

INDIAN INSTITUTE OF TECHNOLOGY, KHARAGPUR

Date FN / AN Time: 2/3 Hrs. Full Marks No. of Students
Autumn / Spring Semester, 20 Deptt. Sub No.
.....Yr. B. Tech.(Hons.) / B. Arch. / M. Sc. Sub. Name

Instructions : Answer Q1. & Any Three (3) from the Remaining Questions (2-5). Do not write illogical statements.

1. [10 × 3 = 30]

Answer with a short justification whether the following statements are *true* or *false*. No credit will be given for writing only *true* or *false*.

- (a) Every language over the alphabet $\{0, 1\}$ is **Turing Recognisable (acceptable)**.
- (b) Every **Turing Decidable** language is a **context-free** language.
- (c) Some context-free languages may have both unambiguous as well as ambiguous grammars.
- (d) Any *non-null* string x of a context-free language can always be **derived** in $2n - 1$ steps (using a suitable grammar), where $|x| = n$.
- (e) If L_1 and L_2 are context-free languages, then so is

$$L = L_1L_2 = \{w : \exists x \in L_1, \exists y \in L_2, w = xy\}.$$

- (f) The following language is **not Turing Decidable**

$$L_{\phi}^{CFG} = \{ \langle G \rangle : G \text{ is a CFG and } L(G) = \phi \}.$$

- (g) The following language is **Turing Recognisable (acceptable)**

$$L_R^{TM} = \{ \langle M, x \rangle : M \text{ is a Turing Machine and } M \text{ does not accept } x \}.$$

- (h) A **linear bounded automaton** can be used as a **decider** of the language $L = \{a^n b^n c^n : n > 1\}$.
- (i) If a language L is **polynomial time mapping reducible** to the language L' and L is **NP-complete**, then so is L' .
- (j) If **P = NP** then every language $L \in P$ except $L = \emptyset$ and $L = \Sigma^*$ is **NP-complete**.

2. [7 + 8]

- (a) Use Cocke-Kasami-Younger (CKY) algorithm to show that **aabbab** can be generated by the grammar $G = (\{S, A, B, C, D\}, \{a, b\}, R, S)$, where

$$\begin{aligned} S &\rightarrow AB \mid BA \mid AC \mid BD \mid SS \\ C &\rightarrow SB \\ D &\rightarrow SA \\ A &\rightarrow a \text{ and } B \rightarrow b \end{aligned}$$

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- (b) Give a context-free grammar for the following language and convert the grammar to Chomsky Normal Form: $L = \{a^k b^m c^n : k, m \geq 1, k \leq n \leq 2k\}$.

3. [10 + 5]

- (a) Design a **single tape deterministic Turing machine** $M = (Q, \Sigma, \Gamma, \delta, q_0, q_A, q_R)$ that accepts the following language in $O(n \log n)$ time : $L = \{a^n b^n : n \geq 0\}$. Clearly specify each component of M and give the state transition function δ as a state transition diagram (give short explanation for each transition). Assume that there is a special **marker** $\$ \notin \Sigma$ at the leftmost square of the tape.
- (b) Give the outline of proof of the following statement:
 $L = \{ \langle G \rangle : G \text{ is a CFG over } \{0, 1\} \text{ and } \{1\}^* \cap L(G) \neq \emptyset \}$ is a **decidable language**.

4. [3 + 7 + 5]

- (a) Prove that the following language is **Turing Recognisable (acceptable)**:
 $L_{HALT}^{TM} = \{ \langle M, x \rangle : M \text{ is a Turing machine and } M \text{ halts on } x \}$.
- (b) Use **diagonalisation** and **reduction to a contradiction** (not the reduction of another language to L_{HALT}^{TM}) to prove that the language L of (4a) is **undecidable**.
- (c) Let L_1 and L_2 be two disjoint languages. Say that language L **separates** L_1 and L_2 if $L_1 \subseteq L$ and $L_2 \subseteq \bar{L}$. Show that two disjoint **co-Turing-Recognisable** languages are separated by some **decidable language**.

5. [4 + 7 + 4]

- (a) Give justification for the following statement:
All deterministic Turing machine models are polynomially equivalent.
- (b) Prove that the following language

$$TRIANGLE = \{ \langle G \rangle : G \text{ is a undirected graph and } G \text{ has a } \mathbf{3}\text{-clique} \},$$

is in the complexity class **P**. A 3-clique is a complete subgraph of three vertices.

- (c) **3SAT** is known to be an NP-complete language. Show by polynomial time mapping reduction that **VERTEX-COVER** is also an NP-complete language.

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