

**Computer Science and Engineering**  
**IIIT Kalyani, West Bengal**

**Compilers Design Laboratory (CS 511)**  
**(Autumn: 2019 - 2020)**  
*3rd Year CSE: 5<sup>th</sup> Semester*

Assignment - 6

Marks: 10

Assignment Out: 11<sup>th</sup> October, 2019 Report on or before: 18<sup>th</sup> October, 2019

1. Consider the following *context-free grammar G*: the *non-terminals* are  $N = \{ AS, BE, D, DL, E, IOS, IS, LS, P, PE, RE, S, SL, TY, VL \}$ , the *terminals* are  $\Sigma = \{ + - * / == < > ( ) = ; \& | \sim \text{else ic fc id if int real do nop print read str then while } \}$ , the *start symbol* is  $P$  and the *production rules* are,

```
P → DL SL
DL → DL D | ε
D → VL : TY
TY → int | real
VL → VL id | id
SL → SL | S
S → AS | IS | LS | IOS | nop
AS → id = E
IS → if BE then SL else SL ;
LS → while BE do SL ;
IOS → print PE | read id
PE → E | str
BE → BE ' | ' BE | BE & BE | ~ BE | ( BE ) | RE
RE → E == E | E < E | E > E
E → E + E | E - E | E * E | E / E | - E | ( E ) | id | ic | fc
```

Most of the terminals have their usual meaning e.g. **id** is an *identifier*, **ic** is an integer constant, **fc** is a floating-point constant, **str** is a string of characters, **nop** for *no operation*, = is an assignment, == is equality, '|' is *logical 'or'*, & is *logical 'and'*, ~ is *'not'* etc. Both '|' and '&' are *left associative*. The precedence relations of logical operators are '| < '&' < '~'.

An *identifier* starts with an English alphabet followed by a sequence of alphabet or decimal digits. An integer constant is a sequence of decimal digits. a floating-point constant is a sequence of decimal digits with a *decimal dot* in it (.12, 12., 1.2).

A comment starts with '//' and is up to the end of line ('\n'). A character *string* is within a pair of quotes ("").

2. Write a *flex* specification of the scanner. The name of the *flex* specification file should be `<roll no>.1`. The command  

```
$ flex <roll no>.1
```

generate the C code for the scanner in `lex.yy.c`. Include the header file `<roll no>.tab.h++` in the definition section of your *flex* specification. This header file, created by *bison*, contains the token names, type of `yylval` etc.
3. Write *bison* specification for the grammar to generate a parser. You need to specify the precedence and associativity of *boolean* and *arithmetic* operators as these parts of the production rules makes the grammar *ambiguous*. The specification file name should be `<roll no>.y++` (this will generate C++ parser code). The command  

```
$ bison -d -v <roll no>.y++
```

generates three files  
`<roll no>.tab.c++`,  
`<roll no>.tab.h++`, and

<roll no>.output.

The .output file contains the description of the LALR(1) automaton.

4. The compiled code of <roll no>.tab.c and lex.yy.c gives the basic parser. Let the name of the executable file be myParser. Prepare a Makefile for the whole process.
5. To send the assignment for evaluation prepare a tar file with the name <roll no>.6.tar which includes three (3) or more files: Makfile, <roll no>.1, <roll no>.y++. Kindly do not put them under a subdirectory while preparing the tar file.
6. A sample input is

```
// This program computes factorial
n fact i : int
read n
i = 1
fact = 1
while i < n | i == n do
    fact = fact * i
    i = i + 1
; // Note the ;
print fact
```

7. You may print the syntax error message by defining yylineno in flex (%option yylineno) and using it with the yyerror() in bison. As an example -

```
// This program computes factorial
n fact i : int
read n
i = 1
fact = 1
while i n | i == n do
    fact = fact * i
    i = i + 1
; // Note the ;
print fact
```

syntax error at line no: 6