Computer Science & Engineering Department IIT Kharagpur Computational Number Theory: CS60094 Lecture IV

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1 Basic Properties of Integers III

1.1 Quadratic Residues

Higher order congruence are more difficult to handle. We shall study a congruence of the form $x^2 \equiv a \pmod{n}$ for odd integer n > 1. Let us consider a general quadratic congruence $ax^2 + bx + c \equiv 0 \pmod{p}$, where p is an odd¹ prime and $a \in \mathbb{Z}_p^*$. It is clear that the gcd(4a, p) = 1. So we write the congruence as

$$4a(ax^2 + bx + c) \equiv 0 \pmod{p}$$

$$\Rightarrow (2ax)^2 + 2 \cdot 2ax \cdot b + b^2 - (b^2 - 4ac) \equiv 0 \pmod{p}$$

$$\Rightarrow (2ax + b)^2 \equiv (b^2 - 4ac) \pmod{p}.$$

If we substitute y for 2ax + b and d for $b^2 - 4ac$, we get $y^2 \equiv d \pmod{p}$. If $x \equiv x_0 \pmod{p}$ is a solution of the original congruence, then $y \equiv 2ax_0 + b \pmod{p}$ is a solution of the transformed congruence. Again if $y \equiv y_0 \pmod{p}$ is a solution of the transformed congruence, then $2ax + b \equiv y_0 \pmod{p}$, i.e. $2ax_0 \equiv y_0 - b \pmod{p}$. The solution of this linear congruence, which always exists as gcd(2a, p) = 1, gives the solution of the original congruence. In general we are interested about odd positive integer n.

<u>Definition 1:</u> Let n be an odd positive integer. An integer a is called a quadratic residue modulo n, if gcd(a, n) = 1 (a mod n belongs to \mathbb{Z}_n^*) and there is an integer b such that $a \equiv b^2 \pmod{n}$, then $x \equiv b \pmod{n}$ is a solution of $x^2 \equiv a \pmod{n}$, and b is called a square root of a modulo n.

There are a's that are not relatively prime to n, but satisfies the congruence $b^2 \equiv a \pmod{n}$. As an example, $6 \equiv 9^2 \pmod{15}$. But 6 is not called *quadratic* residue modulo 15. Quadratic residue and quadratic non-residue are defined for elements of \mathbb{Z}_n^* . Some of these elements are quadratic residue and others are quadratic non-residue.

Example 1. We take n = 13, a prime number.

There are two square roots of all the perfect squares modulo 13. There will always be at least two square roots of 1 modulo n. One is 1 and the other is n-1 as $(n-1)^2 \equiv 1 \pmod{n}$. This also tells us that there will be at least two square roots of any perfect squares modulo n.

Note that half of the elements of \mathbb{Z}_{13}^* are perfect squares modulo 13. They are the *quadratic residues* and the remaining half are the *quadratic non-residues*. The collection of *quadratic residues* forms a subgroup of \mathbb{Z}_n^* .

Example 2. In the second example let us consider a composite number, n = 15.

¹The case of p = 2 is simple as $\mathbb{Z}_2 = \{0, 1\}$. The coefficients a, b can be either 0 or 1. So $x^2 + x + 1 \equiv 0 \pmod{2}$ cannot have any solution, but $x^2 + x \equiv 0 \pmod{2}$ has two solutions.

The elements of $\mathbb{Z}_{15}^* = \{1, 2, 4, 7, 8, 11, 13, 14\}.$

$$1 \equiv 1^{2} \equiv 4^{2} \equiv 11^{2} \equiv 14^{2} \pmod{15}, 4 \equiv 2^{2} \equiv 7^{2} \equiv 8^{2} \equiv 13^{2} \equiv \pmod{15},$$

We observe that only prefect squares modulo 15 are 1 and 4. But each has four square-roots modulo 15. 1 and 4 are the two quadratic residues modulo 15. The quadratic non-residues are 2, 7, 8, 11, 13, 14.

Example 3. In the third example we take n = 9, a power of a prime, where $\overline{\mathbb{Z}_{q}^{*}} = \{1, 2, 4, 5, 7, 8\}.$

In this case also there are only two square roots of 1 modulo 9. So there are two square roots of all other perfect squares (quadratic residues) modulo 9. Half of the elements are quadratic residues.

<u>Definition 2:</u> Let $n, m \in \mathbb{Z}$ and n > 0,

$$(\mathbb{Z}_n^*)^m = \{a^m : a \in \mathbb{Z}_n^*\},\$$

the collection of the m^{th} powers of the elements of \mathbb{Z}_n^* .

It is not difficult to prove the following facts

- 1. $(\mathbb{Z}_n^*)^m$ is a subgroup of \mathbb{Z}_n^* .
- 2. Let $a \in \mathbb{Z}_n^*$ and let $l, m \in \mathbb{Z}$ so that l and m are relatively prime. If $a^l \in (\mathbb{Z}_n^*)^m$, then $a \in (\mathbb{Z}_n^*)^m$.

