



School of Mathematical and Computational Sciences
Indian Association for the Cultivation of Science

Master's/Integrated Master's-PhD Program/ Integrated
Bachelor's-Master's Program/PhD Course

Theory of Computation II: COM 5108

Quiz III (16 November 2023)

Answer All Questions with Proper Justifications

Marks: $6 \times 3 = 18$

1. What is the maximum possible number of configurations of an $f(n) \geq \log n$ space bounded (work tape) TM, where n is the input length.
2. Let $\phi(C_1, C_2, t)$ be a TQBF formula corresponding to a $O(n^k)$ space bounded computation of a TM going from configuration C_1 to configuration C_2 in t steps. An intermediate configuration C_m is such that both C_1 to C_m and C_m to C_2 are reached in $\frac{t}{2}$ steps. So we can write

$$\phi(C_1, C_2, t) = \exists C_m (\phi(C_1, C_m, t/2) \wedge \phi(C_m, C_2, t/2)).$$

Explain why this scheme cannot give a polynomial size formula.

3. Let $A = \{x \in \{0, 1\}^* : \text{number of 0's in } x \text{ is twice the number of 1's in it}\}$. Is $A \in \mathbf{L}$?
4. We define the language

$$H_f = \{ \langle M, x \rangle : \text{the DTM } M \text{ accepts the input } x \text{ in } f(|x|) \text{ steps} \}.$$

We claim that $H_f \notin \mathbf{DTIME}(f(\lfloor \frac{n}{2} \rfloor))$. We prove this by *diagonalization*. Suppose the TM M_f decides H_f within $f(\lfloor \frac{n}{2} \rfloor)$ number of steps. We construct the following TM D_f for diagonalization.

D_f : input: $\langle M \rangle$

Simulate N_f on $\langle M, M \rangle$.

if N_f *accepts*, then **reject** else **accept**.

- (a) Find the running time of D_f on input of length $n = |\langle M \rangle|$.
 - (b) Apply D_f on its own description $\langle D_f \rangle$ and give the argument for contradiction.
5. $A_{cof} = \{ \langle M \rangle : \overline{L(M)} = \Sigma^* \setminus L(M) \text{ is finite} \} \in \Sigma_n^0$. What is the least value of n ? Give a definition of A_{cof} based on a recursive predicate.
 6. People claim that $\mathbf{PH} = \bigcup_{n \geq 0} \Sigma_n^p$ cannot have a complete language. If L is a \mathbf{PH} -complete language it must belong to Σ_i^p for some i . Explain what can be the conclusion in that case?