korn shell prompt	ş
Bourne shell and Korn shell superuser prompt	#

SPARC Assembler for SunOS 5.x

1.1 Introduction

This chapter discusses features of the *SunOS 5.x* SPARC assembler. This document is distributed as part of the developer documentation set with every SunOS operating system release.

This document is also distributed with the on-line documentation set for the convenience of SPARCworksTM and SPARCompilerTM 4.0 users who have products that run on the SunOS 5.x operating system. It is included as part of the SPARCworks/SPARCompiler Floating Point and Common Tools AnswerBook, which is the on-line information retrieval system.

This document contains information from *The SPARC Architecture Manual*, Version 8. Information about Version 9 support is summarized in Appendix E.

1.2 Operating Environment

The SunOS SPARC assembler runs under the SunOS 5.x operating system or the SolarisTM 2.x operating environment. SunOS 5.x refers to SunOS 5.2 operating system and later releases. Solaris 2.x refers to the Solaris 2.2 operating environment and later releases.

1.3 SPARC Assembler for SunOS 4.1 Versus SunOS 5.x

This section describes the differences between the SunOS 4.1 SPARC assembler and the SunOS 5.x SPARC assembler.

1.3.1 Labeling Format

- Symbol names beginning with a dot (.) are assumed to be local symbols.
- Names beginning with an underscore (_) are reserved by ANSI C.

1.3.2 Object File Format

The type of object files created by the SPARC assembler are ELF (Executable and Linking Format) files. These relocatable object files hold code and data suitable for linking with other object files to create an executable file or a shared object file, and are the assembler normal output.

1.3.3 Pseudo-Operations

See Appendix A for a detailed description of the pseudo-operations (pseudo-ops).

1.3.4 Command Line Options

See Appendix C for a detailed description of command line options and a list of SPARC architectures.

Assembler Syntax

The *SunOS 5.x* SPARC assembler takes assembly language programs, as specified in this document, and produces relocatable object files for processing by the *SunOS 5.x* SPARC link editor. The assembly language described in this document corresponds to the SPARC instruction set defined in the *SPARC Architecture Manual* (Version 8 and Version 9) and is intended for use on machines that use the SPARC architecture.

This chapter is organized into the following sections:

- "2.1 Syntax Notation" on page 13
- "2.2 Assembler File Syntax" on page 14
- "2.3 Lexical Features" on page 14
- "2.4 Assembler Error Messages" on page 20

2.1 Syntax Notation

In the descriptions of assembly language syntax in this chapter:

- Brackets ([]) enclose optional items.
- Asterisks (*) indicate items to be repeated zero or more times.
- Braces ({ }) enclose alternate item choices, which are separated from each other by vertical bars (|).
- Wherever blanks are allowed, arbitrary numbers of blanks and horizontal tabs may be used. Newline characters are not allowed in place of blanks.

2.2 Assembler File Syntax

The syntax of assembly language files is:

[line] *

2.2.1 Lines Syntax

The syntax of assembly language lines is:

```
[statement [ ; statement]*] [!comment]
```

2.2.2 Statement Syntax

The syntax of an assembly language statement is:

```
[label:] [instruction]
where:
label
  is a symbol name.
instruction
  is an encoded pseudo-op, synthetic instruction, or instruction.
```

2.3 Lexical Features

This section describes the lexical features of the assembler syntax.

2.3.1 Case Distinction

Uppercase and lowercase letters are distinct everywhere *except* in the names of special symbols. Special symbol names have no case distinction.

2.3.2 Comments

A comment is preceded by an exclamation mark character (!); the exclamation mark character and all following characters up to the end of the line are ignored. C language-style comments ("/*...*/") are also permitted and may span multiple lines.

2.3.3 Labels

A label is either a symbol or a single decimal digit n (0...9). A label is immediately followed by a *colon* (:).

Numeric labels may be defined repeatedly in an assembly file; normal symbolic labels may be defined only once.

A numeric label *n* is referenced after its definition (backward reference) as *n*b, and before its definition (forward reference) as *n*f.

2.3.4 Numbers

Decimal, hexadecimal, and octal numeric constants are recognized and are written as in the C language. However, integer suffixes (such as L) are not recognized.

For floating-point pseudo-operations, floating-point constants are written with 0r or OR (where r or R means REAL) followed by a string acceptable to atof(3); that is, an optional sign followed by a non-empty string of digits with optional decimal point and optional exponent.

The special names Ornan and Orinf represent the special floating-point values Not-A-Number (NaN) and INFinity. Negative Not-A-Number and Negative INFinity are specified as Or-nan and Or-inf.

Note – The names of these floating-point constants begin with the digit zero, *not* the letter "O."

2.3.5 Strings

A string is a sequence of characters quoted with either double-quote mark (") or single-quote mark (') characters. The sequence must not include a *newline* character. When used in an expression, the numeric value of a string is the numeric value of the ASCII representation of its first character.

The suggested style is to use *single quote mark* characters for the ASCII value of a single character, and double quote mark characters for quoted-string operands such as used by pseudo-ops. An example of assembly code in the suggested style is:

```
add %g1,'a'-'A',%g1 ! g1 + ('a' - 'A') --> g1
```

The escape codes described in Table 2–1, derived from ANSI C, are recognized in strings.

TABLE 2-1

Escape Code	Description
\a	Alert
\ b	Backspace
\f	Form feed
\n	Newline (line feed)
\r	Carriage return
\t	Horizontal tab
\v	Vertical tab
\nnn	Octal value nnn
\xnn	Hexadecimal value nn

2.3.6 Symbol Names

The syntax for a symbol *name* is:

```
{ letter | _ | $ | . } { letter | _ | $ | . | digit }*
```

In the above syntax:

- Uppercase and lowercase letters are distinct; the underscore (_), dollar sign (\$), and dot (.) are treated as alphabetic characters.
- Symbol names that begin with a dot (.) are assumed to be local symbols. To simplify debugging, avoid using this type of symbol name in hand-coded assembly language routines.
- The symbol dot (.) is predefined and always refers to the address of the beginning of the current assembly language statement.
- External variable names beginning with the underscore character are reserved by the ANSI C Standard. Do *not* begin these names with the underscore; otherwise, the program will not conform to ANSI C and unpredictable behavior may result.

2.3.7 Special Symbols - Registers

Special symbol names begin with a *percentage sign* (%) to avoid conflict with user symbols. Table 2–2 lists these special symbol names.

TABLE 2-2

Symbol Object	Name	Comment
General-purpose registers	%r0 %r31	
General-purpose global registers	%g0 %g7	Same as %r0 %r7
General-purpose out registers	%00 %07	Same as %r8 %r15
General-purpose local registers	%10 %17	Same as %r16 %r23
General-purpose in registers	%i0 %i7	Same as %r24 %r31
Stack-pointer register	%sp	(%sp = %o6 = %r14)
Frame-pointer register	%fp	(%fp = %i6 = %r30)
Floating-point registers	%f0 %f31	
Floating-point status register	%fsr	
Front of floating-point queue	%fq	
Coprocessor registers	%c0 %c31	
Coprocessor status register	%csr	
Coprocessor queue	%cq	
Program status register	%psr	
Trap vector base address register	%tbr	
Window invalid mask	%wim	
Y register	%y	
Unary operators	%10	Extracts least significant 10 bits
	%hi	Extracts most significant 22 bits
	%r_disp32	Used only in Sun compiler-generated code.
	%r_plt32	Used only in Sun compiler-generated code.
Ancillary state registers	%asr1 %asr31	

There is no case distinction in special symbols; for example,

%PSR

is equivalent to

%psr

The suggested style is to use lowercase letters.

The lack of case distinction allows for the use of non-recursive preprocessor substitutions, for example:

```
#define psr %PSR
```

The special symbols %hi and %lo are true unary operators which can be used in any expression and, as other unary operators, have higher precedence than binary operations. For example:

```
hi a+b = (hi a)+b

a+b = (hi a)+b
```

To avoid ambiguity, enclose operands of the %hi or %10 operators in parentheses. For example:

```
%hi(a) + b
```

2.3.8 Operators and Expressions

The operators described in Table 2–3 are recognized in constant expressions.

TABLE 2-3

Binary	Operators	Unary	Operators
+	Integer addition	+	(No effect)
_	Integer subtraction	_	2's Complement
*	Integer multiplication	~	1's Complement
/	Integer division	%lo(address)	Extract least significant 10 bits as computed by: (address & 0x3ff)
%	Modulo	%hi(address)	Extract most significant 22 bits as computed by: (address >>10)
۸	Exclusive OR	%r_disp32 %r_disp64	Used in Sun compiler-generated code only to instruct the assembler to generate specific relocation information for the given expression.

TABLE 2-3 (Continued)

Binary	Operators	Unary	Operators
<<	Left shift	%r_plt32 %r_plt64	Used in Sun compiler-generated code only to instruct the assembler to generate specific relocation information for the given expression.
>>	Right shift		
&	Bitwise AND		
1	Bitwise OR		

Since these operators have the same precedence as in the C language, put expressions in parentheses to avoid ambiguity.

To avoid confusion with register names or with the %hi, %lo, %r_disp32/64, or %r plt32/64 operators, the modulo operator % must not be immediately followed by a letter or digit. The modulo operator is typically followed by a space or left parenthesis character.

2.3.9 SPARC V9 Operators and Expressions

The following V9 64-bit operators and expressions in Table 2-4 ease the task of converting from V8/V8plus assembly code to V9 assembly code.

TABLE 2-4

Unary	Calculation	Operators
%hh	(address) >> 42	Extract bits 42-63 of a 64-bit word
%hm	((address) >> 32) & 0x3ff	Extract bits 32-41 of a 64-bit word
%lm	(((address) >> 10) & 0x3fffff)	Extract bits 10-31 of a 64-bit word

For example:::

```
sethi %hh (address), %11
or %11, %hm (address), %11
sethi %lm (address), %12
or %12, %lo (address), %12
sllx %11, 32, %11
or %11, %12, %11
```

The V9 high 32-bit operators and expressions are identified in Table 2–5.

TABLE 2-5

Unary	Calculation	Operators
%hix	((((address) ^ 0xffffffffffffff >> 10) &0x4fffff)	Invert every bit and extract bits 10-31
%lox	((address) & 0x3ff 0x1c00	Extract bits 0-9 and sign extend that to 13 bits

For example:

```
%sethi %hix (address), %l1
or %l1, %lox (address), %l1
```

The V9 low 44-bit operators and expressions are identified in Table 2–6.

TABLE 2-6

Unary	Calculation	Operators
%h44	((address) >> 22)	Extract bits 22-43 of a 64-bit word
%m44	((address) >> 12) & 0x3ff	Extract bits 12-21 of a 64-bit word
144	(address) & 0xfff	Extract bits 0-11 of a 64-bit word

For example::

```
%sethi %h44 (address), %l1
or %l1, %m44 (address), %l1
sllx %l1, 12, %l1
or %l1, %144 (address), %l1
```

2.4 Assembler Error Messages

Messages generated by the assembler are generally self-explanatory and give sufficient information to allow correction of a problem.

Certain conditions will cause the assembler to issue warnings associated with delay slots following Control Transfer Instructions (CTI). These warnings are:

- Set synthetic instructions in delay slots
- Labels in delay slots
- Segments that end in control transfer instructions

These warnings point to places where a problem could exist. If you have intentionally written code this way, you can insert an .empty pseudo-operation immediately after the control transfer instruction.

The $% \left(1\right) =\left(1\right) \left(1\right)$. empty pseudo-operation in a delay slot tells the assembler that the delay slot can be empty or can contain whatever follows because you have verified that either the code is correct or the content of the delay slot does not matter.

Executable and Linking Format

The type of object files created by the SPARC assembler version for *SunOS 5.x* are now *Executable and Linking Format* (ELF) files. These relocatable ELF files hold code and data suitable for linking with other object files to create an executable or a shared object file, and are the assembler normal output. The assembler can also write information to standard output (for example, under the -S option) and to standard error (for example, under the -V option). The SPARC assembler creates a default output file when standard input or multiple files are used.

This chapter is organized into the following sections:

- "3.1 ELF Header" on page 24
- "3.2 Sections" on page 25
- "3.3 Locations" on page 32
- "3.5 Relocation Tables" on page 33
- "3.6 Symbol Tables" on page 33
- "3.4 Addresses" on page 32
- "3.7 String Tables" on page 34
- "3.8 Assembler Directives" on page 35

The ELF object file format consists of:

- Header
- Sections
- Locations
- Addresses
- Relocation tables
- Symbol tables
- String tables

For more information, see the *System V Application Binary Interface: SPARC*TM *Processor Supplement*.

3.1 ELF Header

The *ELF header* is always located at the beginning of the ELF file. It describes the ELF file organization and contains the actual sizes of the object file control structures. The initial bytes of an ELF header specify how the file is to be interpreted.

The ELF header contains the following information:

ehsize

ELF header size in bytes.

entry

Virtual address at which the process is to start. A value of 0 indicates no associated entry point.

flag

Processor-specific flags associated with the file.

ident

Marks the file as an object file and provides machine-independent data to decode and interpret the file contents.

machine

Specifies the required architecture for an individual file. A value of 2 specifies SPARC.

phentsize

Size in bytes of entries in the program header table. All entries are the same size.

phnum

Number of entries in program header table. A value of 0 indicates the file has no program header table.

phoff

Program header table file offset in bytes. The value of 0 indicates no program header.

shentsize

Size in bytes of the section header. A section header is one entry in the section header table; all entries are the same size.

shnum

Number of entries in section header table. A value of 0 indicates the file has no section header table.

shoff

Section header table file offset in bytes. The value of 0 indicates no section header.

shstrndx

Section header table index of the entry associated with the section name string table. A value of SHN_UNDEF indicates the file does not have a section name string

type

Identifies the object file type. Table 3–1 describes the reserved object file types.

version

Identifies the object file version.

Table 3–1 shows reserved object file types:

TABLE 3-1

Туре	Value	Description
none	0	No file type
rel	1	Relocatable file
exec	2	Executable file
dyn	3	Shared object file
core	4	Core file
loproc	0xff00	Processor-specific
hiproc	0xffff	Processor-specific

3.2 Sections

A section is the smallest unit of an object that can be relocated. The following sections are commonly present in an ELF file:

- Section header
- Executable text
- Read-only data
- Read-write data
- Read-write uninitialized data (section header only)

Sections do not need to be specified in any particular order. The current section is the section to which code is generated.

These sections contain all other information in an object file and satisfy several conditions.

- 1. Every section must have one section header describing the section. However, a section header does not need to be followed by a section.
- 2. Each section occupies one contiguous sequence of bytes within a file. The section may be empty (that is, of zero-length).
- 3. A byte in a file can reside in only one section. Sections in a file cannot overlap.
- 4. An object file may have inactive space. The contents of the data in the inactive space are unspecified.

Sections can be added for multiple text or data segments, shared data, user-defined sections, or information in the object file for debugging.

Note – Not all of the sections need to be present.

3.2.1 Section Header

The *section header* allows you to locate all of the file sections. An entry in a section header table contains information characterizing the data in a section.

The section header contains the following information:

addr

Address at which the first byte resides if the section appears in the memory image of a process; the default value is 0.

addralign

Aligns the address if a section has an address alignment constraint; for example, if a section contains a double-word, the entire section must be ensured double-word alignment. Only 0 and positive integral powers of 2 are currently allowed. A value of 0 or 1 indicates no address alignment constraints.

entsize

Size in bytes for entries in fixed-size tables such as the symbol table.

flags

One-bit descriptions of section attributes. Table 3–2 describes the section attribute flags.

TABLE 3-2

Flag	Default Value	Description
SHF_WRITE	0x1	Contains data that is writable during process execution.
SHF_ALLOC	0x2	Occupies memory during process execution. This attribute is <i>off</i> if a control section does not reside in the memory image of the object file.
SHF_EXECINSTR	0x4	Contains executable machine instructions.
SHF_MASKPROC	0xf0000000	Reserved for processor-specific semantics.

info

Extra information. The interpretation of this information depends on the section type, as described in Table 3–3.

link

Section header table index link. The interpretation of this information depends on the section type, as described in Table 3–3.

Specifies the section name. An index into the section header string table section specifies the location of a null-terminated string.

offset

Specifies the byte offset from the beginning of the file to the first byte in the section.

Note – If the section type is SHT NOBITS, *offset* specifies the conceptual placement of the file.

size

Specifies the size of the section in bytes.

Note – If the section type is SHT NOBITS, *size* may be non-zero; however, the section still occupies no space in the file.

type

Categorizes the section contents and semantics. Table 3–3 describes the section types.