1.1.1 Quadratic Residue Modulo Odd Prime

We prove following interesting results related to any odd prime p.

Proposition 1. For any odd prime p, and $q \in \mathbb{Z}_p^*$, $q^2 \equiv 1 \pmod{p}$ if and only if $q \equiv 1 \text{ or } q \equiv p - 1 \equiv -1 \pmod{p}.$

Proof: If q = 1 or q = p - 1, then $1^2 \equiv 1 \pmod{p}$ and $(p - 1)^2 \equiv p^2 - 2p + 1 \equiv 1 \pmod{p}$ $1 \pmod{p}$.

In the other direction, let $q^2 \equiv 1 \pmod{p}$, so $p|(q^2 - 1)$. But then p is prime, so p|q-1 or p|q+1. But $q \in \mathbb{Z}_p^*$, we have either q-1=0 or q+1=p. QED.

<u>Proposition 2.</u> If $q \in (\mathbb{Z}_p^*)^2$, where p is an odd prime, then q has exactly two

square roots in \mathbb{Z}_p^* . **Proof:** Let $q \equiv a^2 \pmod{p}$ and also $q \equiv b^2 \pmod{p}$. So we have $a^2 \equiv b^2 \pmod{p}$. p. We multiply both sides by $(b^{-1})^2 \pmod{p}$ and get $(ab^{-1})^2 \equiv 1 \pmod{p}$. By the previous proposition we have $ab^{-1} \equiv 1 \pmod{p}$ or $ab^{-1} \equiv p - 1 \pmod{p}$. So $a \equiv b \pmod{p}$ $a \equiv -b \pmod{p}$. So in \mathbb{Z}_p^* , $a = \pm b$ i.e. there are exactly two square roots. QED.

<u>Proposition 3.</u> For any odd prime p, the size of $(\mathbb{Z}_p^*)^2$ is $\frac{p-1}{2}$. **Proof:** We define the map, $sq: \mathbb{Z}_p^* \to (\mathbb{Z}_p^*)^2, a \mapsto a^2$. Every image has two distinct preimages, so $2|(\mathbb{Z}_p^*)^2| = |\mathbb{Z}_p^*| = p - 1$. QED.

Example 4. Let p = 11, $(\mathbb{Z}_{11}^*)^2 = \{1^2 = 1, 2^2 = 4, 3^2 = 9, 4^2 = 5, 5^2 = 3\}$, and $\mathbb{Z}_{11}^* = \{\pm 1, \pm 2, \pm 3, \pm 4, \pm 5\}.$

If p is an odd prime, half of the elements of \mathbb{Z}_p^* are quadratic residue and other half are quadratic non-residue.

<u>Theorem 4.</u> (Euler's Criterion) Let p be an odd prime and $a \in \mathbb{Z}_p^*$.

- 1. $a^{(p-1)/2} \equiv \pm 1 \pmod{p}$,
- 2. If $a \in (\mathbb{Z}_p^*)^2$ then $a^{(p-1)/2} \equiv 1 \pmod{p}$,

3. If $a \notin (\mathbb{Z}_p^*)^2$ then $a^{(p-1)/2} \equiv -1 \pmod{p}$,

Proof:

- 1. Let $b \equiv a^{(p-1)/2} \pmod{p}$, so $b^2 \equiv a^{p-1} \equiv 1 \pmod{p}$, by Euler's theorem. But we know that if $b^2 \equiv 1 \pmod{p}$, then $b \equiv 1, p - 1 \pmod{p}$.
- 2. $a \equiv b^2 \pmod{p}$. So $a^{(p-1)/2} \equiv b^{p-1} \equiv 1 \pmod{p}$.
- 3. $a \in \mathbb{Z}_p^* \setminus (\mathbb{Z}_p^*)^2$.

We claim that for each $b \in \mathbb{Z}_p^*$ there is a $c \in \mathbb{Z}_p^*$ so that $bc \equiv a \pmod{p}$ and $b \neq c$.

If b = c, then $a \equiv b^2 \pmod{p}$ and $a \in (\mathbb{Z}_p^*)^2$. Our $c = b^{-1}a$ and it is unique.

So we have the product of all elements of \mathbb{Z}_p^* ,

$$\prod_{b,c\in\mathbb{Z}_p^*} (b\times_p c) = a^{(p-1)/2}.$$

We also claim that for each $b \in \mathbb{Z}_p^*$ there is a $c \in \mathbb{Z}_p^*$ so that $bc \equiv 1 \pmod{p}$. We know that there are only two elements in \mathbb{Z}_p^* whose square is 1. They are 1 and p-1. Any other b, c whose product is 1 are distinct. So we have another product of all elements of \mathbb{Z}_p^* ,

$$\left(\prod_{a\in\mathbb{Z}_p^*}a\right) = \left(1\times_p(-1)\times\prod_{b\times_p c=1, b\neq c}b\times_p c\right) = -1$$

Hence the result.

