Computer Science & Engineering Department IIT Kharagpur Computational Number Theory: CS60094 Lecture IX

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QED.

1 Test for prime II

1.1 Primality by Quadratic Residue

It is known from the Euler's Criterion that for an odd prime p, if an integer a is a quadratic residue modulo p,¹ then $a^{\frac{p-1}{2}} \equiv 1 \pmod{p}$. If a is a quadratic non-residue modulo p, then $a^{\frac{p-1}{2}} \equiv p-1 \equiv -1 \pmod{p}$.

Similarly, the Legendre and Jacobi symbol coincides for a prime². $\left(\frac{a}{p}\right) = 1$ if a is a quadratic residue modulo p, and is -1 if a is a quadratic non-residue modulo p.

We can compute $a^{\frac{p-1}{2}} \mod n$ by fast exponentiation and we also have efficient algorithm to compute $\left(\frac{a}{n}\right)$.

Proposition 1. If p is an odd prime, then

$$a^{\frac{p-1}{2}} \times \left(\frac{a}{p}\right) \equiv 1 \pmod{p} \text{ for all } a \in \mathbb{Z}_p^*.$$

The equivalent contrapositive statement is

If $n \geq 3$ is an odd integer, $a \in \mathbb{Z}_n^*$ and $a^{\frac{n-1}{2}} \times (\frac{a}{n}) \not\equiv 1 \pmod{n}$, then n cannot be prime.

<u>Example 1.</u> Let n = 15. The elements of $\mathbb{Z}_{15}^* = \{1, 2, 4, 7, 8, 11, 13, 14\}$. The values of $\left(\frac{a}{15}\right) \times a^7 \mod 15$ are as follows:

a	1	2	4	7	8	11	13	14
$\left(\frac{a}{15}\right) \times a^7 \mod 15$	1	8	4	2	2	4	8	1

So there are large number of witnesses to show that 15 is a composite number. <u>Definition 1</u>: Let $n \ge 3$ be an odd composite number. An integer $a \in \mathbb{Z}_n$ is called an Euler-witness or *E*-witness for *n* if $\left(a^{\frac{n-1}{2}} \mod n\right) \times \left(\frac{a}{n}\right) \ne 1$. Otherwise it is called an *E*-liar,

$$L_n^E = \{ a \in \mathbb{Z}_n^* : \left(a^{\frac{n-1}{2}} \times \left(\frac{a}{n} \right) \right) \equiv 1 \pmod{n} \}$$

<u>Example 2.</u> Consider n = 225 and a = 26. We have $26^{\frac{225-1}{2}} \mod 225 = 26^{112} \mod 225 = 1$ and

$$\left(\frac{26}{225}\right) = \left(\frac{13}{15^2}\right)\left(\frac{2}{15^2}\right) = 1.$$

So 26 is an *E*-liar for the composite 225. But $2^{112} \mod 225 = 196$, so 2 is an *E*-witness of compositeness of 225.

A few other E-liars for 225 are 82, 107, 118, 143 etc.

Proposition 2. Let $n \ge 3$ be an odd composite number. Then every *E*-liar for *n* is also an *F*-liar for *n*. QED.

 $^{{}^{1}}p \not| a$ and a is a perfect square modulo p.

²We have seen that for a composite number n, $\left(\frac{a}{n}\right) = 1$ does not mean that a is a quadratic residue modulo n. As an example $\left(\frac{3}{35}\right) = \left(\frac{3}{5}\right)\left(\frac{3}{7}\right) = (-1) \times (-1) = 1$. But There is no solution of $x^2 \equiv 3 \pmod{35}$.

Proof: Let *a* be an *E*-liar for *n*. So $a^{\frac{n-1}{2}} \times (\frac{a}{n}) \equiv 1 \pmod{n}$. The value of $(\frac{a}{n}) \in \{1, -1\}$ (it cannot be equal to 0 as the product is congruent to 1). So the value of

$$\left(a^{\frac{n-1}{2}} \times \left(\frac{a}{n}\right)\right)^2 \equiv 1 \pmod{n}.$$

The value of $\left(\frac{a}{n}\right)^2 = 1$, implies that $a^{n-1} \equiv 1 \pmod{n}$ i.e. *a* is an *F*-liar for *n*. QED.

If $L_n^E \subseteq L_n^F$, then an *F*-witness of *n* is also an *E*-witness of *n* i.e. $W_n^F \subseteq W_n^E$. We prove that more than half of the elements of \mathbb{Z}_n^* are *E*-witnesses.

Proposition 3. Let $n \ge 3$ be an odd composite number. The set of *E*-liars of n, L_n^E , is a proper subgroup of \mathbb{Z}_n^* . QED.

Proof: We know $L_n^E \subseteq L_n^F \subseteq \mathbb{Z}_n^*$. We prove that L_n^E is closed under the group operation.