TABLE 3-3

Name Value		Description	Interpretation by	
			info	link
null	0	Marks section header as inactive.		
progbits	1	Contains information defined explicitly by the program.		
symtab	2	Contains a symbol table for link editing. This table may also be used for dynamic linking; however, it may contain many unnecessary symbols.	One greater than the symbol table index of the last local symbol.	The section header index of the associated string table.
		Note: Only one section of this type is allowed in a file		
strtab	3	Contains a string table. A file may have multiple string table sections.		
rela	4	Contains relocation entries with explicit addends. A file may have multiple relocation sections.	The section header index of the section to which the relocation applies.	The section header index of the associated symbol table.
hash	5	Contains a symbol rehash table.	0	The section header
		Note: Only one section of this type is allowed in a file		index of the symbol table to which the hash table applies.
dynamic	6	Contains dynamic linking information.	0	The section header index of the string
		Note: Only one section of this type is allowed in a file		table used by entries in the section.
note	7	Contains information that marks the file.		
nobits	8	Contains information defined explicitly by the program; however, a section of this type does not occupy any space in the file.		
rel	9	Contains relocation entries without explicit addends. A file may have multiple relocation sections.	The section header index of the section to which the relocation applies.	The section header index of the associated symbol table.

TABLE 3-3	(Continued)

Name	Value Description	Interpretation by		
			info	link
shlib	10	Reserved.		
dynsym	11	11 Contains a symbol table with a minimal set of symbols for dynamic linking. Note: Only one section of this type is allowed in a file	One greater than the symbol table index of the last local symbol.	The section header index of the associated string table.
loproc	0x7000000	Lower and upper bound of range reserved for processor-specific semantics.		
hiproc	0x7fffffff			
louser	0x80000000	Lower and upper bound of range reserved for application programs.		
hiuser	0xffffffff			
		Note: Section types in this range may be used by an application without conflicting with system-defined section types.		

Note – Some section header table indexes are reserved and the object file will not contain sections for these special indexes.

3.2.2 Predefined User Sections

A section that can be manipulated by the section control directives is known as a user section. You can use the section control directives to change the user section in which code or data is generated. Table 3-4 lists the predefined user sections that can be named in the section control directives.

TABLE 3-4

Section Name	Description
.bss	Section contains uninitialized read-write data.
.comment	Comment section.
.data & .data1	Section contains initialized read-write data.
.debug	Section contains debugging information.

(=	
Section Name	Description
.fini	Section contains runtime finalization instructions.
.init	Section contains runtime initialization instructions.
.rodata & .rodata1	Section contains read-only data.
.text	Section contains executable text.
.line	Section contains line # info for symbolic debugging.
.note	Section contains note information.

3.2.2.1 Creating an .init Section in an Object File

The .init sections contain codes that are to be executed before the the main program is executed. To create an .init section in an object file, use the assembler pseudo-ops shown in Example 3–1.

EXAMPLE 3-1 Creating an .init Section

```
.section ".init"
.align 4
<instructions>
```

At link time, the .init sections in a sequence of .o files are concatenated into an .init section in the linker output file. The code in the .init section are executed before the main program is executed.

Because the whole .init section is treated as a single function body, it is recommented that the only code added to these sections be in the following form:.

call routine_name
nop

The called routine should be located in another section. This will prevent conflicting register and stack usage within the .init sections.

3.2.2.2 Creating a .fini Section in an Object File

. fini sections contain codes that are to be executed after the the main program is executed. To create an .fini section in an object file, use the assembler pseudo-ops shown in Example 3–2.

EXAMPLE 3–2 Creating an .fini Section

```
.section ".fini"
.align 4
<instructions>
```

At link time, the .fini sections in a sequence of .o files are concatenated into a . fini section in the linker output file. The codes in the .fini section are executed after the main program is executed.

Because the whole .fini section is treated as a single function body, it is recommended that the only code added to these section be in the following form:.

```
call routine_name
nop
```

The called routine should be located in another section. This will prevent conflicting register and stack usage within the .fini sections.

3.2.3 Predefined Non-User Sections

Table 3–5 lists sections that are predefined but cannot be named in the section control directives because they are not under user control.

TABLE 3-5

Section Name	Description
".dynamic"	Section contains dynamic linking information.
.dynstr	Section contains strings needed for dynamic linking.
.dynsym	Section contains the dynamic linking symbol table.
.got	Section contains the global offset table.
.hash	Section contains a symbol hash table.
.interp	Section contains the path name of a program interpreter.
.plt	Section contains the procedure linking table.
.relname & .relaname	Section containing relocation information. <i>name</i> is the section to which the relocations apply, that is, ".rel.text", ".rela.text".

TABLE 3–5(Continued)

Section Name	Description
.shstrtab	String table for the section header table names.
.strtab	Section contains the string table.
.symtab	Section contains a symbol table.

3.3 Locations

A *location* is a specific position within a section. Each location is identified by a section and a byte offset from the beginning of the section. The *current location* is the location within the current section where code is generated.

A *location counter* tracks the current offset within each section where code or data is being generated. When a section control directive (for example, the .section pseudo-op) is processed, the location information from the location counter associated with the new section is assigned to and stored with the name and value of the current location.

The current location is updated at the end of processing each statement, but can be updated during processing of data-generating assembler directives (for example, the .word pseudo-op).

Note – Each section has one location counter; if more than one section is present, only one location can be current at any time.

3.4 Addresses

Locations represent *addresses in memory* if a section is allocatable; that is, its contents are to be placed in memory at program runtime. Symbolic references to these locations must be changed to addresses by the SPARC link editor.

3.5 Relocation Tables

The assembler produces a companion *relocation table* for each relocatable section. The table contains a list of relocations (that is, adjustments to data in the section) to be performed by the link editor.

3.6 Symbol Tables

A *symbol table* contains information to locate and relocate symbolic definitions and references. The SPARC assembler creates a symbol table section for the object file. It makes an entry in the symbol table for each symbol that is defined or referenced in the input file and is needed during linking. The symbol table is then used by the SPARC link editor during relocation. The section header contains the symbol table index for the first non-local symbol.

A symbol table contains the following information:

name

Index into the object file symbol string table. A value of zero indicates the symbol table entry has no name; otherwise, the value represents the string table index that gives the symbol name.

value

Value of the associated symbol. This value is dependent on the context; for example, it may be an address, or it may be an absolute value.

size

Size of symbol. A value of 0 indicates that the symbol has either no size or an unknown size.

info

Specifies the symbol type and binding attributes. Table 3–6 and Table 3–7 describe these values.

other

Undefined meaning. Current value is 0.

shndx

Contains the section header table index to another relevant section, if specified. As a section moves during relocation, references to the symbol will continue to point to the same location because the value of the symbol will change as well.

TABLE 3-6

Value	Туре	Description
0	notype	Type not specified.
1	object	<i>Symbol</i> is associated with a data object; for example, a variable or an array.
2	func	<i>Symbol</i> is associated with a function or other executable code. When another object file references a function from a shared object, the link editor automatically creates a procedure linkage table entry for the referenced symbol.
3	section	<i>Symbol</i> is associated with a section. These types of symbols are primarily used for relocation.
4	file	Gives the name of the source file associated with the object file.
13	loproc	Values reserved for processor-specific semantics.
15	hiproc	

Table 3–7 shows the symbol binding attributes.

TABLE 3-7

Value	Binding	Description
0	local	Symbol is defined in the object file and not accessible in other files. Local symbols of the same name may exist in multiple files.
1	global	<i>Symbol</i> is either defined externally or defined in the object file and accessible in other files.
2	weak	Symbol is either defined externally or defined in the object file and accessible in other files; however, these definitions have a lower precedence than globally defined symbols.
13	loproc	Values reserved for processor-specific semantics.
15	hiproc	

3.7 String Tables

A *string table* is a section which contains null-terminated variable-length character sequences, or strings, in the object file; for example, symbol names and file names. The strings are referenced in the section header as indexes into the string table section.

■ A string table index may refer to any byte in the section.

 Empty string table sections are permitted; however, the index referencing this section must contain zero.

A string may appear multiple times and may also be referenced multiple times. References to substrings may exist, and unreferenced strings are allowed.

3.8 Assembler Directives

Assembler directives, or pseudo-operations (pseudo-ops), are commands to the assembler that may or may not result in the generation of code. The different types of assembler directives are:

- Section Control Directives
- Symbol Attribute Directives
- Assignment Directives
- Data Generating Directives
- Optimizer Directives

See Appendix A for a complete description of the pseudo-ops supported by the SPARC assembler.

3.8.1 Section Control Directives

When a section is created, a section header is generated and entered in the ELF object file section header table. The *section control pseudo-ops* allow you to make entries in this table. Sections that can be manipulated with the section control directives are known as *user sections*. You can also use the section control directives to change the user section in which code or data is generated.

Note – The *symbol table, relocation table,* and *string table* sections are created implicitly. The section control pseudo-ops cannot be used to manipulate these sections.

The section control directives also create a section symbol which is associated with the location at the beginning of each created section. The section symbol has an offset value of zero.

3.8.2 Symbol Attribute Directives

The *symbol attribute* pseudo-ops declare the symbol type and size and whether it is local or global.

3.8.3 Assignment Directive

The *assignment* directive associates the value and type of expression with the symbol and creates a symbol table entry for the symbol. This directive constitutes a *definition* of the symbol and, therefore, must be the only definition of the symbol.

3.8.4 Data Generating Directives

The data generating directives are used for allocating storage and loading values.

Converting Files to the New Format

4.1 Introduction

This chapter discusses how to convert existing SunOS 4.1 SPARC assembly files to the *SunOS 5.x* SPARC assembly file format.

4.2 Conversion Instructions

- Remove the leading underscore (_) from symbol names. The *Solaris 2.x* SPARCompilers do not prepend a leading underscore to symbol names in the users' programs as did the SPARCompilers that ran under SunOS 4.1.
- Prefix local symbol names with a dot (.). Local symbol names in the *SunOS 5.x* SPARC assembly language begin with a dot (.) so that they will not conflict with user programs' symbol names.
- Change the usage of the pseudo-op .seg to .section, for example, change .seg data to .section .data. See Appendix A for more information.

4.3 Examples

Figure 4–1 shows how to convert an existing 4.1 file to the new format. The lines that are different in the new format are marked with change bars.

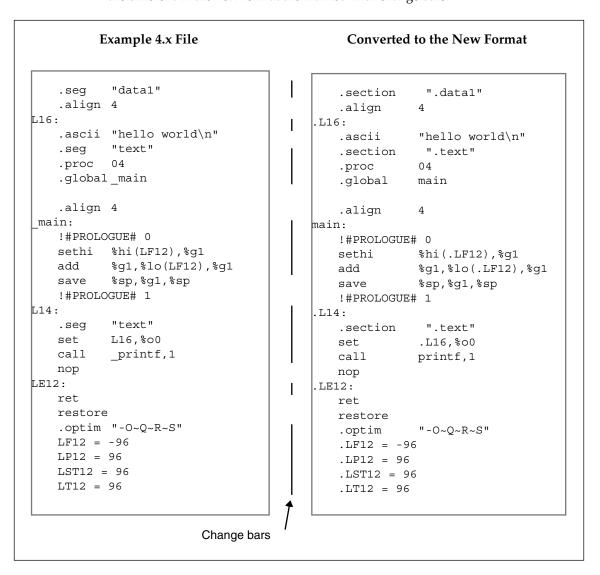


FIGURE 4–1 Converting a 4.x File to the New Format

Instruction-Set Mapping

The tables in this chapter describe the relationship between hardware instructions of the SPARC architecture, as defined in *The SPARC Architecture Manual* and the assembly language instruction set recognized by the *SunOS 5.x* SPARC assembler.

- "5.1 Table Notation" on page 39
- "5.2 Integer Instructions" on page 41
- "5.3 Floating-Point Instruction" on page 48
- "5.4 Coprocessor Instructions" on page 50
- "5.5 Synthetic Instructions" on page 51

The SPARC-V9 instruction set is described in Appendix E.

5.1 Table Notation

Table 5–1 shows the table notation used in this chapter to describe the instruction set of the assembler. The following notations are commonly suffixed to assembler mnemonics (uppercase letters refer to SPARC architecture instruction names.

TABLE 5-1

Notations	Describes	Comment
address	$reg_{rs1} + reg_{rs2}$ $reg_{rs1} + const13$ $reg_{rs1} - const13$	Address formed from register contents, immediate constant, or both.
	const13 + reg _{rs1}	
	const13	

(Continued) TABLE 5-1

Notations	Describes	Comment
asi		Alternate address space identifier; an unsigned 8-bit value. It can be the result of the evaluation of a symbol expression.
const13		A signed constant which fits in 13 bits. It can be the result of the evaluation of a symbol expression.
const22		A constant which fits in 22 bits. It can be the result of the evaluation of a symbol expression.
creg	%c0 %c31	Coprocessor registers.
freg	%f0 %f31	Floating-point registers.
imm7		A signed or unsigned constant that can be represented in 7 bits (it is in the range -64 127). It can be the result of the evaluation of a symbol expression.
reg	%r0 %r31	General purpose registers.
	%g0 %g7	Same as %r0 %r7 (Globals)
	%00 %07	Same as %r8 %r15 (Outs)
	%10 %17	Same as %r16 %r23 (Locals)
	%i0 %i7	Same as %r24 %r31 (Ins)
reg _{rd}		Destination register.
reg _{rs1} , reg _{rs2}		Source register 1, source register 2.
reg_or_imm	reg _{rs2} , const13	Value from either a single register, or an immediate constant.
regaddr	reg _{rs1} reg _{rs1} + reg _{rs2}	Address formed with register contents only.
Software_trap_number	reg _{rs1} + reg _{rs2}	A value formed from register contents, immediate
	$reg_{rs1} + imm7$	constant, or both. The resulting value must be in the range 0127, inclusive.
	reg_{rs1} - imm7	
	uimm7	
	imm7 + reg _{rs1}	
uimm7		An unsigned constant that can be represented in 7 bits (it is in the range 0 127). It can be the result of the evaluation of a symbol expression.

5.2 Integer Instructions

The notations described in Table 5–2 are commonly suffixed to assembler mnemonics (uppercase letters for architecture instruction names).

TABLE 5-2

Notation	Description
a	Instructions that deal with alternate space
b	Byte instructions
С	Reference to coprocessor registers
d	Doubleword instructions
f	Reference to floating-point registers
h	Halfword instructions
q	Quadword instructions
sr	Status register

Table 5–3 outlines the correspondence between SPARC hardware integer instructions and SPARC assembly language instructions.

The syntax of individual instructions is designed so that a destination operand (if any), which may be either a register or a reference to a memory location, is always the last operand in a statement.