QED.

Example 5. We consider p = 13 and $5 \in \mathbb{Z}_{13}^* \setminus (\mathbb{Z}_{13}^*)^2$.

$$5 = 1 \times_{13} 5 = 2 \times_{13} 9 = 3 \times_{13} 6 = 4 \times_{13} 11 = 7 \times_{13} 10 = 8 \times_{13} 12.$$

Also

$$1 = 2 \times_{13} 7 = 3 \times_{13} 9 = 4 \times_{13} 10 = 5 \times_{13} 8 = 6 \times_{13} 11.$$

So we have $5^{\frac{13-1}{2}} \equiv 1 \times_{13} \times_{13} (-1) \times_{13} 1^{\frac{13-3}{2}} \equiv -1 \pmod{13}$. The conclusion of the Euler's criterion is

$$a \in (\mathbb{Z}_p^*)^2$$
 if and only if $a^{(p-1)/2} \equiv 1$.

We have a byproduct of our earlier proof.

<u>Theorem 5.</u> (Wilson's Theorem) If p is a prime, then $(p-1)! \equiv -1 \pmod{p}$. **Proof:** We have already proved that for an odd prime $\prod_{a \in \mathbb{Z}_p^*} a = -1$. This is also true for 2 where 1 is same as $-1 \mod 2$. QED.

Proposition 6. (Converse of Wilson's Theorem) If n is a positive integer greater than 1 and $(n-1)! \equiv -1 \pmod{n}$, then n is prime.

Proof: If n is not a prime, then n = ab, where 1 < a, b < n. From the given condition we see that n|[(n-1)!+1]. So a|[(n-1)!+1]. But then a|(n-1)! implies that a|1 - a contradiction. QED.

<u>Proposition 7.</u> Let p be an odd prime and $a, b \in \mathbb{Z}_p^*$. If none of a, b are in $(\mathbb{Z}_p^*)^2$, then $ab \in (\mathbb{Z}_p^*)^2$.

Proof: We have

$$(a \times_p b)^{(p-1)/2} \equiv a^{(p-1)/2} \times_p b^{(p-1)/2} \equiv -1 \times_p -1 \equiv 1 \pmod{p}.$$

So by the Euler's criterion $ab \in (\mathbb{Z}_p^*)^2$.

QED.

1.1.2 Quadratic Residue Modulo Power of Odd Prime

Let p be an odd prime and k > 0 be an integer. We are interested about the solution of $x^2 \equiv a \pmod{p^k}$ in $\mathbb{Z}_{p^k}^*$. We have already seen characterisation in case of k = 1. There are similar theorems for k > 1.

<u>Proposition 8.</u> Let p be any odd prime and let k be any positive integer. For all $q \in \mathbb{Z}_{p^k}^*$, $q^2 \equiv 1 \pmod{p^k}$ if and only if q = 1 or $q = p^k - 1 \equiv -1 \pmod{p^k}$. **Proof:** If q = 1 or $q = p^k - 1$, then $1^2 \equiv 1 \pmod{p^k}$ and $(p^k - 1)^2 \equiv p^{2k} - 2p^k + 1 \equiv 1 \pmod{p^k}$.

In the other direction, let $q^2 \equiv 1 \pmod{p^k}$, so $p^k | (q^2 - 1)$, implies p | (q - 1)(q + 1). p is prime, so p | q - 1 or p | q + 1. But p cannot divide both q - 1 as well as q + 1; otherwise p divides (q + 1) - (q - 1) = 2. But that is impossible as p is an odd prime. So p^k divides either q - 1 or q + 1. But then $q \in \mathbb{Z}_{p^k}^*$. So if $p^k | (q - 1)$, then q - 1 = 0, and if $p^k | (q + 1)$, then $q + 1 = p^k$ i.e. $q = p^k - 1$. QED.

Following sequence of propositions are similar to the case of k = 1. We leave them as exercise.

<u>Proposition 9.</u> For any odd prime p and positive integer k, if $q \in (\mathbb{Z}_{p^k}^*)^2$, then q has exactly two square roots in $\mathbb{Z}_{p^k}^*$.

<u>Proposition 10.</u> For any odd prime p and a positive integer k, the size of $(\mathbb{Z}_{p^k}^*)^2$ is $\frac{\phi(p^k)}{2}$.

Proposition 11. (Generalisation of Euler's Theorem)

Let p be an odd prime, k be a positive integer and $a \in \mathbb{Z}_{p^k}^*$.