Let
$$a, b \in L_n^E$$

$$(a \cdot b)^{\frac{n-1}{2}} \times \left(\frac{a \cdot b}{n}\right) \equiv \left(a^{\frac{n-1}{2}} \times \left(\frac{a}{n}\right)\right) \cdot \left(b^{\frac{n-1}{2}} \times \left(\frac{b}{n}\right)\right) \equiv 1 \cdot 1 \equiv 1 \pmod{n}.$$

Finally we show that all elements of \mathbb{Z}_n^* are not *E*-liars. There is at least one *E*-witness in \mathbb{Z}_n^* . This will imply L_n^E is a proper subgroup. We consider two cases.

Case I: Let $n = p^k \cdot m$, where p is an odd prime, $k \ge 2$ and m is an odd number relatively prime to p.

(a) If m = 1, we choose a = p + 1 and claim that a is an F-witness of n. So it is also an E-witness of n.

 $gcd(a,n)=gcd(p+1,p^k)=1$ implies that $a\in\mathbb{Z}_n^*.$ Now we show that a is an F-witness.

If a is an F-liar, then $a^{n-1} \equiv 1 \pmod{n}$, implies that $a^{n-1} \equiv 1 \pmod{p^2}$ as p^2 is a divisor of n. So we have

$$a^{n-1} \equiv (1+p)^{n-1} \equiv 1 + (n-1)p + \sum_{2 \le i \le n-1} \binom{n-1}{i} p^i \equiv 1 + (n-1)p \pmod{p^2}.$$

 $p^2|a^{n-1}-1$, so $p^2|(n-1)p$. Hence p|n-1. But that is impossible as p|n. So a is an *E*-witness.

(b) If $m \ge 3$, we take a as a solution of the pair of congruence

$$\begin{array}{rcl} x &\equiv& 1+p \ (\mathrm{mod} \ p^2), \\ x &\equiv& 1 \ (\mathrm{mod} \ m) \end{array}$$

By the CRT there is a solution $a, 1 \le a < p^2 \cdot m \le n$.

As $a \equiv 1 + p \pmod{p^2}$, $a \equiv 1 + p \pmod{p}$, implies that p|a - 1. Hence $gcd(a, p^k) = 1$.

Similarly, m|a-1, so gcd(a,m) = 1. As $n = p^k \cdot m$, the value of gcd(a,n) cannot be larger than 1. Hence $a \in \mathbb{Z}_n^*$.

Our proof that this a cannot be an F-liar is similar to the case of m = 1.

Case II: n may be square-free and is a product of several distinct primes. Let $n = p \cdot m$ where p is an odd prime and $m \ge 3$ is an odd square free integer so that $p \not| m$.

Let $b \in \mathbb{Z}_p^*$ be a quadratic non-residue modulo p i.e. $\left(\frac{b}{p}\right) = -1$. We consider the following congruence

$$\begin{array}{rcl} x & \equiv & b \pmod{p}, \\ x & \equiv & 1 \pmod{m}. \end{array}$$

By CRT there is a solution a of this pair of congruence, $1 \le a . We prove that <math>a \in \mathbb{Z}_n^*$ and a is an *E*-witness.

p|a-b and $1 \leq b < p,$ so $p \not |a.$ Also gcd(a,m)=1, hence $gcd(a,n)=gcd(a,p\cdot m)=1.$ So $a\in \mathbb{Z}_n^*.$ We also have

$$\left(\frac{a}{n}\right) = \left(\frac{a}{p}\right) \cdot \left(\frac{a}{m}\right) = \left(\frac{b}{p}\right) \cdot \left(\frac{1}{m}\right) = -1 \cdot 1 = -1.$$

If a is an *E*-liar, then $a^{\frac{n-1}{2}} \equiv -1 \pmod{n}$. On the other hand m is a divisor of n, so $a^{\frac{n-1}{2}} \equiv -1 \pmod{m}$. But that contradicts the fact that $a \equiv 1 \pmod{m}$. So a is an *E*-witness of n. QED.

The size of *E*-liar is $\leq \frac{\phi(n)}{2} \leq \frac{n-2}{2}$. So at least half of the elements of \mathbb{Z}_n^* are *E*-witnesses.

1.2 Solovay-Strassen Test

R. Solovay and V. Strassen proposed the following randomized test in 1977.

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\begin{split} \text{isprimeSS}(n) \ // \ n \ \text{is odd} &\geq 3\\ a \leftarrow rand\{2, \cdots, n-2\}\\ \text{if } a^{\frac{n-1}{2}} \times \left(\frac{a}{n}\right) \ \text{mod} \ n \neq 1\\ \text{return 0}\\ \text{return 1} \end{split}
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References

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