Note - In Table 5-3,

- Braces ({ }) indicate optional arguments.Braces are not literally coded.
- Brackets ([]) indicate indirection: the contents of the addressed memory location
- are being read from or written to.
 - Brackets are coded literally in the assembly language. Note that the usage of brackets described in Chapter 2 differs from the usage of these brackets.
- All Bicc and Bfcc instructions described may indicate that the annul bit is to be set by appending ", a" to the opcode mnemonic; for example,

[&]quot;bgeu,a label"

TABLE 5-3

Opcode	Mnemonic	Argument List	Operation	Comments
ADD	add	reg _{rs1} , reg_or_imm, reg _{rd}	Add	
ADDcc	addcc	reg _{rs1} , reg_or_imm, reg _{rd}	Add and modify icc	
ADDX	addx	reg _{rs1} , reg_or_imm, reg _{rd}	Add with carry	
ADDXcc	addxcc	reg _{rs1} , reg_or_imm, reg _{rd}		
AND	and	reg _{rs1} , reg_or_imm, reg _{rd}	And	
ANDcc	andcc	reg _{rs1} , reg_or_imm, reg _{rd}		
ANDcc	andn	reg _{rs1} , reg_or_imm, reg _{rd}		
ANDNcc	andcc	reg _{rs1} , reg_or_imm, reg _{rd}		
BN	bn{,a}	label	Branch on integer condition codes	branch never
BNE	bne{,a}	label		synonym: bnz
BE	be{,a}	label		synonym: bz
BG	bg{,a}	label		
BLE	ble{,a}	label		
BGE	bge{,a}	label		
BI	bl{,a}	label		
BGU	bgu{,a}	label		
BLEU	bleu{,a}	label		
BCC	bcc{,a}	label		synonym: bgeu
BCS	bcs{,a}	label		synonym: blu
BPOS	bpos{,a}	label		
BNEG	bneg{,a}	label		
BVC	bvc{,a}	label		
BVS	bvs{,a}	label		
BA	ba{,a}	label		synonym: b
CALL	call	label	Call subprogram	

(Continued) TABLE 5-3

Opcode	Mnemonic	Argument List	Operation	Comments
СВссс	cbn{,a}	label	Branch on coprocessor	branch never
	cb3{,a}	a} label condition codes		
	cb2{,a}	label		
	cb23{,a}	label		
	cb1{,a}	label		
	cb13{,eo}	label		
	cb12{,a}	label		
	cb123{,a}	label		
	cb0{,a}	label		
	cb03{,a}	label		
	cb02{,a}	label		
	cb023{,a}	label		
	cb01{,a}	label		
	cb013{,a}	label		
	cb012{,a}	label		
	cba{,a}	label		
FBN	fbn{,a}	label	Branch on floating-point	branch never
FBU	fbu{,a}	label	condition codes	
FBG	fbg{,a}	label		
FBUG	fbug{,a}	label		
FBL	fbl{,a}	label		
FBUL	fbul{,a}	label		
FBLG	fblg{,a}	label		
FBNE	fbne{,a}	label		synonym: fbnz
FBE	fbe{,a}	label		synonym: fbz

TABLE 5–3(Continued)

Opcode	Mnemonic	Argument List	Operation	Comments
FBUE	fbue{,a}	label		
FBGE	fbge{,a}	label		
FBUGE	fbuge{,a}	label		
FBLE	fble{,a}	label		
FBULE	fbule{,a}	label		
FBO	fbo{,a}	label		
FBA	fba{,a}	label		
FLUSH	flush	address	Instruction cache flush	
JMPL	jmpl	address, reg _{rd}	Jump and link	
LDSB	ldsb	[address], reg _{rd}	Load signed byte	
LDSH	ldsh	[address], reg _{rd}	Load signed halfword	
LDSTUB	ldstub	[address], reg _{rd}	Load-store unsigned byte	
LDUB	ldub	[address], reg _{rd}	Load unsigned byte	
LDUH	lduh	[address], reg _{rd}	Load unsigned halfword	
LD	ld	[address], reg _{rd}	Load word	
LDD	ldd	[address], reg _{rd}	Load double word	reg _{rd} must be even
LDF	ld	[address], freg _{rd}		
LDFSR	ld	[address], %fsr	Load floating-point register	
LDDF	ldd	[address], freg _{rd}	Load double floating-point	freg _{rd} must be even
LDC	ld	[address], creg _{rd}	Load coprocessor	
LDCSR	ld	[address], %csr	Load double coprocessor	
LDDC	ldd	[address], creg _{rd}		
LDSBA	ldsba	[regaddr]asi, reg _{rd}	Load signed byte from	
LDSHA	ldsha	[regaddr]asi, reg _{rd}	alternate space	
LDUBA	lduba	[regaddr]asi, reg _{rd}		
LDUHA	lduha	[regaddr]asi, reg _{rd}		
LDA	lda	[regaddr]asi, reg _{rd}		
LDDA	ldda	[regaddr]asi, reg _{rd}		reg _{rd} must be even
LDSTUBA	ldstuba	[regaddr]asi, reg _{rd}		

(Continued) TABLE 5-3

Opcode	Mnemonic	Argument List	Operation	Comments
MULScc	mulscc	reg _{rs1} , reg_or_imm, reg _{rd}	Multiply step (and modify icc)	
NOP	nop		No operation	
OR	or	reg _{rs1} , reg_or_imm, reg _{rd}	Inclusive or	
ORcc	orcc	reg _{rs1} , reg_or_imm, reg _{rd}		
ORN	orn	reg _{rs1} , reg_or_imm, reg _{rd}		
ORNcc	orncc	reg _{rs1} , reg_or_imm, reg _{rd}		
RDASR	rd	%asrn _{rs1} , reg _{rd}		
RDY	rd	%y, reg _{rd}		See synthetic instructions.
RDPSR	rd	%psr, reg _{rd}		See synthetic instructions.
RDWIM	rd	%wim, reg _{rd}		See synthetic instructions.
RDTBR	rd	%tbr, reg _{rd}	%tbr, reg _{rd}	
RESTORE	restore	reg _{rs1} , reg_or_imm, reg _{rd}	reg _{rs1} , reg_or_imm, reg rd	
RETT	rett	address	Return from trap	
SAVE	save	reg _{rs1} , reg_or_imm, reg _{rd}	reg_or_imm, reg _{rd}	
SDIV	sdiv	reg _{rs1} , reg_or_imm, reg _{rd}	Signed divide	
SDIVcc	sdivcc	reg _{rs1} , reg_or_imm, reg _{rd}	Signed divide and modify icc	
SMUL	smul	reg _{rs1} , reg_or_imm, reg _{rd}	Signed multiply	
SMULcc	smulcc	reg _{rs1} , reg_or_imm, reg _{rd}		
SETHI	sethi	const22, reg _{rd}	Set high 22 bits of register	
	sethi	%hi(value), reg _{rd}		See synthetic instructions.
SLL	sll	reg _{rs1} , reg_or_imm, reg _{rd}	Shift left logical	
SRL	srl	reg _{rs1} , reg_or_imm, reg _{rd}	Shift right logical	
SRA	sra	reg _{rs1} , reg_or_imm, reg _{rd}	Shift right arithmetic	

(Continued) TABLE 5-3

Opcode	Mnemonic	Argument List	Operation	Comments
STB	stb	reg _{rd} , [address]	reg _{rd} , [address] Store byte	
STH	sth	reg _{rd} , [address]	reg _{rd} , [address] Store half-word	
ST	st	reg _{rd} , [address]		
STD	std	reg _{rd} , [address]		reg _{rd} Must be even
STF	st	freg _{rd} , [address]		
STDF	std	freg _{rd} , [address]		
STFSR	st	%fsr,[address]	Store floating-point status register	freg _{rd} Must be even
STDFQ	std	%fq,[address]	Store double floating-point queue	
STC	st	creg _{rd} , [address]	Store coprocessor	creg _{rd} Must be even
STDC	std	creg _{rd} , [address]		creg _{rd} Must be even
STCSR	st	%csr,[address]		
STDCQ	std	%cq, [address]	Store double coprocessor	
STBA	stba	reg _{rd} [regaddr]asi	Store byte into alternate space	Synonyms: stuba, stsba
STHA	stha	reg _{rd} [regaddr]asi		Synonyms: stuha, stsha
STA	sta	reg _{rd} , [regaddr]asi		
STDA	stda	reg _{rd} , [regaddr]asi		reg _{rd} Must be even
SUB	sub	reg _{rs1} , reg_or_imm, reg _{rd}	Subtract	
SUBcc	subcc	reg _{rs1} , reg_or_imm, reg _{rd}	Subtract and modify icc	
SUBX	subx	reg _{rs1} , reg_or_imm, reg _{rd}	reg _{rs1} , reg_or_imm, reg _{rd} Subtract with carry	
SUBXcc	subxcc	reg _{rs1} , reg_or_imm, reg _{rd}		
SWAP	swap	[address], reg _{rd}	Swap memory word	
SWAPA	swapa	[regaddr]asi, reg _{rd}	with register	
Ticc	tn	software_trap_number	r Trap on integer condition code	Trap never

(Continued) TABLE 5-3

Opcode	Mnemonic	Argument List	Operation	Comments
	tne	software_trap_number	Note: Trap numbers 16-31 are reserved for the user. Currently-defined trap numbers are those defined in /usr/include/sys/trap.h	Synonym: tnz
	te	software_trap_number		Synonym: tz
	tg	software_trap_number		
	tle	software_trap_number		
	tge	software_trap_number		
	tl	software_trap_number		
	tgu	software_trap_number		
	tleu	software_trap_number		Synonym: tcc
	tlu	software_trap_number		Synonym: tcc
	tgeu	software_trap_number		
	tpos	software_trap_number		
	tneg	software_trap_number		
	tvc	software_trap_number		Synonym: t
	tvs	software_trap_number		
	ta	software_trap_number		
TADDcc	taddcc	reg _{rs1} , reg_or_imm, reg _{rd}	Tagged add and modify icc	
TSUBcc	tsubcc	reg _{rs1} , reg_or_imm, reg _{rd}		
TADDccTV	taddcctv	reg _{rs1} , reg_or_imm, reg _{rd}	Tagged add and modify icc	
TSUBccTV	tsubcctv	reg _{rs1} , reg_or_imm, reg _{rd}	and trap on overflow	
UDIV	udiv	reg _{rs1} , reg_or_imm, reg _{rd}	Unsigned divide	
UDIVcc	udivcc	reg _{rs1} , reg_or_imm, reg _{rd}	Unsigned divide and modify	
UMUL	umul	reg _{rs1} , reg_or_imm, reg _{rd}	Unsigned multiply	
UMULcc	umulcc	reg _{rs1} , reg_or_imm, reg _{rd}	Unsigned multiply and modify icc	
UNIMP	unimp	const22	Illegal instruction	
WRASR	wr	reg_or_imm, %asrn _{rs1}		

TABLE 5–3(Continued)

Opcode	Mnemonic	Argument List	Operation	Comments
WRY	wr	reg _{rs1} , reg_or_imm, %y		See synthetic instructions
WRPSR	wr	reg _{rs1} , reg_or_imm, %psr		See synthetic instructions
WRWIM	wr	reg _{rs1} , reg_or_imm, %wim		See synthetic instructions
WRTBR	wr	reg _{rs1} , reg_or_imm, %tbr		See synthetic instructions
XNOR	xnor	reg _{rs1} , reg_or_imm, reg _{rd}	Exclusive nor	
XNORcc	xnorcc	reg _{rs1} , reg_or_imm, reg _{rd}		
XOR	xor	reg _{rs1} , reg_or_imm, reg _{rd}	Exclusive or	
XORcc	xorcc	reg _{rs1} , reg_or_imm, reg _{rd}		

5.3 Floating-Point Instruction

Table 5–4 shows floating-point instructions. In cases where more than numeric type is involved, each instruction in a group is described; otherwise, only the first member of a group is described.

TABLE 5-4

SPARC	Mnemonic ¹	Argument List	Description
FiTOs	fitos	freg _{rs2} , freg _{rd}	Convert integer to single
FiTOd	fitod	freg _{rs2} , freg _{rd}	Convert integer to double
FiTOq	fitoq	freg _{rs2} , freg _{rd}	Convert integer to quad
FsTOi	fstoi	freg _{rs2} , freg _{rd}	Convert single to integer
FdTOi	fdtoi	freg _{rs2} , freg _{rd}	Convert double to integer

Types of Operands are denoted by the following lower-case letters:

- i integer
- s single
- d double
- $\mathtt{q} \ \text{quad}$

(Continued) TABLE 5-4

SPARC	Mnemonic ¹	Argument List	Description
FqTOi	fqtoi	freg _{rs2} , freg _{rd}	Convert quad to integer
FsTOd	fstod	freg _{rs2} , freg _{rd}	Convert single to double
FsTOq	fstoq	freg _{rs2} , freg _{rd}	Convert single to quad
FdTOs	fdtos	freg _{rs2} , freg _{rd}	Convert double to single
FdTOq	fdtoq	freg _{rs2} , freg _{rd}	Convert double to quad
FqTOd	fqtod	freg _{rs2} , freg _{rd}	Convert quad to double
FqTOs	fqtos	freg _{rs2} , freg _{rd}	Convert quad to single
FMOVs	fmovs	freg _{rs2} , freg _{rd}	Move
FNEGs	fnegs	freg _{rs2} , freg _{rd}	Negate
FABSs	fabss	freg _{rs2} , freg _{rd}	Absolute value
FSQRTs	fsqrts	freg _{rs2} , freg _{rd}	Square root
FSQRTd	fsqrtd	freg _{rs2} , freg _{rd}	
FSQRTq	fsqrtq	freg _{rs2} , freg _{rd}	
FADDs	fadds	freg _{rs1} , freg _{rs2} , freg _{rd}	Add
FADDd	faddd	freg _{rs1} , freg _{rs2} , freg _{rd}	
FADDq	faddq	freg _{rs1} , freg _{rs2} , freg _{rd}	
FSUBs	fsubs	freg _{rs1} , freg _{rs2} , freg _{rd}	Subtract
FSUBd	fsubd	freg _{rs1} , freg _{rs2} , freg _{rd}	
FSUBq	fsubq	freg _{rs1} , freg _{rs2} , freg _{rd}	
FMULs	fmuls	freg _{rs1} , freg _{rs2} , freg _{rd}	Multiply
FMULd	fmuld	freg _{rs1} , freg _{rs2} , freg _{rd}	
FMULq	fmulq	freg _{rs1} , freg _{rs2} , freg _{rd}	
FdMULq	fmulq	freg _{rs1} , freg _{rs2} , freg _{rd}	Multiply double to quad
FsMULd	fsmuld	freg _{rs1} , freg _{rs2} , freg _{rd}	Multiply single to double

Types of Operands are denoted by the following lower-case letters:

i integer

s single

d double

 $^{{\}tt q} \, \textbf{quad} \,$

TABLE 5-4	(Continued)
IADLE 3-4	(Communea)

	1	T	
SPARC	Mnemonic ¹	Argument List	Description
FDIVs	fdivs	freg _{rs1} , freg _{rs2} , freg _{rd}	Divide
FDIVd	fdivd	$freg_{rs1}$, $freg_{rs2}$, $freg_{rd}$	
FDIVq	fdivq	$freg_{rs1}$, $freg_{rs2}$, $freg_{rd}$	
FCMPs	fcmps	freg _{rs1} , freg _{rs2}	Compare
FCMPd	fcmpd	$freg_{rs1}, freg_{rs2}$	
FCMPq	fcmpq	$freg_{rs1}, freg_{rs2}$	
FCMPEs	fcmpes	$freg_{rs1}, freg_{rs2}$	Compare, generate exception if
FCMPEd	fcmped	$freg_{rs1}, freg_{rs2}$	not ordered
FCMPEq	fcmpeq	$freg_{rs1}$, $freg_{rs2}$	

Types of Operands are denoted by the following lower-case letters:

5.4 Coprocessor Instructions

All *coprocessor-operate* (cpop*n*) instructions take all operands from and return all results to coprocessor registers. The data types supported by the coprocessor are coprocessor-dependent. Operand alignment is also coprocessor-dependent. Coprocessor-operate instructions are described in Table 5–5.

If the EC (PSR_enable_coprocessor) field of the processor state register (PSR) is 0, or if a coprocessor is not present, a cpop*n* instruction causes a *cp_disabled* trap.

The conditions that cause a *cp_exception* trap are coprocessor-dependent.

TABLE 5-5

SPARC	Mnemonic	Argument List	Name	Comments
CPop1	cpop1	opc, reg _{rs1} , reg _{rs2} , reg _{rd}	Coprocessor operation	
CPop2	cpop2	opc, reg _{rs1} , reg _{rs2} , reg _{rd}	Coprocessor operation	May modify ccc

i integer

s single

d double

q quad

5.5 Synthetic Instructions

Table 5–6 describes the mapping of synthetic instructions to hardware instructions.