- 1. $a^{\phi(p^k)/2} \equiv \pm 1 \pmod{p^k}$,
- 2. If $a \in (\mathbb{Z}_{p^k}^*)^2$ then $a^{\phi(p^k)/2} \equiv 1 \pmod{p^k}$,
- 3. If $a \notin (\mathbb{Z}_{p^k}^*)^2$ then $a^{\phi(p^k)/2} \equiv -1 \pmod{p^k}$,

Proposition 12. (Generalisation of Wilson's Theorem)

If p is an odd prime and k is a positive integer, then $\prod_{a \in \mathbb{Z}_{p^k}^*} a \equiv -1 \pmod{p^k}$. <u>Proposition 13.</u> Let p be an odd prime and k be a positive integer and $a, b \in \mathbb{Z}_{p^k}^* \setminus (\mathbb{Z}_{p^k}^*)^2$, then $ab \in (\mathbb{Z}_{p^k}^*)^2$.

Finally we have the following interesting proposition.

<u>Proposition 14.</u> If p is an odd prime, k is a positive integer and $a \in \mathbb{Z}_{p^k}^*$, then \overline{a} is a quadratic residue modulo p if and only if it is a quadratic residue modulo p^k .

Proof: Let *a* be a quadratic residue modulo p^k i.e. $gcd(a, p^k) = 1$ and there is an integer *b* so that $a \equiv b^2 \pmod{p^k}$. So we have gcd(a, p) = 1 and $a \equiv b^2 \pmod{p}$. So *a* is quadratic residue (perfect square) modulo *p*.

Let *a* is not a quadratic residue modulo p^k . If p|a, then *a* is not a quadratic residue modulo *p*. So we assume that $p \not|a$. Using the generalised Euler's criterion we have $a^{\phi(p^k)/2} \equiv -1 \pmod{p^k}$. This implies that $a^{\phi(p^k)/2} \equiv -1 \pmod{p}$. We use the Fermat's little theorem

$$a \equiv a^p \equiv (a^p)^p \equiv \dots \equiv a^{p^{\kappa-1}} \pmod{p}.$$

By substituting we get,

$$-1 \equiv a^{\phi(p^k)/2} \equiv a^{p^{k-1}(p-1)/2} \equiv (a^{p^{k-1}})^{(p-1)/2} \equiv a^{(p-1)/2} \pmod{p}.$$

So a is a quadratic non-residue modulo p.

QED.

1.1.3 Quadratic Residue Modulo n

Now we consider the general case of odd n, a product of odd primes.

<u>Proposition 15.</u> Let n be an odd integer greater than 1. The prime decomposition of $n = p_1^{e_1} \cdots p_k^{e_k}$. If $a \in \mathbb{Z}_n^*$ is a perfect square, then a has 2^k square-roots.

Example 6. Let us look at Example (1.1) where $n = 15 = 3^1 \times 5^1$. The Chinese Remainder Map is $f : \mathbb{Z}_{15} \to \mathbb{Z}_3 \times \mathbb{Z}_5$ is $n \mapsto (n \mod 3, n \mod 5)$. We know that

there are two perfect squares in \mathbb{Z}_{15} and they are 1 and 4. We have f(1) = (1, 1)and f(4) = (1, 4). There are two square roots of 1 modulo 3, and they are 1 and 2. Similarly there are two square roots of 1 modulo 5, they are 1 and 4. And two square roots of 4 modulo 5, they are 2 and 3.

So $f(1) = (1, 1) = (1^2 \mod 3, 1^2 \mod 5) = (1^2 \mod 3, 4^2 \mod 5) = (2^2 \mod 3, 1^2 \mod 5) = (2^2 \mod 3, 4^2 \mod 5)$. If $b^2 \equiv 1 \pmod{15}$, then $f(b^2) = (b^2 \mod 3, b^2 \mod 5) = ((b \mod 3)^2 \mod 3, (b \mod 5)^2 \mod 5)$. As $f(b) = (b \mod 3, b \mod 5)$, the values of b are $f^{-1}(1, 1) = 1$, $f^{-1}(1, 4) = 4$, $f^{-1}(2, 1) = 11$ and $f^{-1}(2, 4) = 14$. So the square roots of 1 modulo 15 are 1, 4, 11, 14. **Proof:** We consider the Chinese Remainder Map,

$$f: \mathbb{Z}_n \to \mathbb{Z}_{p_1^{e_1}} \times \cdots \times \mathbb{Z}_{p_h^{e_k}}$$

We know that the restriction of f to \mathbb{Z}_n^* is also a bijection.

$$f:\mathbb{Z}_n^*\to\mathbb{Z}_{p_1^{e_1}}^*\times\cdots\times\mathbb{Z}_{p_k^{e_k}}^*.$$

Let $a \in (\mathbb{Z}_n^*)^2$, a perfect square modulo n i.e. $a \equiv b^2 \pmod{n}$, for some $b \in \mathbb{Z}_n^*$. We have $f(a) = (a \mod p_1^{e_1}, \cdots, a \mod p_k^{e_k}) = (a_1, \cdots, a_k) \in \mathbb{Z}_{p_1^{e_1}}^{*} \times \cdots \times \mathbb{Z}_{p_k^{e_k}}^{*}$. Let $f(b) = (b \mod p_1^{e_1}, \cdots, b \mod p_k^{e_k}) = (b_1, \cdots, b_k) \in \mathbb{Z}_{p_1^{e_1}}^{*} \times \cdots \times \mathbb{Z}_{p_k^{e_k}}^{*}$.

