

# Formal Language and Automata Theory (CS21004)

## Tutorial - IX

Class: CSE 2<sup>nd</sup> Year

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**Exercise 1.** Consider the following CFG in CNF and show using the CYK table construction that the string 'aaaaa' is in the language of the grammar  $L(G)$ .

$$\begin{aligned}S &\rightarrow AB \mid BC \\A &\rightarrow BA \mid a \\B &\rightarrow DD \mid b \\D &\rightarrow AB \mid a\end{aligned}$$

where  $S$  is the start symbol,  $\{S, A, B, D\}$  are the non terminals and  $\{a, b\}$  are the terminals.

**Exercise 2.** A language  $L$  is *prefix-closed* if for every  $x \in L$ , every prefix of  $x$  is also in  $L$ . Let  $L$  be an *infinite, prefix-closed, context-free language*. Show that  $L$  has an *infinite regular subset*.

**Exercise 3.** Design a DTM that accepts the language  $L_1 = \{a^i b^j c^k : i, j, k \geq 1, i + j = k\}$ .

**Exercise 4.** Design a DTM that converts a string  $a^n$  to the binary representation of  $n$ . In fact it transforms an input  $\overbrace{aaa\dots aa}^n$  to binary representation of  $n$ .

**Exercise 5.** What will be the running time complexity of a singly-infinite tape DTM that recognises the language  $L = \{x \in \{a, b\}^* : x = x^R\}$ ? Can you think of a Turing machine model where it can be better?

**Exercise 6.** Convert the machine of *exercise 3* so that either it will accept a string or it will loop-forever. The language is  $L_1 = \{a^i b^j c^k : i, j, k \geq 1, i + j = k\}$ .