TABLE 5-6

Synthetic Instruction	on	Hardware Equivaler	nt(s)	Comment
btst	reg_or_imm, reg _{rs1}	andcc	reg _{rs1} , reg_or_imm, %g0	Bit test
bset	reg_or_imm, reg _{rd}	or	reg _{rd} , reg_or_imm, reg _{rd}	Bit set
bclr	reg_or_imm, reg _{rd}	andn	reg _{rd} , reg_or_imm, reg _{rd}	Bit clear
btog	reg_or_imm, reg _{rd}	xor	reg _{rd} , reg_or_imm, reg _{rd}	Bit toggle
call	reg_or_imm	jmpl	reg_or_imm, %07	
clr	reg _{rd}	or	%g0, %g0, reg _{rd}	Clear (zero) register
clrb	[address]	stb	%g0, [address]	Clear byte
clrh	[address]	st	%g0, [address]	Clear halfword
clr	[address]	st	%g0, [address]	Clear word
cmp	reg, reg_or_imm	subcc	reg _{rs1} , reg_or_imm, %g0	Compare
dec	reg _{rd}	sub	reg _{rd} , 1, reg _{rd}	Decrement by 1
dec	const13, reg _{rd}	sub	reg _{rd} , const13, reg _{rd}	Decrement by const13
deccc	reg_{rd}	subcc	reg _{rd} , 1, reg _{rd}	Decrement by 1 and set
deccc	const13, reg _{rd}	subcc	reg _{rd} , const13, reg _{rd}	Decrement by const13 and set icc
inc	reg _{rd}	add	reg _{rd} , 1, reg _{rd}	Increment by 1
inc	const13, reg _{rd}	add	reg _{rd} , const13, reg _{rd}	Increment by const13
inccc	$reg_{\rm rd}$	addcc	reg _{rd} , 1, reg _{rd}	Increment by 1 and set
inccc	const13, reg _{rd}	addcc	reg _{rd} , const13, reg _{rd}	Increment by const13 and set icc
jmp	address	jmpl	address, %g0	

(Continued) TABLE 5-6

Synthetic Instru	uction	Hardware Equiva	lent(s)	Comment
mov	reg_or_imm,reg _{rd}	or	%g0, reg_or_imm, reg _{rd}	
mov	%y, reg _{rs1}	rd	%y, reg _{rs1}	
mov	%psr, reg _{rs1}	rd	%psr, reg _{rs1}	
mov	%wim, $reg_{ m rs1}$	rd	%wim, $reg_{ m rs1}$	
mov	%tbr, reg_{rs1}	rd	%tbr, reg _{rs1}	
mov	reg_or_imm, %y	wr	%g0,reg_or_imm,%y	
mov	reg_or_imm, %psr	wr	%g0,reg_or_imm,%psr	
mov	reg_or_imm, %wim	wr	%g0,reg_or_imm,%wim	
mov	reg_or_imm,%tbr	wr	%g0,reg_or_imm,%tbr	
not	reg _{rs1} , reg _{rd}	xnor	reg _{rs1} , %g0, reg _{rd}	One's complement
not	reg _{rd}	xnor	reg _{rd} , %g0, reg _{rd}	One's complement
neg	reg _{rs1} , reg _{rd}	sub	%g0, reg _{rs2} , reg _{rd}	Two's complement
neg	reg _{rd}	sub	%g0, reg _{rd} , reg _{rd}	Two's complement
restore		restore	%g0, %g0, %g0	Trivial restore
save		save	%g0, %g0, %g0	Trivial save
				trivial save should only be used in supervisor code!
set	value,reg _{rd}	or	%g0, value, reg _{rd}	if -4096 ≤value ≤ 4095
				Do not use the set synthetic instruction in an instruction delay slot.
set	value,reg _{rd}	sethi	%hi(value), reg _{rd}	if ((value & 0x3ff) == 0)
set	value, reg _{rd}	sethi	%hi(value), reg _{rd} ; reg _{rd} ,	otherwise
		or	%lo(value), reg _{rd}	Do not use the set synthetic instruction in an instruction delay slot.
skipz		bnz,a .+8		if z is set, ignores next instruction
skipnz		bz,a .+8		if z is not set, ignores next instruction
tst	reg	orcc	<i>reg</i> _{rs1} , %g0, %g0	test

5.6 V8/V9 Natural Pseudo Instructions

Table 5–7 describes the V8/V9 natural pseudo instructions that will help increase the portability of your assembly code from V8/V8plus to V9.

TABLE 5-7

	-xarch=		
Pseudo Instructions	V8/V8plus ¹	V9	
ldn	ld	ldx	
stn	st	stx	
ldna	lda	ldxa	
stna	sta	stxa	
setn	set	setx	
setnhi	sethi	setxhi	
casn	cas	casx	
slln	sll	sllx	
srln	srl	srlx	
sran	sra	srax	
clrn	clr	clrx	

Indicates default setting

Note – Depending on the value set for the -xarch option, the assembler substitutes the appropriate pseudo instruction.

Pseudo-Operations

The pseudo-operations listed in this appendix are supported by the SPARC assembler.

A.1 Alphabetized Listing with Descriptions

```
.alias
Turns off the effect of the preceding .noalias pseudo-op. (Compiler-generated
only.)
```

.align boundary
Aligns the location counter on a boundary where (("location counter" mod boundary) == 0); boundary may be any power of 2.

.ascii string [, string"]
Generates the given sequence(s) of ASCII characters.

.asciz string [, string] *
 Generates the given sequence(s) of ASCII characters. This pseudo-op appends a
 null (zero) byte to each string.

.byte *8bitval* [, *8bitval*] *
Generates (a sequence of) initialized bytes in the current segment.

.common *symbol*, *size* [, *sect_name*] [, *alignment*] Provides a tentative definition of *symbol*. *Size* bytes are allocated for the object represented by *symbol*.

- If the symbol is not defined in the input file and is declared to be *local* to the file, the symbol is allocated in *sect_name* and its location is optionally aligned to a multiple of *alignment*. If *sect_name* is not given, the symbol is allocated in the uninitialized data section (*bss*). Currently, only .bss is supported for the section name. (.data is not currently supported.)
- If the symbol is not defined in the input file and is declared to be *global*, the SPARC link editor allocates storage for the symbol, depending on the definition of *symbol_name* in other files. Global is the default binding for common symbols.
- If the symbol is defined in the input file, the definition specifies the location of the symbol and the tentative definition is overridden.

.double Orfloatval [, Orfloatval] *

Generates (a sequence of) initialized double-precision floating-point values in the current segment. *floatval* is a string acceptable to atof(3); that is, an optional sign followed by a non-empty string of digits with optional decimal point and optional exponent.

.empty

Suppresses assembler complaints about the next instruction presence in a delay slot when used in the delay slot of a Control Transfer Instruction (CTI).

Some instructions should not be in the delay slot of a CTI. See the *SPARC Architecture Manual* for details.

.file string

Creates a symbol table entry where *string* is the symbol name and STT_FILE is the symbol table type. *string* specifies the name of the source file associated with the object file.

.global symbol [, symbol] * .globl symbol [, symbol] *

Declares each *symbol* in the list to be global; that is, each symbol is either defined externally or defined in the input file and accessible in other files; default bindings for the symbol are overridden.

- A global symbol definition in one file will satisfy an undefined reference to the same global symbol in another file.
- Multiple definitions of a defined global symbol is not allowed. If a defined global symbol has more than one definition, an error will occur.
- A global psuedo-op oes not need to occur before a definition, or tentative definition, of the specified symbol.

Note – This pseudo-op by itself does not define the symbol.

.half 16bitval [, 16bitval] *

Generates (a sequence of) initialized halfwords in the current segment. The location counter must already be aligned on a halfword boundary (use .align 2).

.ident string

Generates the null terminated string in a comment section. This operation is equivalent to:

.pushsection .comment

.asciz string

.popsection

.local symbol [, symbol] *

Declares each *symbol* in the list to be local; that is, each symbol is defined in the input file and not accessible in other files; default bindings for the symbol are overridden. These symbols take precedence over *weak* and *global* symbols.

Since local symbols are not accessible to other files, local symbols of the same name may exist in multiple files.

Note – This pseudo-op by itself does not define the symbol.

.noalias %reg1, %reg2

%reg1 and %reg2 will not alias each other (that is, point to the same destination)
until a .alias pseudo-op is issued. (Compiler-generated only.)

.nonvolatile

Defines the end of a block of instruction. The instructions in the block may not be permuted. This pseudo-op has no effect if:

- The block of instruction has been previously terminated by a Control Transfer Instruction (CTI) or a label
- There is no preceding .volatile pseudo-op

.nword 64bitval [, 64bitval] *

If -xarch=v8/v8plus then assembler interprets the instruction as .word. If -xarch=v9 the assembler interprets the instruction as .xword.

.optim string

This pseudo-op changes the optimization level of a particular function. (Compiler-generated only.)

.popsection

Removes the top section from the section stack. The new section on the top of the stack becomes the current section. This pseudo-op and its corresponding .pushsection command allow you to switch back and forth between the named sections.

.proc n

Signals the beginning of a *procedure* (that is, a unit of optimization) to the peephole optimizer in the SPARC assembler; *n* specifies which registers will contain the return value upon return from the procedure. (Compiler-generated only.)

.pushsection sect_name [, attributes]

Moves the named section to the top of the section stack. This new top section then becomes the current section. This pseudo-op and its corresponding <code>.popsection</code> command allow you to switch back and forth between the named sections.

.quad Orfloatval [, Orfloatval] *

Generates (a sequence of) initialized quad-precision floating-point values in the current segment. *floatval* is a string acceptable to atof(3); that is, an optional sign followed by a non-empty string of digits with optional decimal point and optional exponent.

 $\mbox{\bf Note}$ – The . quad command currently generates quad-precision values with only double-precision significance.

```
.reserve symbol, size [, sect_name [, alignment]]
```

Defines *symbol*, and reserves *size* bytes of space for it in the *sect_name*. This operation is equivalent to:

```
.pushsection sect_name
  .align alignment
symbol:
  .skip size
  .popsection
```

If a section is not specified, space is reserved in the current segment.

```
.section section_name [, attributes] Makes the specified section the current section.
```

The assembler maintains a section stack which is manipulated by the section control directives. The current section is the section that is currently on top of the stack. This pseudo-op changes the top of the section stack.

- If section_name does not exist, a new section with the specified name and attributes is created.
- If section_name is a non-reserved section, attributes must be included the first time it is specified by the .section directive.

See the sections "3.2.2 Predefined User Sections" on page 29 and "3.2.3 Predefined Non-User Sections" on page 31 in Chapter 3 for a detailed description of the reserved sections. See Table 3–2 in Chapter 3 for a detailed description of the section attribute flags.

Attributes can be:

```
#write | #alloc | #execinstr
```

.seq section_name

Note – This pseudo-op is currently supported for compatibility with existing SunOS 4.1 SPARC assembly language programs. This pseudo-op has been replaced by the . section pseudo-op.

Changes the current section to one of the predefined user sections. The assembler will interpret the following SunOS 4.1 SPARC assembly directive: to be the same as the following *SunOS 5.x* SPARC assembly directive:

```
.seg text, .seg data, .seg data1, .seg bss,
.section .text, .section .data, .section .data1,
.section .bss.
```

Predefined user section names are changed in SunOS 5.x.

.single Orfloatval [, Orfloatval] *
Generates (a sequence of) initialized single-precision floating-point values in the
current segment.

Note - This operation does not align automatically.

.size symbol, expr

Declares the symbol size to be *expr. expr* must be an absolute expression.

.skip n

Increments the location counter by n, which allocates n bytes of empty space in the current segment.

.stabn <various parameters>

The pseudo-op is used by *Solaris 2.x* SPARCompilers only to pass debugging information to the symbolic debuggers.

.stabs <various parameters>

The pseudo-op is used by *Solaris 2.x* SPARCompilers only to pass debugging information to the symbolic debuggers.

.type symbol, type

Declares the type of symbol, where *type* can be:

#object

#function

#no_type

See Table 3–6 in Chapter 3, for detailed information on symbols.

.uahalf 16bitval [, 16bitval] *

Generates a (sequence of) 16-bit value(s).

Note – This operation does not align automatically.

.uaword 32bitval [, 32bitval] *

Generates a (sequence of) 32-bit value(s).

Note – This operation does not align automatically.

.version string

Identifies the minimum assembler version necessary to assemble the input file. You can use this pseudo-op to ensure assembler-compiler compatibility. If *string* indicates a newer version of the assembler than this version of the assembler, a fatal error message is displayed and the SPARC assembler exits.

.volatile

Defines the beginning of a block of instruction. The instructions in the section may not be changed. The block of instruction should end at a .nonvolatile pseudo-op and should not contain any Control Transfer Instructions (CTI) or labels. The volatile block of instructions is terminated after the last instruction preceding a CTI or label.

.weak symbol [, symbol]

Declares each *symbol* in the list to be defined either externally, or in the input file and accessible to other files; default bindings of the symbol are overridden by this directive.

Note the following:

- A weak symbol definition in one file will satisfy an undefined reference to a global symbol of the same name in another file.
- Unresolved weak symbols have a default value of zero; the link editor does not resolve these symbols.
- If a *weak* symbol has the same name as a defined *global* symbol, the weak symbol is ignored and no error results.

Note – This pseudo-op does not itself define the symbol.

.word 32bitval [, 32bitval] *

Generates (a sequence of) initialized words in the current segment.

Note – This operation does not align automatically.

.xword 64bitval [, 64bitval] *

Generates (a sequence of) initialized 64-bit values in the current segment.

Note – This operation does not align automatically.

.xstabs <various parameters>

The pseudo-op is used by *Solaris 2.x* SPARCompilers only to pass debugging information to the symbolic debuggers.

symbol =*expr*

Assigns the value of *expr* to *symbol*.

Examples of Pseudo-Operations

This appendix shows some examples of ways to use various pseudo-ops.

B.1 Example 1

This example shows how to use the following pseudo-ops to specify the bindings of variables in C:

```
common, .global, .local, .weak
```

The following C definitions/declarations:

```
int foo1 = 1;
#pragma weak foo2 = foo1
static int foo3;
static int foo4 = 2;
```

can be translated into the following assembly code:

EXAMPLE B-1

```
.pushsection ".data"

.global foo1 ! int foo1 = 1
    .align 4

foo1:
    .word 0x1
    .type foo1,#object ! foo1 is of type data object,
    .size foo1,4 ! with size = 4 bytes

.weak foo2 ! #pragma weak foo2 = foo1
foo2 = foo1

.local foo3 ! static int foo3
```

EXAMPLE B-1 (Continued)

```
.common foo3,4,4

.align 4 ! static int foo4 = 2
foo4:
.word 0x2
.type foo4,#object
.size foo4,4

.popsection
```

B.2 Example 2

This example shows how to use the pseudo-op .ident to generate a string in the .comment section of the object file for identification purposes.

```
.ident "acomp: (CDS) SPARCompilers 2.0 alpha4 12 Aug 1991"
```

B.3 Example 3

The pseudo-ops shown in this example are .align, .global, .type, and .size.

The following C subroutine:

```
int sum(a, b)
    int a, b;
{
    return(a + b);
}
```

can be translated into the following assembly code:

```
! delay slot of retl

.type sum,#function ! sum is of type function
.size sum,.-sum ! size of sum is the diff

! of current location

! counter and the initial

! definition of sum
```

B.4 Example 4

The pseudo-ops shown in this example are .section, .ascii, and .align. The example calls the printf function to output the string "hello world".

```
.section
                        ".data1"
    .aliqn
.L16:
    .ascii
                  "hello world\n\0"
    .section
                       ".text"
    .global
                      main
main:
   save
               %sp,-96,%sp
              .L16,%o0
    set
               printf,1
    call
    nop
    restore
```

B.5 Example 5

This example shows how to use the .volatile and .nonvolatile pseudo-ops to protect a section of handwritten assembly code from peephole optimization.

Using the Assembler Command Line

This appendix is organized into the following secitons:

- "C.1 Assembler Command Line" on page 67
- "C.2 Assembler Command Line Options" on page 68
- "C.3 Disassembling Object Code" on page 71

C.1 Assembler Command Line

You invoke the assembler command line as follows:

```
as [options] [inputfile] ...
```

Note – The language drivers (such as cc and f77) invoke the assembler command line with the fbe command. You can use either the as or fbe command to invoke the assembler command line.

The as command translates the assembly language source files, *inputfile*, into an executable object file, *objfile*. The SPARC assembler recognizes the filename argument *hyphen* (-) as the standard input. It accepts more than one file name on the command line. The input file is the concatenation of all the specified files. If an invalid option is given or the command line contains a syntax error, the SPARC assembler prints the error (including a synopsis of the command line syntax and options) to standard error output, and then terminates.

The SPARC assembler supports macros, #include files, and symbolic substitution through use of the C preprocessor cpp. The assembler invokes the preprocessor before assembly begins if it has been specified from the command line as an option. (See the -P option.)

C.2 Assembler Command Line Options

-h

This option generates extra symbol table information for the source code browser.

- If the as command line option -P is set, the cpp preprocessor also collects browser information.
- If the as command line option -m is set, this option is ignored as the m4 macro processor does not generate browser data.

For more information about the SPARCworks SourceBrowser, see the *Browsing Source Code* manual.

-Dname -Dname=def

When the -P option is in effect, these options are passed to the cpp preprocessor without interpretation by the as command; otherwise, they are ignored.

-Ipath

When the -P option is in effect, this option is passed to the cpp preprocessor without interpretation by the as command; otherwise, it is ignored.

-K PIC

This option generates position-independent code. This option has the same functionality as the -k option under the SunOS 4.1 SPARC assembler.

Note - - K PIC and - K pic are equivalent.

- L

Saves all symbols, including temporary labels that are normally discarded to save space, in the ELF symbol table.