Note that in a Chinese remainder map, $f : \mathbb{Z}_N \to \mathbb{Z}_{n_1} \times \cdots \times \mathbb{Z}_{n_k}$, where $\{n_i\}_{i=1}^k$ are pairwise relatively prime, and $N = \prod_{i=1}^k n_i$, $f(xy) = ((xy) \mod n_1, \cdots, (xy) \mod n_k) = ((x \mod n_1)(y \mod n_1) \mod n_1, \cdots, (x \mod n_k)(y \mod n_k) \mod n_k) \mod n_k) = (x_1y_1, \cdots, x_ky_k).$

So we have

$$(a_1, \cdots a_k) = f(a) = f(b^2) = (b_1^2, \cdots, b_k^2).$$

So we have perfect squares $a_i \equiv b_i^2 \pmod{p_i^{e_i}}$, for all $i = 1, \dots, k$.

On the other hand, if we have perfect square, $a_i \equiv b_i^2 \pmod{p_i^{e_i}}$, for all $i = 1, \dots, k$, then let us call $f^{-1}(b_1, \dots, b_k) = b \in \mathbb{Z}_n^*$ (the restriction of f is also a bijection). So we have

$$f(b^2) = (b^2 \mod p_1^{e_1}, \cdots, b^2 \mod p_k^{e_k}) = (b_1^2, \cdots, b_k^2) = (a_1, \cdots, a_k) = f(a).$$

As f is a bijection, $a \equiv b^2 \pmod{n}$ i.e. a is a perfect square in \mathbb{Z}_n^* . This shows that

$$a \in (\mathbb{Z}_n^*)^2$$
 if and only if $a_i \in (\mathbb{Z}_{p_1^{e_i}}^{*_{e_i}})^2$, for $i = 1, \cdots, k$.

Each perfect square in $\mathbb{Z}_{p_1^{e_i}}^{*}$ has two square roots, so *a* has 2^k square roots.

This gives us the size of $(\mathbb{Z}_n^*)^2$.

$$(\mathbb{Z}_n^*)^2| = \prod_{i=1}^k |(\mathbb{Z}_{p_1^{e_i}}^*)^2| = \prod_{i=1}^k \phi(p_1^{e_i})/2 = \phi(n)/2^k$$

We formally conclude that any element $a \equiv b^2 \in (\mathbb{Z}_n^*)^2$ has 2^k square roots. Let $a \equiv b^2 \pmod{n}$ and $a \equiv c^2 \pmod{n}$. So we have $b^2 \equiv c^2 \pmod{n}$. This amounts to saying that $b_i^2 \equiv c_i^2 \pmod{n_i}$, for $i = 1, \dots, k$. So each $b_i \equiv \pm c_i \pmod{n_i}$ and there are altogether 2^k possibilities. QED.

1.1.4 Testing of Quadratic Residuosity

We wish to test whether an integer a is quadratic residue modulo n. If $gcd(a, n) \neq 1$, then a by definition is not a quadratic residue modulo n. So we assume that a and n are relatively prime. We consider the following three cases:

1. *n* is an odd prime: We compute the value of $a^{\frac{p-1}{2}} \mod n$. This can be done using repeated squaring algorithm given below

It computes $a^e \mod n$ where $a \in \mathbb{Z}_n$ and e is a positive integer. Let the binary representation of $e = (e_{k-1}e_{k-2}\cdots e_1e_0)$.

$$\begin{array}{l} \operatorname{modExpN}(a,e,n) \\ exp \leftarrow 1 \\ s \leftarrow a \bmod n \\ \texttt{while } e \geq 1 \\ \texttt{if } (e \bmod 2) = 1 \texttt{ then } exp \leftarrow (exp \times s) \bmod n \\ s \leftarrow s^2 \bmod n \\ e \leftarrow e \div 2 \\ \texttt{return } exp \end{array}$$

 $k = \lceil \log_2 e \rceil$, so the loop is executed k times with k squaring and $\leq k$ multiplications over \mathbb{Z}_n . So the running time is $O(\log e(\log n)^2)$. If 1 < e < n, then it is $O(\log n)^3$.

We shall see afterward that there is better method for testing *quadratic* residuosity for an odd prime.

- 2. $n = p^k$ where p is an odd prime: We have already proved that a is a quadratic residue modulo p^k if and only if a is a quadratic residue modulo p. So this can also be done efficiently.
- 3. n is an odd integer: If the prime factorisation of n is known, then we can use the previous method to determine whether a is a quadratic residue modulo p for every prime factor of n. Then using this fact we can conclude about the quadratic residuocity of a modulo n (Chinese remainder theorem). But if the factorisation is not given, there is no efficient algorithm known to test quadratic residuocity. Factorisation is believed to be a hard problem.

We shall see that the computation of *Jacobi symbol*, for which efficient algorithm is known, gives partial answer.

If it is known that a is a quadratic residue modulo p, an odd prime, it is necessary to find one b such that $b^2 \equiv a \pmod{p}$. We shall address this problem afterward.

1.1.5 Square Roots of p-1

Following theorem characterises the odd primes p such that p-1 or -1 is a quadratic residue modulo p. This has some interesting applications.

<u>Proposition 16.</u> Let p be an odd prime. $p - 1 \in (\mathbb{Z}_p^*)^2$ if and only if $p \equiv 1 \pmod{4}$ i.e. p = 4k + 1.

Proof: By the Euler's criterion, p-1 is a quadratic residue modulo p if and only if $(p-1)^{(p-1)/2} \equiv 1 \pmod{p}$. If p is of the form 4k+1, then (p-1)/2 = 2k, an even number. So

$$(p-1)^{2k} \equiv (-1)^{2k} \equiv 1 \pmod{p}.$$

If $p \equiv 3 \pmod{p}$, then p = 4k + 3 and (p - 1)/2 = 2k + 1, an odd number. So,

$$(p-1)^{2k+1} \equiv (-1)^{2k+1} \equiv -1 \pmod{p}.$$

QED.

Proposition 17. There are infinitely many primes $p \equiv 1 \pmod{4}$.

Proof: Let there be finite number of such primes, p_1, \dots, p_k , and let $n = 4m^2 + 1$, where $m = p_1 \dots p_k$. Let p be a prime factor of n. Clearly p is not equal to any one of p_1, \dots, p_k . We have $(2m)^2 \equiv -1 \pmod{p}$. So -1 is a quadratic residue of p and by our previous theorem, $p \equiv 1 \pmod{4}$. This contradicts our assumption. QED.

<u>Proposition 18.</u> (Thue's Lemma) Let p be a prime and a is an integer such that $p \not| a$. There exists two integers x and y, such that (i) $0 < |x|, |y| < \sqrt{p}$, and (ii) $ax \equiv y \pmod{p}$.

<u>Example 7.</u> Let n = 13 and a = 7. We have $7x = y \pmod{13}$. (2,1) is a solution of the congruence satisfying $0 < 1, 2 < \sqrt{13}$.

Proof: Let

$$A = \{au - v : u, v \in \mathbb{Z} \land 0 \le u, v \le |\sqrt{p}|\}.$$

Clearly there are $\lfloor \sqrt{p} \rfloor + 1$ integers in the interval. So, the number of ordered pairs (u, v) corresponding to the elements of A is greater than p. By the pigeonhole principle there are two distinct ordered pairs (u_1, v_1) and (u_2, v_2) such that $au_1 - v_1 \equiv au_2 - v_2 \pmod{p}$. So we have $a(u_1 - u_2) \equiv (v_1 - v_2) \pmod{p}$. This gives a solution of $ax \equiv y \pmod{p}$, where $x = u_1 - u_2$ and $y = v_1 - v_2$.

Both $|x|, |y| < \sqrt{p}$ (as a prime cannot be a perfect square). If one of x or y is 0, the congruence $ax \equiv y \pmod{p}$ implies that the other one will also be 0. But both $x = u_1 - u_2$ and $y = v_1 - v_2$ cannot be 0 as the ordered pairs are distinct. So both x and y are non-zero. QED.

Theorem 19. (Fermat)

An odd prime p is expressible as sum of two squares if and only if $p \equiv 1 \pmod{4}$.

Proof: If $p = a^2 + b^2$, then one of a or b is odd and the other one is even. We assume that a = 2c and b = 2d + 1. So $a^2 \equiv 0 \pmod{4}$ and $b^2 \equiv 1 \pmod{4}$, implies that $p = a^2 + b^2 \equiv 1 \pmod{4}$.

If $p \equiv 1 \pmod{4}$, -1 is a quadratic residue modulo p. So we have an integer a so that $a^2 \equiv -1 \pmod{p}$, where gcd(p, a) = 1.

At this point we invoke the Thue's lemma. There is a solution (X_0, Y_0) of $ax \equiv y \pmod{p}$ such that $0 < |X_0|, |Y_0| < \sqrt{p}$. So we have

$$\begin{array}{rcl} aX_0 &\equiv & Y_0 (\bmod \ p), \\ (aX_0)^2 &\equiv & Y_0^2 (\bmod \ p), \\ a^2X_0^2 &\equiv & Y_0^2 (\bmod \ p), \\ -X_0^2 &\equiv & Y_0^2 (\bmod \ p), \ a^2 \equiv -1 (\bmod \ p) \\ X_0^2 + Y_0^2 &\equiv & 0 (\bmod \ p). \end{array}$$

As $p|(X_0^2 + Y_0^2), X_0^2 + Y_0^2 = kp$, where $k \ge 1$. But $0 < |X_0|, |Y_0| < \sqrt{p}$. So, $X_0^2 + Y_0^2 < 2p$, implies that k = 1. QED.