– m

This option runs m4 macro preprocessing on input. The m4 preprocessor is more powerful than the C preprocessor (invoked by the -P option), so it is more useful for complex preprocessing. See the m4 (1) man page for more information about the m4 macro-processor.

-n

Suppress all warnings while assembling.

-o outfile

Takes the next argument as the name of the output file to be produced. By default, the .s suffix, if present, is removed from the input file and the .o suffix is appended to form the ouput file name.

- P

Run *cpp*, the C preprocessor, on the files being assembled. The preprocessor is run separately on each input file, not on their concatenation. The preprocessor output is passed to the assembler.

$-Q\{y|n\}$

This option produces the "assembler version" information in the comment section of the output object file if the y option is specified; if the n option is specified, the information is suppressed.

- 0

This option causes the assembler to perform a quick assembly. Many error-checks are not performed when -q is specified.

Note – This option disables many error checks. It is recommended that you do *not* use this option to assemble handwritten assembly language.

-S[a|b|c|1|A|B|C|L]

Produces a disassembly of the emitted code to the standard output. Adding each of the following characters to the -S option produces:

- a disassembling with address
- b disassembling with ".bof"
- c disassembling with comments
- 1 disassembling with line numbers

Capital letters turn the switch off for the corresponding option.

- s

This option places all stabs in the ".stabs" section. By default, stabs are placed in "stabs.excl" sections, which are stripped out by the static linker 1d during final execution. When the -s option is used, stabs remain in the final executable because ".stab" sections are not stripped out by the static linker 1d.

- T

This is a migration option for SunOS 4.1 assembly files to be assembled on SunOS 5.x systems. With this option, the symbol names in SunOS 4.1 assembly files will be interpreted as SunOS 5.x symbol names. This option can be used in conjunction with the -S option to convert SunOS 4.1 assembly files to their corresponding SunOS 5.x versions.

-Uname

When the -P option is in effect, this option is passed to the cpp preprocessor without interpretation by the as command; otherwise, it is ignored.

- V

This option writes the version information on the standard error output.

-xarch=v7

This option instructs the assembler to accept instructions defined in the SPARC version 7 (V7) architecture. The resulting object code is in ELF format.

-xarch=v8

This option instructs the assembler to accept instructions defined in the SPARC-V8 architecture. The resulting object code is in ELF format. The quad-precision floating-point instructions are allowed; however when the program is executed these instructions cause a hardware exception called "trap" (an illegal instruction trap). The kernel has the trap handler to emulate the quad precision floating-point arithmetic. Consequently, all quad precision arithmetic is performed by the emulator in the kernel.

-xarch=v8a

This option instructs the assembler to accept instructions defined in the SPARC-V8 architecture, less the fsmuld instruction. The resulting object code is in ELF format. The quad-precision floating-point instructions are allowed; however when the program is executed these instructions cause a hardware exception called "trap" (an illegal instruction trap). The kernel has the trap handler to emulate the quad precision floating-point arithmetic. Consequently, all quad precision arithmetic is performed by the emulator in the kernel. This is the default choice of the <code>-xarch=</code> options.

-xarch=v8plus

This option instructs the assembler to accept instructions defined in the SPARC-V9 architecture. The resulting object code is in ELF format. The quad-precision floating-point instructions are allowed; however when the program is executed these instructions cause a hardware exception called "trap" (an illegal instruction trap). The kernel has the trap handler to emulate the quad precision floating-point arithmetic. Consequently, all quad precision arithmetic is performed by the emulator in the kernel. It will not execute on a Solaris V8 system (a machine with a V8 processor). It will execute on a Solaris V8+ system. This combination is a SPARC 64-bit processor and a 32-bit OS. For more information regarding SPARC-V9 instructions, see Appendix E.

-xarch=v8plusa

This option instructs the assembler to accept instructions defined in the SPARC-V9 architecture, plus the instructions in the Visual Instruction Set (VIS). The resulting object code is in V8+ ELF format. It will not execute on a Solaris V8 system. It will execute on a Solaris V8+ system. For more information about VIS instructions, see the *UltraSPARC Programmer's Reference Manual* and the *UltraSPARC User's Guide*.

The quad-precision floating-point instructions are allowed; however when the program is executed these instructions cause a hardware exception called "trap" (an illegal instruction trap). The kernel has the trap handler to emulate the quad precision floating-point arithmetic. Consequently, all quad precision arithmetic is performed by the emulator in the kernel.

-xarch=v9

This option limits instruction set to the SPARC-V9 architecture. The resulting .o object files are in 64-bit ELF format and can only be linked with other object files in the same format. The resulting executable can only be run on a 64-bit SPARC processor running 64-bit Solaris 7 with the 64-bit kernel.

Note – This option is available only on Solaris 7.

-xarch=v9a

This option limits instruction set to the SPARC-V9 architecture, adding the Visual Instruction Set (VIS) and extensions specific to UltraSPARC processors. The resulting .o object files are in 64-bit ELF format and can only be run on a 64-bit SPARC processor running 64-bit Solaris 7 with the 64-bit kernel.

Note – This option is available only on Solaris 7.

C.3 Disassembling Object Code

The dis program is the object code disassembler for ELF. It produces an assembly language listing of the object file. For detailed information about this function, see the man page dis(1).

An Example Language Program

The following code shows an example C language program; the second example code shows the corresponding assembly code generated by SPARCompiler C 3.0.2 that runs on the *Solaris 2.x* operating environment. Comments have been added to the assembly code to show correspondence to the C code.

The following C Program computes the first n Fibonacci numbers:

EXAMPLE D-1 C Program Example Source

```
/* a simple program computing the first n Fibonacci numbers */
extern unsigned * fibonacci();
#define MAX FIB REPRESENTABLE 49
/* compute the first n Fibonacci numbers */
unsigned * fibonacci(n)
    int n;
 static unsigned fib_array[MAX_FIB_REPRESENTABLE] = {0,1};
 unsigned prev_number = 0;
 unsigned curr_number = 1;
 if (n >= MAX_FIB_REPRESENTABLE) {
   printf("Fibonacci(%d) cannot be represented in a 32 bit word\n", n);
   exit(1);
  for (i = 2; i < n; i++) {
   fib_array[i] = prev_number + curr_number;
   prev_number = curr_number;
   curr number = fib array[i];
 return(fib_array);
```

```
main()
{
  int n, i;
  unsigned * result;

  printf("Fibonacci(n):, please enter n:\n");
  scanf("%d", &n);

  result = fibonacci(n);
  for (i = 1; i <= n; i++)
    printf("Fibonacci (%d) is %u\n", i, *result++);
}</pre>
```

The C SPARCompiler generates the following assembler output for the Fibonacci number C source. Annotation has been added to help you understand the code.

EXAMPLE D-2 Assembler Output From C Source

```
a simple program computing the first n Fibonacci numbers,
  showing various pseudo-operations, sparc instructions, synthetic instructions
! pseudo-operations:
                          .align, .ascii, .file, .global, .ident, .proc, .section,
!
              .size, .skip, .type, .word
  sparc instructions: add, bg, bge, bl, ble, ld, or, restore, save, sethi, st
1
  synthetic instructions: call, cmp, inc, mov, ret
    .file
             "fibonacci.c"
                                     ! the original source file name
    .section
               ".text"
                                   ! text section (executable instructions)
           79
    .proc
                              ! subroutine fibonacci, it's return
                       ! value will be in %i0
                                   ! fibonacci() can be referenced
                       ! outside this file
                              ! align the beginning of this section
    .align
                        ! to word boundary
fibonacci:
    save
            %sp,-96,%sp
                                  ! create new stack frame and register
                        ! window for this subroutine
   if (n >= MAX_FIB_REPRESENTABLE) { */
                        ! note, C style comment strings are
                        ! also permitted
           %i0,49
                                ! n >= MAX FIB REPRESENTABLE ?
    cmp
                        ! note, n, the 1st parameter to
                        ! fibonacci(), is stored in %i0 upon
                        ! entry
          .L77003
    mov
           0,%i2
                                ! initialization of variable
```

EXAMPLE D-2 Assembler Output From C Source (Continued)

```
! prev_number is executed in the
                       ! delay slot
/* printf("Fibonacci(%d) cannot be represented in a 32 bits word\n", n); */
           %hi(.L20),%o0 ! if branch not taken, call p

%lo(.L20),%o0 ! set up 1st, 2nd argument in %o0, %o1;
   sethi
                                  ! if branch not taken, call printf(),
   or %00,%lo(.L20),%00
         printf,2
   call
   mov
         %i0,%o1
                               ! registers used as arguments
/* exit(1); */
   call exit,1
   mov
         1,%00
                          ! initialize variables before the loop
.L77003:
/* for (i = 2; i < n; i++) { */
   mov 1,%i4
                              ! curr number = 1
   mov
         2,%i3
                             ! i = 2
   cmp %i3,%i0
                              ! i <= n?
         .L77006
                               ! if not, return
   bge
   sethi %hi(.L16+8),%o0
                                    ! use %i5 to store fib_array[i]
   add %00,%lo(.L16+8),%i5
                         ! loop body
.LY1:
/* fib_array[i] = prev_number + curr_number; */
   add %i2,%i4,%i2
                               ! fib_array[i] = prev_number+curr_number
        %i2,[%i5]
   st.
/* prev_number = curr_number; */
   mov %i4,%i2
                              ! prev_number = curr_number
/* curr_number = fib_array[i]; */
   ld [%i5],%i4
                           ! curr_number = fib_array[i]
   inc
         %i3
                            ! i++
        %i3,%i0
                               ! i <= n?
   cmp
   bl
                           ! if yes, repeat loop
        .LY1
   inc 4,%i5
                             ! increment ptr to fib array[]
.L77006:
/* return(fib_array); */
   sethi %hi(.L16),%o0
                                   ! return fib array in %i0
   add
          %00,%lo(.L16),%i0
   ret
                             ! destroy stack frame and register
   restore
                      ! window
   .tvpe
            fibonacci, #function
                                    ! fibonacci() is of type function
            fibonacci, (.-fibonacci)
                                      ! size of function:
   .size
                      ! current location counter minus
                       ! beginning definition of function
           18
                             ! main program
   .proc
   .global main
   .align
            4
main:
           %sp,-104,%sp
                                 ! create stack frame for main()
/* printf("Fibonacci(n):, please input n:\n"); */
          %hi(.L31),%o0 ! call printf, with 1st arg in %o0
   sethi
   call
           printf,1
   or %00,%lo(.L31),%00
```

```
/* scanf("%d", &n); */
  sethi %hi(.L33),%o0
   call scanf,2
   add %fp,-4,%o1
/* result = fibonacci(n); */
   call fibonacci,1
   ld [%fp-4],%o0
                   ! some initializations before the for-
                   ! loop, put the variables in registers
/* for (i = 1; i <= n; i++) */
   mov 1,%i5
                          ! %i5 <-- i
                           ! %i4 <-- result
   mov
        %o0,%i4
   sethi %hi(.L38),%o0
                              ! %i2 <-- format string for printf
   add %o0,%lo(.L38),%i2
   ld [\$fp-4], \$o0 ! test if (i <= n) ?
   cmp %i5,%o0
                           ! note, n is stored in [%fp-4]
        .LE27
   bg
   nop
.LY2:
                      ! loop body
/* printf("Fibonacci (%d) is %u\n", i, *result++); */
  ld [%i4],%o2
                       ! call printf, with (*result) in %o2,
       %i5,%o1
   mov
                           ! i in %o1, format string in %o0
   call printf,3
        %i2,%o0
   mov
   inc
        %i5
                        ! i++
       [%fp-4],%o0
                          ! i <= n?
   1 d
   cmp %i5,%o0
       .LY2
   ble
   inc 4,%i4
                        ! result++
.LE27:
   ret
   restore
   .type
          main, #function
                        ! type and size of main
        main,(.-main)
   .size
   .section ".data"
                     ! switch to data section
                   ! (contains initialized data)
   .align
.L16:
/* static unsigned fib_array[MAX_FIB_REPRESENTABLE] = {0,1}; */
  .align 4 ! initialization of first 2 elements
   .word
         0
                       ! of fib_array[]
   .align 4
   .word 1
   .skip
         188
          .L16,#object
   .type
                              ! storage allocation for the rest of
                  ! fib_array[]
```

EXAMPLE D-2 Assembler Output From C Source (Continued)

```
.section ".data1"
                            ! the ascii string data are entered
                       ! into the .data1 section;
                       ! #alloc: memory would be allocated
                       ! for this section during run time
                       ! #write: the section contains data
                          that is writeable during process
                            execution
   .aliqn
.L20:
                            ! ascii strings used in the printf stmts
             "Fibonacci(%d) cannot be represented in a 32 bit w"
   .ascii
   .ascii
             "ord\n\0"
   .align
             4
                              ! align the next ascii string to word
                       ! boundary
.L31:
   .ascii
            "Fibonacci(n):, please enter n:\n\0"
   .align
.L33:
   .ascii
             "%d\0"
   .align
.L38:
   .ascii
             "Fibonacci (%d) is %u\n\0"
             "acomp: (CDS) SPARCompilers 2.0 05 Jun 1991"
   .ident
                       ! an idenitfication string produced
                       ! by the compiler to be entered into
                       ! the .comment section
```

APPENDIX **E**

SPARC-V9 Instruction Set

This appendix describes changes made to the SPARC instruction set due to the SPARC-V9 architecture. Application software for the 32-bit SPARC-V8 (Version8) architecture can execute, unchanged, on SPARC-V9 systems.

This appendix is organized into the following sections:

- "E.1 SPARC-V9 Changes" on page 79
- "E.2 SPARC-V9 Instruction Set Changes" on page 81
- "E.3 SPARC-V9 Instruction Set Mapping" on page 84
- "E.4 SPARC-V9 Floating-Point Instruction Set Mapping" on page 93
- "E.5 SPARC-V9 Synthetic Instruction-Set Mapping" on page 95
- "E.6 UltraSPARC and VIS Instruction Set Extensions" on page 97

E.1 SPARC-V9 Changes

The SPARC-V9 architecture differs from SPARC-V8 architecture in the following areas, expanded below: registers, alternate space access, byte order, and instruction set.

E.1.1 Registers

These registers have been deleted:

PSR	Processor State Register
TBR	Trap Base Register

TABLE E-1	(Continued)	١
IABLE E-I	Continuea	J

WIM Window Invalid Mask

These registers have been widened from 32 to 64 bits:

TABLE E-2

Integer registers

All state registers FSR, PC, nPC, and Y

Note – FSR Floating-Point State Register: fcc1, fcc2, and fcc3 (added floating-point condition code) bits are added and the register widened to 64-bits.

These SPARC-V9 registers are within a SPARC-V8 register field:

TABLE E-3

CCR	Condition Codes Register
CWP	Current Window Pointer
PIL	Processor Interrupt Level
TBA	Trap Base Address
TT[MAXTL]	Trap Type
VER	Version

These are registers that have been added.

ASI	Address Space Identifier
CANRESTORE	Restorable Windows
CANSAVE	Savable windows
CLEANWIN	Clean Windows
FPRS	Floating-point Register State
OTHERWIN	Other Windows
PSTATE	Processor State

TABLE E-4(Continued)	
TICK	Hardware clock tick-counter
TL	Trap Level
TNPC[MAXTL]	Trap Next Program Counter
TPC[MAXTL]	Trap Program Counter
TSTATE[MAXTL]	Trap State
WSTATE	Windows State

Also, there are sixteen additional double-precision floating-point registers, f[32] .. f[62]. These registers overlap (and are aliased with) eight additional quad-precision floating-point registers, f[32] .. f[60]

The SPARC-V9, CWP register is decremented during a RESTORE instruction, and incremented during a SAVE instruction. This is the opposite of PSR.CWP's behavior in SPARC-V8. This change has no effect on nonprivileged instructions.

E.1.2 Alternate Space Access

Load- and store-alternate instructions to one-half of the alternate spaces can now be included in user code. In SPARC-V9, loads and stores to ASIs 00_{16} .. $7f_{16}$ are privileged; those to ASIs 80_{16} .. FF_{16} are nonprivileged. In SPARC-V8, access to alternate address spaces is privileged.

E.1.3 Byte Order

SPARC-V9 supports both little- and big-endian byte orders for data accesses only; instruction accesses are always performed using big-endian byte order. In SPARC-V8, all data and instruction accesses are performed in big-endian byte order.

E.2 SPARC-V9 Instruction Set Changes

Application software written for the SPARC-V8 processor runs unchanged on a SPARC-V9 processor.

E.2.1 Extended Instruction Definitions to Support the 64-bit Model

TABLE E-5

FCMP, FCMPE	Floating-Point Compare—can set any of the four floating-point condition codes.
LDFSR, STFSR	Load/Store FSR- only affect low-order 32 bits of FSR
LDUW, LDUWA	Same as LD, LDA in SPARC-V8
RDASR/WRASR	Read/Write State Registers - access additional registers
SAVE/RESTORE	
SETHI	
SRA, SRL, SLL, Shifts	Split into 32-bit and 64-bit versions
Tcc	(was Ticc) Operates with either the 32-bit integer condition codes (icc), or the 64-bit integer condition codes (xcc) $\frac{1}{2}$

All other arithmetic operations operate on 64-bit operands and produce 64-bit results.

E.2.2 Added Instructions to Support 64 bits

F[sdq]TOx	Convert floating point to 64-bit word
FxTO[sdq]	Convert 64-bit word to floating point
FMOV[dq]	Floating-Point Move, double and quad
FNEG[dq]	Floating-point Negate, double and quad
FABS[dq]	Floating-point Absolute Value, double and quad
LDDFA, STDFA, LDFA, STFA	Alternate address space forms of LDDF, STDF, LDF, and STF
LDSW	Load a signed word
LDSWA	Load a signed word from an alternate space
LDX	Load an extended word
LDXA	Load an extended word from an alternate space
LDXFSR	Load all 64 bits of the FSR register

TABLE E-6	(Continued)
STX	Store an extended word
STXA	Store an extended word into an alternate space
STXFSR	Store all 64 bits if the FSR register

E.2.3 Added Instructions to Support High-Performance System Implementation

TABLE E-7

BPcc	Branch on integer condition code with prediction
BPr	Branch on integer register contents with prediction
CASA, CASXA	Compare and Swap from an alternate space
FBPfcc	Branch on floating-point condition code with prediction
FLUSHW	Flush windows
FMOVcc	Move floating-point register if condition code is satisfied
FMOVr	Move floating-point register if integer register satisfies condition
LDQF(A), STQF(A)	Load/Store Quad Floating-point (in an alternate space)
MOVcc	Move integer register if condition code is satisfied
MOVr	Move integer register if register contents satisfy condition
MULX	Generic 64-bit multiply
POPC	Population count
PREFETCH, PREFETCHA	Prefetch Data
SDIVX, UDIVX	Signed and Unsigned 64-bit divide

E.2.4 Deleted Instructions

Coprocessor loads and stores	
RDTBR and WRTBR	TBR no longer exists. It is replaced by TBA, which can be read/written with RDPR/WRPR instructions

TABLE E-8 (Conti	nued)
RDWIM and WRWIM	WIM no longer exists. WIM has been replaced by several register-window registers
REPSR and WRPSR	PSR no longer exists. It has been replaced by several separate registers that are read/written with other instructions
RETT	Return from trap (replace by DONE/RETRY)
STDFQ	Store Double from Floating-point Queue (replaced by the RDPR FQ instruction

E.2.5 Miscellaneous Instruction Changes

TABLE E-9

IMPDEPn	(Changed) Implementation-dependent instructions (replace SPARC-V8 CPop instructions)
MEMBAR	(Added) Memory barrier (memory synchronization support)

E.3 SPARC-V9 Instruction Set Mapping

Opcode	Mnemonic	Argument List	Operation	Comments
вра	ba{,a}	%icc or %xcc, label	(Branch on cc with prediction)	1
	{,pt ,pn}		Branch always	
BPN	bn{,a}	%icc or %xcc, label	Branch never	0
	{,pt ,pn}			
BPNE	bne{,a}	%icc or %xcc, label	Branch on not equal	not Z
	{,pt ,pn}			
BPE	be{,a}	%icc or %xcc, label	Branch on equal	Z
	{,pt ,pn}			
BPG	bg{,a}	%icc or %xcc, label	Branch on greater	not (Z or (N
	{,pt ,pn}			xor V))

(Continued) TABLE E-10

Opcode	Mnemonic	Argument List	Operation	Comments
BPLE	ble{,a}	%icc or %xcc, label	Branch on less or equal	Z or (N xor V)
	{,pt ,pn}			
BPGE	bge{,a}	%icc or %xcc, label	Branch on greater or	not (N xor V)
	{,pt ,pn}		equal	
BPL	bl{,a}	%icc or %xcc, label	Branch on less	N xor V
	{,pt ,pn}			
BPGU	bgu{,a}	%icc or %xcc, label	Branch on greater	not (C or Z)
	{,pt ,pn}		unsigned	
BPLEU	bleu{,a}	%icc or %xcc, label	Branch on less or equal	C or Z
	{,pt ,pn}		unsigned	
BPCC	bcc{,a}	%icc or %xcc, label	Branch on carry clear	not C
	{,pt ,pn}		(greater than or equal, unsigned)	
BPCS	bcs{,a}	%icc or %xcc, label	Branch on carry set (less	С
	{,pt ,pn}		than, unsigned)	
BPPOS	bpos{,a}	%icc or %xcc, label	Branch on positive	not N
	{,pt ,pn}			
BPNEG	bneg{,a}	%icc or %xcc, label	Branch on negative	N
	{,pt ,pn}			
BPVC	bvc{,a}	%icc or %xcc, label	Branch on overflow clear	not V
	{,pt ,pn}			
BPVS	bvs{,a}	%icc or %xcc, label	Branch on overflow set	V
	{,pt ,pn}			
BRZ	brz{,a}	reg _{rs1} , label	Branch on register zero	Z
	{,pt ,pn}			
BRLEZ	brlez{,a}	reg _{rs1} , label	Branch on register less	N or Z
	{,pt ,pn}		than or equal to zero	
BRLZ	brlz{,a}	reg _{rs1} , label	Branch on register less	N
	{,pt ,pn}		than zero	

TABLE E-10(Continued)

Opcode	Mnemonic	Argument List	Operation	Comments
BRNZ	brnz{,a} {,pt ,pn}	reg _{rs1} , label	Branch on register not zero	not Z
BRGZ	brgz{,a} {,pt ,pn}	reg_{rs1} , label	Branch on register greater than zero	not (N or Z)
BRGEZ	<pre>brgez{,a} {,pt ,pn}</pre>	reg _{rs1} , label	Branch on register greater than or equal to zero	not N
CASA	casa	[reg _{rs1}]imm_asi,reg _{rs2} ,reg _{rd} [reg _{rs1}]%asi,reg _{rs2} ,reg _{rd}	Compare and swap word from alternate space	
CASXA	casxa	$[reg_{\rm rs1}]imm_asi,reg_{\rm rs2},reg_{\rm rd}$ $[reg_{\rm rs1}]\%asi,reg_{\rm rs2},reg_{\rm rd}$	Compare and swap extended from alternate space	
FBPA	fba{,a}	%fccn, label	(Branch on cc with prediction)	1
	{,pt ,pn}		Branch never	
FBPN	fbn{,a}	%fccn, label	Branch always	0
	{,pt ,pn}			
FBPU	fbu{,a}	%fccn, label	Branch on unordered	U
	{,pt ,pn}			
FBPG	fbg{,a}	%fccn, label	Branch on greater	G
	{,pt ,pn}			
FBPUG	fbug{,a}	%fccn, label	Branch on unordered or	G or U
	{,pt ,pn}		greater	
FBPL	fbl{,a}	%fccn, label	Branch on less	L
	{,pt ,pn}			
FBPUL	fbul{,a}	%fccn, label	Branch on unordered or	L or U
	{,pt ,pn}		less	
FBPLG	fblg{,a}	%fccn, label	Branch on less or greater	L or G
	{,pt ,pn}			
FBPNE	fbne{,a}	%fccn, label	Branch on not equal	L or G or U
	{,pt ,pn}			

(Continued) TABLE E-10

Opcode	Mnemonic	Argument List	Operation	Comments
FBPE	fbe{,a}	%fccn, label	Branch on equal	Е
	{,pt ,pn}			
FBPUE	fbue{,a}	%fccn,label	Branch on unordered or	E or U
	{,pt ,pn}		equal	
FBPGE	fbge{,a}	%fccn, label	Branch on greater or	E or G
	{,pt ,pn}		equal	
FBPUGE	fbuge{,a}	%fccn, label	Branch on unordered or	E or G or U
	{,pt ,pn}		greater or equal	
FBPLE	fble{,a}	%fccn, label	Branch on less or equal	E or L
	{,pt ,pn}			
FBPULE	fbule{,a}	%fccn, label	Branch on unordered or	E or L or u
	{,pt ,pn}		less or equal	
FBPO	fbo{,a}	%fccn, label	Branch on ordered	E or L or G
	{,pt ,pn}			
FLUSHW	flushw		Flush register windows	
FMOVA	fmov	%icc or %xcc, $freg_{rs2}$, $freg_{rd}$	(Move on integer cc)	1
	{s,d,q}a	17 01327 014	Move always	
FMOVN	fmov	%icc or %xcc, freg _{rs2} , freg _{rd}	Move never	0
	$\{s,d,q\}n$, 6,527 6,4		
FMOVNE	fmov	%icc or %xcc, freg _{rs2} , freg _{rd}	Move if not equal	not Z
	$\{s,d,q\}$ ne			
FMOVE	fmov	%icc or %xcc, freg _{rs2} , freg _{rd}	Move if equal	Z
	{s,d,q}e			
FMOVG	fmov	%icc or %xcc, $freg_{rs2}$, $freg_{rd}$	Move if greater	not (Z or (N
	{s,d,q}g			xor V))
FMOVLE	fmov	%icc or %xcc, $freg_{rs2}$, $freg_{rd}$	Move if less or equal	Z or (N xor V)
	{s,d,q}le			
FMOVGE	fmov	%icc or %xcc, $freg_{rs2}$, $freg_{rd}$	Move if greater or equal	not (N xor V)
	{s,d,q}ge			

TABLE E-10(Continued)

Opcode	Mnemonic	Argument List	Operation	Comments
FMOVL	fmov	%icc or %xcc, freg _{rs2} , freg _{rd}	Move if less	N xor V
	{s,d,q}l			
FMOVGU	fmov	%icc or %xcc, $freg_{rs2}$, $freg_{rd}$	Move if greater unsigned	not (C or Z)
	{s,d,q}gu			
FMOVLEU	fmov	%icc or %xcc, $freg_{rs2}$, $freg_{rd}$	Move if less or equal	C or Z
	{s,d,q}leu		unsigned	
FMOVCC	fmov	%icc or %xcc, $freg_{rs2}$, $freg_{rd}$	Move if carry clear	not C
	{s,d,q}cc		(greater or equal, unsigned)	
FMOVCS	fmov	%icc or %xcc, freg _{rs2} , freg _{rd}	Move if carry set (less	С
	{s,d,q}cs		than, unsigned)	
FMOVPOS	fmov	%icc or %xcc, $freg_{rs2}$, $freg_{rd}$	Move if positive	not N
	{s,d,q}pos			
FMOVNEG	fmov	%icc or %xcc, $freg_{rs2}$, $freg_{rd}$	Move if negative	N
	{s,d,q}neg			
FMOVVC	fmov	%icc or %xcc, $freg_{rs2}$, $freg_{rd}$	Move if overflow clear	not V
	{s,d,q}vc			
FMOVVS	fmov	%icc or %xcc, $freg_{rs2}$, $freg_{rd}$	Move if overflow set	V
	{s,d,q}vs			
FMOVRZ	fmovr	reg_{rs1} , $freg_{rs2}$, $freg_{rd}$	(Move f-p register on cc)	
1110 1112	{s,d,q}e	8rs1/1/8rs2/1/8rd	Move if register zero	
FMOVRLEZ	fmovr	reg_{rs1} , $freg_{rs2}$, $freg_{rd}$	Move if register less than	
	{s,d,q}lz	C131 7 C132 7 C14	or equal zero	
FMOVRLZ	fmovr	reg _{rs1} , freg _{rs2} , freg _{rd}	Move if register less than	
	{s,d,q}lz	0131-7 0132-7 01a	zero	

(Continued) TABLE E-10

Opcode	Mnemonic	Argument List	Operation	Comments
FMOVRNZ	fmovr	reg_{rs1} , $freg_{rs2}$, $freg_{rd}$	Move if register not zero	
FMOVRGZ	{s,d,q}ne	reg_{rs1} , $freg_{rs2}$, $freg_{rd}$	Move if register greater	
FMOVRGEZ	fmovr	reg_{rs1} , $freg_{rs2}$, $freg_{rd}$	than zero	
	{s,d,q}gz		Move if register greater than or equal to zero	
	fmovr		•	
	{s,d,q}gez			
FMOVFA	fmov{s,d,q}a	%fccn, freg _{rs2} , freg _{rd}	(Move on floating-point	1
FMOVFN	$fmov{s,d,q}n$	%fccn,freg _{rs2} ,freg _{rd}	cc)	0
FMOVFU	fmov{s,d,q}u	%fccn, freg _{rs2} , freg _{rd}	Move always	U
FMOVFG	fmov{s,d,q}g	%fccn, freg _{rs2} , freg _{rd}	Move never	G
FMOVFUG	fmov{s,d,q}ug	%fccn, freg _{rs2} , freg _{rd}	Move if unordered	G or U
FMOVFL	fmov{s,d,q}l	%fccn, freg _{rs2} , freg _{rd}	Move if greater	L
FMOVFUL	fmov{s,d,q}ul	%fccn, freg _{rs2} , freg _{rd}	Move if unordered or greater	L or U
FMOVFLG	fmov{s,d,q}lg	%fccn, freg _{rs2} , freg _{rd}	Move if less	L or G
FMOVFNE	$fmov{s,d,q}ne$	%fccn, freg _{rs2} , freg _{rd}	Move if unordered or	L or G or U
FMOVFE	fmov{s,d,q}e	%fccn, freg _{rs2} , freg _{rd}	less	E
FMOVFUE	fmov{s,d,q}ue	%fccn, freg _{rs2} , freg _{rd}	Move if less or greater	E or U
FMOVFGE	fmov{s,d,q}ge	%fccn, freg _{rs2} , freg _{rd}	Move if not equal	E or G
FMOVFUGE	fmov{s,d,q}uge	%fccn, freg _{rs2} , freg _{rd}	Move if equal	E or G or U
FMOVFLE	$fmov{s,d,q}le$	%fccn, freg _{rs2} , freg _{rd}	Move if unordered or	E or L
FMOVFULE	<pre>fmov{s,d,q}ule</pre>	%fccn, freg _{rs2} freg _{rd}	equal Move if greater or equal	E or L or u
FMOVFO	fmov{s,d,q}o	%fccn, freg _{rs2} , freg _{rd}	Move if greater or equal	E or L or G
			greater or equal	
			Move if less or equal	
			Move if unordered or less or equal	
			Move if ordered	
LDSW	ldsw	[address], reg _{rd}	Load a signed word	
LDSWA	ldsw	[regaddr] imm_asi, reg _{rd}	Load signed word from alternate space	

TABLE E-10(Continued)

Opcode	Mnemonic	Argument List	Operation	Comments
LDX	ldx	[address], reg _{rd}	Load extended word	
LDXA	ldxa	[regaddr] imm_asi, reg _{rd}	Load extended word	
LDXFSR	ldxa	[reg_plus_imm] %asi, reg _{rd}	from alternate space	
	ldx	[address], %fsr	Load floating-point state register	
MEMBAR	membar	membar_mask	Memory barrier	