1.1.6 Computation of Fermat's Two Square

The proof of Fermat's two-square theorem depends on Thue's Lemma and the square-root of -1 modulo the prime p which is of the form 4k + 1. The proof of Theu's Lemma depends on pigeon-hole principle, and in that form it is not computable.

But we can us the extended GCD algorithm to compute (X_0, Y_0) as a solution of $ax \equiv y \pmod{p}$. Consider the following is the sequence of remainders r_i , $i = 0, \dots, k, k + 1$, and Bezout's coefficients $x_i, y_i, i = 0, \dots, k$, computed by the extended GCD algorithm.

$$(r_0 = p, 1, 0), (r_1 = a, 0, 1), \cdots, (r_i, x_i, y_i) \cdots (r_k, x_k, y_k),$$

where $r_{k+1} = 0$. As $gcd(p, a) = r_k = 1 = px_k + ay_k$. In general $r_i = px_i + ay_i$. The computation steps are as usual.

$$\begin{aligned} r_{i-1} &= r_i q_i + r_{i+1}, \\ x_{i+1} &= x_{i-1} - x_i q_i, \\ y_{i+1} &= y_{i-1} - y_i q_i, \quad i = 1, \cdots, k. \end{aligned}$$

We continue the computation as long as $r_i \ge \sqrt{p}$ and stop at $r_i < \sqrt{p}$. This is possible as $r_0 = p > 0 = r_{k+1}$. On termination we set $Y_0 = r_i$ and $X_0 = y_i$. We know that $r_i = x_i p + y_i a$ i.e. $Y_0 = x_i p + X_0 a$, so we have $aX_0 \equiv Y_0 \pmod{n}$ where $0 < Y_0 < \sqrt{p}$.

We prove that $0 < |X_0| < \sqrt{p}$. Consider following two equations.

$$r_{i-1} = px_{i-1} + ay_{i-1}, (1)$$

$$r_i = px_i + ay_i. (2)$$

 $(1) \times y_i - (2) \times y_{i-1}$ gives us,

$$p(x_{i-1}y_i - x_iy_{i-1}) = y_ir_{i-1} - y_{i-1}r_i$$

$$p \times (-1)^{i-1} = y_ir_{i-1} - y_{i-1}r_i$$

$$p = |y_ir_{i-1} - y_{i-1}r_i|$$

$$= |y_i|r_{i-1} + |y_{i-1}|r_i, \ y_iy_{i-1} \le 0,$$

$$y_i \text{ and } y_{i-1} \text{ have opposite signs}$$

$$\ge |y_i|r_{i-1}.$$

We can prove by induction that $x_{i-1}y_i - x_iy_{i-1} = (-1)^{i-1}$. So,

$$|X_0| = |y_i| \le \frac{p}{r_{i-1}} < \frac{p}{r_i} = \frac{p}{\sqrt{p}} = \sqrt{p}.$$

<u>Example 8.</u> Let p = 83, a = 34, $\lfloor \sqrt{83} \rfloor = 9$ So it Following is the table for extended GCD computation.

i	r_i	x_i	y_i	q_i
0	83	1	0	
1	34	0	1	2
2	15	1	-2	2
3	4	-2	5	

So we have $Y_0 = r_3 = 4$ and $X_0 = y_3 = 5$ such that $0 < X_0, Y_0 < \sqrt{83}$. We have $83 \times (-2) + 34 \times 5 = 4$. A solution of $34x \equiv y \pmod{83}$.

Now we turn our attention to the computation of the square-root of -1 modulo prime $p \equiv 1 \pmod{4}$. We want to compute an element $a \in \mathbb{Z}_p^*$ so that $a^2 \equiv -1 \pmod{p}$. If we can find an element $b \in \mathbb{Z}_p^* \setminus (\mathbb{Z}_p^*)^2$, we may take $a = b^{\frac{p-1}{4}}$, as $a^2 \equiv (b^{\frac{p-1}{4}})^2 = b^{\frac{p-1}{2}} \equiv -1 \pmod{p}$ (Euler's criterion).

We know that half of the elements of \mathbb{Z}_p^* are quadratic non-residue. So we can use the following randomised algorithm.

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\begin{array}{c} \textit{sqrt-1}(p) \\ \texttt{do} \\ b \leftarrow \textit{rand}\{1, \cdots, p-1\} \\ a \leftarrow b^{(p-1)/4} \\ \texttt{while} \ (a^2 \bmod p \neq p-1) \\ \texttt{return} \ a \end{array}
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The probability of picking a quadratic non-residue is $\frac{1}{2}$. So the expected number of times the loop is executed is 2. The probability that the algorithm has not found a quadratic non-residue after k iterations is $1/2^k$. The algorithm when terminates gives the correct a. But its running time is a random variable that is bounded. This type of algorithms are known as Las Vegas algorithm.