(Continued) TABLE E-10

Opcode	Mnemonic	Argument List	Operation	Comments
MOVA	mova	%icc or %xcc, reg_or_imm11,	(Move integer register on	1
MOVN	movn	reg_{rd}	cc)	C
MOVNE	movne	%icc or %xcc, reg_or_imm11,	Move always	not Z
MOVE	move	reg_{rd}	Move never	Z
MOVG	movq	%icc or %xcc, reg_or_imm11, reg _{rd}	Move if not equal	not (Z or (N
MOVLE	movle	%icc or %xcc, reg_or_imm11,	Move if equal	xor V)
MOVGE		reg _{rd}	Move if greater	Z or (N xor V)
	movge	%icc or %xcc, reg_or_imm11,	Move if less or equal	not (N xor V)
MOVL	movl	reg _{rd}	Move if greater or equal	N xor V
MOVGU	movgu	%icc or %xcc, reg_or_imm11,	Move if less	not (C or Z)
MOVLEU	movleu	reg _{rd}	Move if greater unsigned	C or Z
MOVCC	movcc	%icc or %xcc, reg_or_imm11, reg _{rd}	Move if less or equal	not C
MOVCS	movcs	%icc or %xcc, reg_or_imm11,	unsigned	
MOVPOS	movpos	reg _{rd}	Move if carry clear	C
MOVNEG	movneg	%icc or %xcc, reg_or_imm11,	(greater or equal,	not N
MOVVC	movvc	reg_{rd}	unsigned)	N
MOVVS	movvs	%icc or %xcc, reg_or_imm11, reg _{rd}	Move if carry set (less than, unsigned)	not V V
		%icc or %xcc, reg_or_imm11,	Move if positive	,
		reg_{rd}	Move if negative	
		%icc or %xcc, reg_or_imm11,	Move if overflow clear	
		reg _{rd}	Move if overflow set	
		%icc or %xcc, reg_or_imm11, reg _{rd}		
		%icc or %xcc, reg_or_imm11, reg _{rd}		
		%icc or %xcc, reg_or_imm11, reg _{rd}		
		%icc or %xcc, reg_or_imm11, reg _{rd}		

(Continued) TABLE E-10

Opcode	Mnemonic	Argument List	Operation	Comments
MOVFA	mova	%fccn, reg_or_imm11, reg _{rd}	(Move on floating-point	1
MOVFN	movn	%fccn,reg_or_imm11,reg _{rd}	cc)	0
MOVFU	movu	%fccn,reg_or_imm11,reg _{rd}	Move always	U
MOVFG	movg	%fccn,reg_or_imm11,reg _{rd}	Move never	G
MOVFUG	movug	%fccn, reg_or_imm11, reg_rd	Move if unordered	G or U
MOVFL	movl	%fccn, reg_or_imm11, reg_rd	Move if greater	L
MOVFUL	movul	%fccn, reg_or_imm11, reg _{rd}	Move if unordered or greater	L or U
MOVFLG	movlg	%fccn, reg_or_imm11, reg _{rd}	Move if less	L or G
MOVFNE	movne	%fccn, reg_or_imm11, reg _{rd}	Move if unordered or	L or G or U
MOVFE	move	%fccn, reg_or_imm11, reg _{rd}	less	E E or U
MOVFUE	movue	%fccn, reg_or_imm11, reg _{rd}	Move if less or greater	E or G
MOVFGE	movge	%fccn, reg_or_imm11, reg _{rd}	Move if not equal	E or G or U
MOVFUGE	movuge	%fccn,reg_or_imm11,reg _{rd}	Move if equal	E or L
MOVFLE	movle	%fccn, reg_or_imm11, reg _{rd}	Move if unordered or	E or L or u
MOVFULE	movule	%fccn, reg_or_imm11, reg _{rd}	equal Move if greater or equal	E or L or G
MOVFO	movo	%fccn, reg_or_imm11, reg_rd	Move if unordered or	
			greater or equal	
			Move if less or equal	
			Move if unordered or less or equal	
			Move if ordered	
MOVRZ	movre	reg _{rs1} , reg_or_imm10,reg _{rd}	(Move register on	Z
MOVRLEZ	movrlez	reg _{rs1} , reg_or_imm10,reg _{rd}	register cc)	N or Z
MOVRLZ	movrlz	reg _{rs1} , reg_or_imm10,reg _{rd}	Move if register zero	N
MOVRNZ	movrnz	reg _{rs1} , reg_or_imm10,reg _{rd}	Move if register less than or equal to zero	not Z
MOVRGZ	movrgz	reg _{rs1} , reg_or_imm10,reg _{rd}	Move if register less than	N nor Z
MOVRGEZ	movrgez	reg _{rs1} , reg_or_imm10,reg _{rd}	zero	not N
			Move if register not zero	
			Move if register greater than zero	
			Move if register greater than or equal to zero	

TABLE E-10 (Continued)

Opcode	Mnemonic	Argument List	Operation	Comments
MULX	mulx	reg _{rs1} , reg_or_imm,reg _{rd}	(Generic 64-bit Multiply) Multiply (signed or unsigned)	See SDIVX and UDIVX
POPC	popc	reg_or_imm, reg _{rd}	Population count	
PREFETCH	prefetch	[address], prefetch_dcn [regaddr]	Prefetch data	See The SPARC
PREFETCHA	prefetcha	imm_asi, prefetch_fcn [rev_plus_imm] %asi.		architecture manual, version
	prefetcha		ancinate space	9
SDIVX	sdivx	reg _{rs1} , reg_or_imm,reg _{rd}	(64-bit signed divide) Signed Divide	See MULX and UDIVX
STX	stx	reg _{rd} , [address]	Store extended word	
STXA	stxa	reg _{rd} , [address] imm_asi	Store extended word	
STXFSR	stxa	reg _{rd} , [reg_plus_imm] %asi %tsr, [address]	into alternate space	
	stx		Store floating-point register (all 64-bits)	
UDIVX	udivx	reg _{rs1} , reg_or_imm, reg _{rd}	(64-bit unsigned divide) Unsigned divide	See MULX and SDIVX

E.4 SPARC-V9 Floating-Point Instruction Set Mapping

SPARC-V9 floating-point instructions are shown in the following table.

TABLE E-11

SPARC	Mnemonic ¹	Argument List	Description
F[sdq]TOx	fstox	$freg_{rs2}$, $freg_{rd}$	Convert floating point to 64-bit integer
	fdtox	$freg_{rs2}$, $freg_{rd}$	
	fqtox	freg _{rs2} , freg _{rd}	
	fstoi	$freg_{rs2}$, $freg_{rd}$	Convert floating-point to 32-bit integer
	fdtoi	freg _{rs2} , freg _{rd}	
	fqtoi	freg _{rs2} , freg _{rd}	
FxTO[sdq]	fxtos	freg _{rs2} , freg _{rd}	Convert 64-bit integer to floating point
	fxtod	freg _{rs2} , freg _{rd}	
	fxtoq	freg _{rs2} , freg _{rd}	
	fitos	freg _{rs2} , freg _{rd}	Convert 32-bit integer to floating point
	fitod	freg _{rs2} , freg _{rd}	
	fitoq	freg _{rs2} , freg _{rd}	
FMOV[dq]	fmovd	freg _{rs2} , freg _{rd}	Move double
	fmovq	freg _{rs2} , freg _{rd}	Move quad
FNEG[dq]	fnegd	freg _{rs2} , freg _{rd}	Negate double
	fnegq	freg _{rs2} , freg _{rd}	Negate quad
FABS [dq]	fabsd	freg _{rs2} , freg _{rd}	Absolute value double
	fabsq	freg _{rs2} , freg _{rd}	Absolute value quad
LDFA	lda	[regaddr] imm_asi, freg _{rd}	Load floating-point register from
LDDFA	lda	[reg_plus_imm] %asi, freg _{rd}	alternate space
LDQFA	ldda	[regaddr] imm_asi, freg _{rd}	Load double floating-point register from alternate space.
	ldda	[reg_plus_imm] %asi, freg _{rd}	Load quad floating-point register from
	ldqa	[regaddr] imm_asi, freg _{rd}	alternate space
	ldqa	[reg_plus_imm] %asi, freg _{rd}	

Types of Operands are denoted by the following lower-case letters:

i 32-bit integer

x 64-bit integer

s single

d double

 $^{{\}tt q} \, \text{quad} \,$

TABLE E-11 (Continued)

SPARC	Mnemonic ¹	Argument List	Description
STFA	sta	freg _{rd} , [regaddr] imm_asi	Store floating-point register to alternate
STDFA	sta	freg _{rd} , [reg_plus_imm] %asi	space
STQFA	stda	freg _{rd} , [regaddr] imm_asi	Store double floating-point register to alternate space
	stda	freg _{rd} , [reg_plus_imm] %asi	Store quad floating-point register to
	stqa	freg _{rd} , [regaddr] imm_asi	alternate space
	stqa	freg _{rd} , [reg_plus_imm] %asi	

Types of Operands are denoted by the following lower-case letters:

E.5 SPARC-V9 Synthetic Instruction-Set Mapping

Here is a mapping of synthetic instructions to hardware equivalent instructions.

TABLE E-12

Synthetic Instruction		Hardware Equivalent(s)		Comment
cas	[reg _{rsl}], reg _{rs2} , reg _{rd}	casa	[reg _{rs1}]ASI_P, reg _{rs2} , reg _{rd}	Compare & swap (cas)
casl	$[reg_{rsl}]$, reg_{rs2} , reg_{rd}	casa	[reg _{rsl}]ASI_P_L, reg _{rs2} ,	cas little-endian
casx	$[reg_{rsl}]$, reg_{rs2} , reg_{rd} $[reg_{rsl}]$, reg_{rs2} , reg_{rd}	casxa	reg_{rd} $[reg_{rsl}]$ ASI_P, reg_{rs2} , reg_{rd} $[reg_{rsl}]$ ASI_P_L, reg_{rs2} , reg_{rd}	cas extended cas little-endian, extended
clrx	[address]	stx	%g0, [address]	Clear extended word
clruw	reg _{rs1} , reg _{rd}	srl srl	reg_{rs1} , %g0, reg_{rd} reg_{rd} , %g0, reg_{rd}	Copy and clear upper word Clear upper word
iprefetch	label	bn, pt	%xcc, label	Instruction prefetch,

i 32-bit integer

x 64-bit integer

s single

d double

 $[\]mathtt{q} \, \text{quad} \,$

TABLE E-12(Continued)

Synthetic Instruction		Hardware Equivalent(s)		Comment
mov	%y, reg _{rd}	rd	%y, reg _{rd}	
mov	%asrn, reg _{rd}	rd	%asrn, reg _{rd}	
mov	reg_or_imm, %asrn	wr	%g0, reg_or_imm, %asrn	
ret		jmpl	%i7+8, %g0	Return from subroutine
retl		jmpl	%o7+8, %g0	Return from leaf subroutine
setn	value, r1, r2	for -xarch=v9s	ame as setx value r1, r2	
		for -xarch=v8 s	ame as set value r2	
setnhi	value, r1, r2	for -xarch=v9 s	ame as setxhi <i>value r1</i> ,	
		for -xarch=v8 s	ame as sethi value r2	
setuw	value,reg _{rd}	sethi	%hi(value), reg _{rd}	(value & 3FF ₁₆)==0
		or	%g0, value, reg _{rd}	when $0 \le \text{value} \le 4095$
		sethi	%hi(value), reg _{rd;}	(otherwise)
		or	reg _{rd} , %10(value), reg _{rd}	Do not use setuw in a DCTI delay slot.
setsw	value,reg _{rd}	sethi	%hi(value), reg _{rd}	value>=0 and (value &
		or	%g0, value, reg _{rd}	$3FF_{16}$)==0
		sethi	%hi(value), reg _{rd}	-4096 ≤ value ≤ 4095
		sra	reg _{rd} , %g0, reg _{rd}	if (value<0) and ((value & 3FF)==0)
		sethi	%hi(value), reg _{rd;}	(otherwise, if value>=0)
		or	reg _{rd} , %10(value), reg _{rd}	(otherwise, if value<0)
		sethi	%hi(value), reg _{rd} ;	Do not use setsw in a
		or	reg _{rd} , %10(value), reg _{rd}	CTI delay slot.
		sra	reg _{rd} , %g0, reg _{rd}	
setx	value, r1, r2	sethi	%hh(value), r1	
		or	r1, %hm(value), r1	
		sethi	%lm(value), r2	
		or	r2, %lo(value), r2	
		sllx	r1, 32, r1	
		or	r1, r2, r2	

TABLE E-12 (Continued)

Synthetic Instruction		Hardware Equivalent(s)		Comment
setxhi	value r1, r2	sethi	%hh(value), r1	
		or	r1, %hm(value), r1	
		sethi	%lm(value), r2	
		sllx	r1, 32, r1	
		or	r1, r2, r2	
signx	reg _{rsl,} reg _{rd}	sra	reg _{rsl,} %g0, reg _{rd}	Sign-extend 32-bit value
signx	reg _{rd}	sra	reg _{rsl,} %g0, reg _{rd} reg _{rd,} %g0, reg _{rd}	to 64 bits

E.6 UltraSPARC and VIS Instruction Set **Extensions**

This section describes extensions that require SPARC-V9. The extensions support enhanced graphics functionality and improved memory access efficiency.

Note – SPARC-V9 instruction set extensions used in executables may not be portable to other SPARC-V9 systems.

E.6.1 Graphics Data Formats

The overhead of converting to and from floating-point arithmetic is high, so the graphics instructions are optimized for short-integer arithmetic. Image components are 8 or 16 bits. Intermediate results are 16 or 32 bits.

E.6.2 Eight-bit Format

A 32-bit word contains pixels of four unsigned 8-bit integers. The integers represent image intensity values (α , G, B, R). Support is provided for band interleaved images (store color components of a point), and band sequential images (store all values of one color component).

E.6.3 Fixed Data Formats

A 64-bit word contains four 16-bit signed fixed-point values. This is the fixed 16-bit data format.

A 64-bit word contains two 8-bit signed fixed-point values. This is the fixed 32-bit data format.

Enough precision and dynamic range (for filtering and simple image computations on pixel values) can be provided by an intermediate format of fixed data values. Pixel multiplication is used to convert from pixel data to fixed data. Pack instructions are used to convert from fixed data to pixel data (clip and truncate to an 8-bit unsigned value). The FPACKFIX instruction supports conversion from 32-bit fixed to 16-bit fixed. Rounding is done by adding one to the rounding bit position. You should use floating-point data to perform complex calculations needing more precision or dynamic range.

E.6.4 SHUTDOWN Instruction

All outstanding transactions are completed before the SHUTDOWN instruction completes.

TABLE E-13

SPARC	Mnemonic	Argument List	Description
SHUTDOWN	shutdown		shutdown to enter power down mode

E.6.5 Graphics Status Register (GSR)

You use ASR 0x13 instructions RDASR and WRASR to access the Graphics Status Register.

SPARC	Mnemonic	Argument List	Description
RDASR	rdasr	%gsr, reg _{rd}	read GSR
WRASR	wrasr	reg _{rs1} , reg_or_imm, %gsr	write GSR

E.6.6 Graphics Instructions

Unless otherwise specified, floating-point registers contain all instruction operands. There are 32 double-precision registers. Single-precision floating-point registers contain the pixel values, and double-precision floating-point registers contain the fixed values.

The opcode space reserved for the Implementation-Dependent Instruction1 (IMPDEP1) instructions is where the graphics instruction set is mapped.

Partitioned add/subtract instructions perform two 32-bit or four 16-bit partitioned adds or subtracts between the source operands corresponding fixed point values.

TABLE E-15

SPARC	Mnemonic	Argument List	Description
FPADD16	fpadd16	freg _{rs1} , freg _{rs2} , freg _{rd}	four 16-bit add
FPADD16S	fpadd16s	freg _{rs1} , freg _{rs2} , freg _{rd}	two 16-bit add
FPADD32	fpadd32	freg _{rs1} , freg _{rs2} , freg _{rd}	two 32-bit add
FPADD32S	fpadd32s	freg _{rs1} , freg _{rs2} , freg _{rd}	one 32-bit add
FPSUB16	fpsub16	freg _{rs1} , freg _{rs2} , freg _{rd}	four 16-bit subtract
FPSUB16S	fpsub16s	freg _{rs1} , freg _{rs2} , freg _{rd}	two 16-bit subtract
FPSUB32	fpsub32	freg _{rs1} , freg _{rs2} , freg _{rd}	two 32-bit subtract
FPSUB32S	fpsub32s	freg _{rs1} , freg _{rs2} , freg _{rd}	one 32-bit subtract

Pack instructions convert to a lower pixel or precision fixed format.

TABLE E-16

SPARC	Mnemonic	Argument List	Description
FPACK16	fpack16	freg _{rs2} , freg _{rd}	four 16-bit packs
FPACK32	fpack32	freg _{rs1} , freg _{rs2} , freg _{rd}	two 32-bit packs
FPACKFIX	fpackfix	freg _{rs2} , freg _{rd}	four 16-bit packs
FEXPAND	fexpand	freg _{rs2} , freg _{rd}	four 16-bit expands
FPMERGE	fpmerge	$freg_{rs1}$, $freg_{rs2}$, $freg_{rd}$	two 32-bit merges

Partitioned multiply instructions have the following variations.

TABLE E-17

SPARC	Mnemonic	Argument List	Description
FMUL8x16	fmul8x16	freg _{rs1} , freg _{rs2} , freg _{rd}	8x16-bit partition
FMUL8x16AU	fmul8x16au	freg _{rs1} , freg _{rs2} , freg _{rd}	8x16-bit upper α partition
FMUL8x16AL	fmul8x16al	freg _{rs1} , freg _{rs2} , freg _{rd}	8x16-bit lower α partition
FMUL8SUx16	fmul8sux16	freg _{rs1} , freg _{rs2} , freg _{rd}	upper 8x16-bit partition
FMUL8ULx16	fmul8ulx16	freg _{rs1} , freg _{rs2} , freg _{rd}	lower unsigned 8x16-bit
FMULD8SUx16	fmuld8sux16	freg _{rs1} , freg _{rs2} , freg _{rd}	partition
FMULD8ULx16	fmuld8ulx16	freg _{rs1} , freg _{rs2} , freg _{rd}	upper 8x16-bit partition
			lower unsigned 8x16-bit partition

Alignment instructions have the following variations.

TABLE E-18

SPARC	Mnemonic	Argument List	Description
ALIGNADDRESS	alignaddr	reg _{rs1} , reg _{rs2} , reg _{rd}	find misaligned data access address
ALIGNADDRESS_LITTLE	alignaddrl	reg _{rs1} , reg _{rs2} , reg _{rd}	same as above, but little-endian
FALIGNDATA	faligndata	freg _{rs1} , freg _{rs2} , freg _{rd}	do misaligned data, data alignment

Logical operate instructions perform one of sixteen 64-bit logical operations between rs1 and rs2 (in the standard 64-bit version).

TABLE E-19

SPARC	Mnemonic	Argument List	Description
FZERO	fzero	freg _{rd}	zero fill
FZEROS	fzeros	freg _{rd}	zero fill, single precision
FONE	fone	freg _{rd}	one fill
FONES	fones	freg _{rd}	one fill, single precision
FSRC1	fsrc1	$freg_{rs1}$, $freg_{rd}$	copy src1

(Continued) TABLE E-19

SPARC	Mnemonic	Argument List	Description
FSRC1S	fsrc1s	freg _{rs1} , freg _{rd}	copy src1, single precision
FSRC2	fsrc2	freg _{rs2} , freg _{rd}	copy src2
FSRC2S	fsrc2s	freg _{rs2} , freg _{rd}	copy src2, single precision
FNOT1	fnot1	freg _{rs1} , freg _{rd}	negate src1, 1's complement
FNOT1S	fnot1s	freg _{rs1} , freg _{rd}	same as above, single precision
FNOT2	fnot2	freg _{rs2} , freg _{rd}	negate src2, 1's complement
FNOT2S	fnot2s	freg _{rs2} , freg _{rd}	same as above, single precision
FOR	for	freg _{rs1} , freg _{rs2} , freg _{rd}	logical OR
FORS	fors	freg _{rs1} , freg _{rs2} , freg _{rd}	logical OR, single precision
FNOR	fnor	freg _{rs1} , freg _{rs2} , freg _{rd}	logical NOR
FNORS	fnors	freg _{rs1} , freg _{rs2} , freg _{rd}	logical NOR, single precision
FAND	fand	freg _{rs1} , freg _{rs2} , freg _{rd}	logical AND
FANDS	fands	freg _{rs1} , freg _{rs2} , freg _{rd}	logical AND, single precision
FNAND	fnand	freg _{rs1} , freg _{rs2} , freg _{rd}	logical NAND
FNANDS	fnands	freg _{rs1} , freg _{rs2} , freg _{rd}	logical NAND, single precision
FXOR	fxor	freg _{rs1} , freg _{rs2} , freg _{rd}	logical XOR
FXORS	fxors	freg _{rs1} , freg _{rs2} , freg _{rd}	logical XOR, single precision
FXNOR	fxnor	freg _{rs1} , freg _{rs2} , freg _{rd}	logical XNOR
FXNORS	fxnors	freg _{rs1} , freg _{rs2} , freg _{rd}	logical XNOR, single precision
FORNOT1	fornot1	freg _{rs1} , freg _{rs2} , freg _{rd}	negated src1 OR src2
FORNOT1S	fornot1s	freg _{rs1} , freg _{rs2} , freg _{rd}	same as above, single precision
FORNOT2	fornot2	freg _{rs1} , freg _{rs2} , freg _{rd}	src1 OR negated src2
FORNOT2S	fornot2s	freg _{rs1} , freg _{rs2} , freg _{rd}	same as above, single precision
FANDNOT1	fandnot1	freg _{rs1} , freg _{rs2} , freg _{rd}	negated src1 AND src2
FANDNOT1S	fandnot1s	freg _{rs1} , freg _{rs2} , freg _{rd}	same as above, single precision
FANDNOT2	fandnot2	freg _{rs1} , freg _{rs2} , freg _{rd}	src1 AND negated src2
FANDNOT2S	fandnot2s	freg _{rs1} , freg _{rs2} , freg _{rd}	same as above, single precision

Pixel compare instructions compare fixed-point values in rs1 and rs2 (two 32 bit or four 16 bit)

TABLE E-20

SPARC	Mnemonic	Argument List	Description
FCMPGT16	fcmpgt16	freg _{rs1} , freg _{rs2} , reg _{rd}	4 16-bit compare, set rd if src1>src2
FCMPGT32	fcmpgt32	freg _{rs1} , freg _{rs2} , reg _{rd}	2 32-bit compare, set rd if src1>src2
FCMPLE16	fcmple16	freg _{rs1} , freg _{rs2} , reg _{rd}	4 16-bit compare, set rd if src1≤src2
FCMPLE32	fcmple32	freg _{rs1} , freg _{rs2} , reg _{rd}	2 32-bit compare, set rd if src1≤src2
FCMPNE16	fcmpne16	freg _{rs1} , freg _{rs2} , reg _{rd}	4 16-bit compare, set rd if src1≠src2
FCMPNE32	fcmpne32	freg _{rs1} , freg _{rs2} , reg _{rd}	2 32-bit compare, set rd if src1≠src2
FCMPEQ16	fcmpeq16	freg _{rs1} , freg _{rs2} , reg _{rd}	4 16-bit compare, set rd if src1=src2
FCMPEQ32	fcmpeq32	freg _{rs1} , freg _{rs2} , reg _{rd}	2 32-bit compare, set rd if src1=src2

Edge handling instructions handle the boundary conditions for parallel pixel scan line loops.

TABLE E-21

SPARC	Mnemonic	Argument List	Description
EDGE8	edge8	reg_{rs1} , reg_{rs2} , reg_{rd}	8 8-bit edge boundary processing
EDGE8L	edge81	reg_{rs1} , reg_{rs2} , reg_{rd}	same as above, little-endian
EDGE16	edge16	reg_{rs1} , reg_{rs2} , reg_{rd}	4 16-bit edge boundary processing
EDGE16L	edge16l	reg_{rs1} , reg_{rs2} , reg_{rd}	same as above, little-endian
EDGE32	edge32	reg_{rs1} , reg_{rs2} , reg_{rd}	2 32-bit edge boundary processing
EDGE32L	edge321	reg_{rs1} , reg_{rs2} , reg_{rd}	same as above, little-endian

Pixel component distance instructions are used for motion estimation in video compression algorithms.

TABLE E-22

SPARC	Mnemonic	Argument List	Description
PDIST	pdist	freg _{rs1} , freg _{rs2} , freg _{rd}	8 8-bit components, distance between

The three-dimensional array addressing instructions convert three-dimensional fixed-point addresses (in rs1) to a blocked-byte address. The result is stored in rd.

TABLE E-23

SPARC	Mnemonic	Argument List	Description
ARRAY8	array8	reg _{rs1} , reg _{rs2} , reg _{rd}	convert 8-bit 3-D address to blocked
ARRAY16	array16	reg _{rs1} , reg _{rs2} , reg _{rd}	byte address
ARRAY32	array32	reg _{rs1} , reg _{rs2} , reg _{rd}	same as above, but 16-bit
		0131 0132 014	same as above, but 32-bit

E.6.7 Memory Access Instructions

These memory access instructions are part of the SPARC-V9 instruction set extensions.

TABLE E-24

SPARC	imm_asi	Argument List	Description
STDFA	ASI_PST8_P	stda freg _{rd} , [freg _{rs1}] reg _{mask} ,	eight 8-bit conditional stores to:
STDFA	ASI_PST8_S	imm_asi	primary address space
STDFA	ASI_PST8_PL		secondary address space
STDFA	ASI_PST8_SL		primary address space, little endian secondary address space, little endian
STDFA	ASI_PST16_P		four 16-bit conditional stores to:
STDFA	ASI_PST16_S		primary address space
STDFA	ASI_PST16_PL		secondary address space
STDFA	ASI_PST16_SL		primary address space, little endian secondary address space, little endian
STDFA	ASI_PST32_P		two 32-bit conditional stores to:
STDFA	ASI_PST32_S		primary address space
STDFA	ASI_PST32_PL		secondary address space
STDFA	ASI_PST32_SL		primary address space, little endian secondary address space, little endian

Note – To select a partial store instruction, use one of the partial store ASIs with the STDA instruction.

TABLE E-25

SPARC	imm_asi	Argument List	Description
LDDFA STDFA	ASI_FL8_P	ldda [reg_addr] imm_asi, freq _{rd} stda freq _{rd} , [reg_addr] imm_asi	8-bit load/store from/to: primary address space
LDDFA STDFA	ASI_FL8_S	ldda [reg_plus_imm] %asi, freq _{rd} stda [reg_plus_imm] %asi	secondary address space
LDDFA STDFA	ASI_FL8_PL		primary address space, little endian
LDDFA STDFA	ASI_FL8_SL		secondary address space, little endian
LDDFA	ASI_FL16_P		16-bit load/store from/to: primary address space
LDDFA STDFA	ASI_FL16_S		secondary address space
LDDFA STDFA	ASI_FL16_PL		primary address space, little endian
LDDFA STDFA	ASI_FL16_SL		secondary address space, little endian

Note – To select a short floating-point load and store instruction, use one of the short ASIs with the LDDA and STDA instructions.

TABLE E-26

SPARC	imm_asi	Argument List	Description
LDDA	ASI_NUCLEUS_QUAD_LDD	[reg_addr] imm_asi, reg _{rd}	128-bit atomic load
LDDA	ASI_NUCLEUS_QUAD_LDD_L	[reg_plus_imm] %asi, reg _{rd}	128-bit atomic load, little endian

(Continued) TABLE E-26

SPARC	imm_asi	Argument List	Description
LDDFA STDFA	ASI_BLK_AIUP	ldda [reg_addr] imm_asi, freq _{rd} stda freq _{rd} , [reg_addr] imm_asi	64-byte block load/store from/to: primary address space, user privilege
LDDFA STDFA	ASI_BLK_AIUS	ldda [reg_plus_imm] %asi, freq _{rd} stda freg _{rd} , [reg_plus_imm] %asi	secondary address space, user privilege.
LDDFA STDFA	ASI_BLK_AIUPL		primary address space, user privilege, little endian
LDDFA	ASI_BLK_AIUSL		secondary address space, user privilege little endian
LDDFA STDFA	ASI_BLK_P		primary address space
LDDFA	ASI_BLK_S		secondary address space
LDDFA STDFA	ASI_BLK_PL		primary address space, little endian
LDDFA	ASI_BLK_SL		secondary address space, little endian
LDDFA	ASI_BLK_COMMIT_P		64-byte block commit store to primary address space
LDDFA STDFA	ASI_BLK_COMMIT_S		64-byte block commit store to secondary address space

Note – To select a block load and store instruction, use one of the block transfer ASIs with the LDDA and STDA instructions.

Index

addresses, 32 .alias, 55 .align, 55 as command, 67 .ascii, 55 .asciz, 55 assembler command line, 67 assembler command line options, 68 assembler directives, 35 types, 35 assembly language, 13 lines, 14 statements, 14 syntax notation, 13 assignment directive, 36	comment lines, multiple, 15 .common, 55 constants, 15 decimal, 15 floating-point, 15 hexadecimal, 15 octal numeric, 15 Control Transfer Instructions (CTI), 20 converting existing object files, 37 coprocessor instruction, 50 cp_disabled trap, 50 cp_exception trap, 50 current location, 32 current section, 25
B binary operations, 18 .byte, 55 byte order for V9, 81	D -D option, 68 data generating directives, 36 default output file, 23 dis program, 71 disassembling object code, 71 .double, 56
case distinction, 14 case distinction, in special symbols, 18 cc language driver, 67 command-line options, 68 comment lines, 15	E ELF header, 24, 25 ehsize, 24 entry, 24 flag, 24 ident, 24

ELF header (Continued)	1
machine, 24	-I option, 68
phentsize, 24	ident, 57
phnum, 24	instruction set, used by assembler, 39
phoff, 24	instruction set changes (V9), 81
shentsize, 24	instruction set extensions (V9), 97
shnum, 24	instructions
shoff, 25	assembly language, 41
shstrndx, 25	hardware integer, 41
type, 25	integer instructions, 41
version, 25	integer suffixes, 15
.empty pseudo-operation, 21	invoking, as command, 67
.empty, 56	
error messages, 20	
escape codes, in strings, 16	
Executable and Linking Format (ELF) files, 12,	K
23	-K option, 68
expressions, 18	•
expressions, SPARC-V9, 19	
	L
_	-L option, 68
F	labeling format, 12
f77 language driver, 67	labels, 15
fbe command, 67	language drivers, 67
features, lexical, 14	lexical features, 14
.file, 56	lines syntax, 14
file syntax, 14	local, 57
floating-point instructions, 48	location counter, 32
floating-point pseudo-operations, 15	locations, 32
•	
G	M
.global, 56	-m option, 68
.globl, 56	multiple comment lines, 15
	multiple files, on, 67
	multiple sections, 26
н	multiple strings, in string table, 35
••	
.half, 57	
hardware instructions, SPARC architecture, 39	A.I
hardware integer, assembly language	N
instructions, 41	.noalias pseudo-op, 55
hyphen (-), 67	.noalias, 57
	.nonvolatile, 57
	numbers, 15
	numeric labels, 15

0	section header, 26, 29
-o option, 69	addr, 26
object file format, 12	addralign, 26
object files	entsize, 26
type, 12, 23	flags, 26
operators, 18	info, 27
operators, SPARC-V9, 19	link, 27
optim, 57	name, 27
options, command-line, 68	offset, 27
•	size, 27
	type, 27
	sections, 25
P	.seg, 59
-P option, 69	.single, 59
percentage sign (%), 16	.size, 59
popsection, 58	.skip, 59
predefined non-user sections, 31	SPARC-V9, 79
predefined user sections, 29	8-bit format, 97
.proc, 58	alternate space access, 81
pseudo-operations, 55	byte order, 81
pseudo-ops, examples of, 63	fixed data formats, 98
pushsection, 58	floating-point instructions, 93
	graphics data formats, 97
	instruction set changes, 81
	instruction set extensions, 97
Q	instruction set mapping, 84
-Q option, 69	registers, 79
-q option, 69	synthetic instruction set, 95
.quad, 58	SPARC-V9, 64-bit expressions, 19
	SPARC-V9, 64-bit operators, 19
	special floating-point values, 15
	special names, floating point values, 15
R	special symbols, 16
references, 7	.stabn, 60
registers, 16	.stabs, 60
relocatable files, 12, 23	statement syntax, 14
relocation tables, 33	string tables, 34
.reserve, 58	strings, 15
	multiple in string table, 35
	multiple references in string table, 35
	suggested style, 15
S	unreferenced in string table, 35
-S option, 69	sub-strings in string table, references to, 35
-s option, 69	symbol, 61
-sb option, 68	symbol attribute directives, 35
section, 58	symbol names, 16
section control directives, 35	symbol table, 33
section control pseudo-ops, 35	info, 33

```
symbol table (Continued)
                                                     -xarch=v8plusa option, 70
  name, 33
                                                     .xstabs, 61
  other, 33
  shndx, 33
  size, 33
  value, 33
symbol tables, 33
syntax notation, 13
synthetic instructions, 51
-T option, 69
table notation, 39
trap numbers, reserved, 47
.type, 60
U
-U option, 70
.uahalf, 60
.uaword, 60
unary operators, 18
user sections, 35
/usr/include/sys/trap.h, 47
-V option, 70
.version, 60
.volatile, 60
W
.weak, 61
.word, 61
X
-xarch=v7 option, 70
-xarch=v8 option, 70
-xarch=v8a option, 70
-xarch=v8plus option, 70
```

110 SPARC Assembly Language Reference Manual • May 2002