Modular exponentiation is the costly part of computation and we have seen that it takes $O(\log n)^3$) time.

We are now ready to express a prime $p \equiv 1 \pmod{4}$ as a sum of two squares in the following way. The input is p, where $p \equiv 1 \pmod{4}$

- 1. Find $a \in \mathbb{Z}_p^*$ such that $a^2 \equiv -1 \pmod{p}$.
- 2. Take a and run the modified extended-GCD algorithm to compute (X_0, Y_0) .

<u>Example 9.</u> Let $p = 977 = 4 \times 244 + 1$. Take $19 \in \mathbb{Z}_{977}^* \setminus (\mathbb{Z}_{977}^*)^2$ so that $19^{(977-1)/4} \equiv 725 \pmod{977}$ and $725^2 \equiv 976 \pmod{977}$. We have $\sqrt{977} > 31$. The run of the extended GCD algorithm on (p, a) is as follows:

i	r_i	x_i	y_i	q_i
0	977	1	0	-
1	725	0	1	1
2	252	1	-1	2
3	221	-2	3	1
4	31	3	-4	• • •

At this point we stop computation where $Y_0 = 31$ and $X_0 = -4$. We express $977 = Y_0^2 + X_0^2 = 31^2 + 4^2$.

<u>Proposition 20.</u> A prime $p \equiv 1 \pmod{4}$ can be represented uniquely as a sum of two squares (ignoring sign and order).

Proof: Let $p = a^2 + b^2 = c^2 + d^2$. We rewrite it as

$$a^{2}d^{2} + b^{2}d^{2} - b^{2}c^{2} - b^{2}d^{2} = (a^{2} + b^{2})d^{2} - (c^{2} + d^{2})b^{2} = p(d^{2} - b^{2}) \equiv 0 \pmod{p}.$$

So we have $(ad)^2 - (bc)^2 \equiv 0 \pmod{p}$ i.e. $(ad + bc) \equiv 0 \pmod{p}$ or $(ad - bc) \equiv 0 \pmod{p}$. But we know that $a, b, c, d < \sqrt{p}$. So there are two possibilities either (i) ad - bc = 0, or (ii) ad + bc = p. The second condition gives us

$$p^{2} = (a^{2} + b^{2})(c^{2} + d^{2}) = (ad + bc)^{2} + (ac - bd)^{2} = p^{2} + (ac - bd)^{2}.$$

It is equivalent to ac = bd.

We have two conditions (i) ad - bc = 0, or (ii) ac = bd.

If we consider the first condition ad = bc. We know that gcd(a, b) = 1, as p is prime. So a|c i.e. c = ak and we get ad = kab implies that d = kb.

But then $p = c^2 + d^2 = k^2(a^2 + b^2)$ and p is prime. So k = 1, which implies that a = c and b = d.

Similarly from the other condition we get the same result. QED.

<u>Proposition 21.</u> If n be a positive integer written as $n = N^2 m$, where m is square free, then n can be represented as sum of two squares if and only if m does not contain a prime factor of the form 4k + 3.

Proof: Let *m* has no prime factor of the form 4k+3: if m = 1, then $n = N^2+0^2$. If $m = p_1 \cdots p_k$, where p_i is either 2 or any prime of the form 4k+1. So each $p_i = a_i^2 + b_i^2$. Given two such primes p_i and p_j we have

$$p_i p_j = (a_i^2 + b_i^2)(a_j^2 + b_j^2) = (a_i a_j + b_i b_j)^2 + (a_i b_j - b_i a_j)^2.$$

So by induction we can prove that m can be expressed as $a^2 + b^2$. And finally $n = (aN)^2 + (bN)^2$.

Let $n = N^2 m$ can be written as $a^2 + b^2$. If m = 1, there is nothing to prove. Let m > 1 and also let gcd(a, b) = d, a = dA, b = dB. We have

$$a^{2} + b^{2} = d^{2}(A^{2} + B^{2}) = n = N^{2}m.$$

m is square free so $d^2 | N^2$. So we have

$$A^2 + B^2 = \frac{N^2}{d^2}m = qp$$

where p is an odd prime factor of m. So,

$$A^2 + B^2 \equiv 0 \pmod{p}.$$

As gcd(A, B) = 1, either A or B is relatively prime to p. Otherwise p will divide both of them and the $gcd(A, B) \ge p$, a contradiction. Let A be relatively prime to p. So we have

$$AA' \equiv 1 \pmod{p}$$

So we have $(A^2 + B^2)(A')^2 \equiv 0 \pmod{p}$, implies that $(AA')^2 + (BA')^2 \equiv 1 + (BA')^2 \equiv 0 \pmod{p}$. So -1 is a quadratic residue of p implies that $p \equiv 1 \pmod{4}$. QED.

<u>Proposition 22.</u> A positive integer can be represented as a sum of two squares if and only if its prime factors of the form 4k + 3 occurs in even power.

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