

**CS11001/CS11002**  
**Programming and Data**  
**Structures (PDS)**  

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**(Theory: 3-1-0)**

**Important Announcements**

- Next class on next Wednesday as usual
- First class test on Wednesday (6<sup>th</sup> Sep.)
  - 1 hour duration (6pm-7pm)
  - Please assemble half an hour early

# Recursions and Functions

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## What are recursions?

- Expressing an entity in terms of itself is called recursion.
  - But remember it is not a circular definition.
    - here a solution to a bigger problem is expressed in terms of solutions to smaller problems.
    - a general mathematical concept
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## Iterative vs Recursive

- The factorial of 0 is 1, and the factorial of any positive integer  $n$  is the product of all integers from 1 to  $n$ .
- The factorial of 0 is 1. The factorial of any positive integer  $n$  is the product of  $n$  and the factorial of  $(n-1)$ .
  - the first definition is iterative, while the second one is recursive.
  - The factorial of  $(n-1)$  is solved in the same way, until the factorial of 0 is asked for, which is 1.

## The C function

```
int factorial ( int n )
{
    if ( n < 0 ) return (-1); /* Error condition */
    if ( n == 0 ) return (1); /* Base case */
    return(n * factorial(n-1)); /* Recursive call */
}
```

## A recursive function

- A function that calls itself, or which calls another function, which calls the first one is called recursive function.
- The later definition can be extended to more than two functions.
- However there are two important criteria:
  - Each time a function call occurs, it must be closer, **in some sense** to the solution.
  - There must be a decision criteria for stopping the process or computation: this being called “the escape hatch”.

Careless recursive coding can lead to infinite loops!

```
■ int Infinite(void)
  {
    Infinite();
    return 1;
  }
```

This will lead to an infinite loop! More technically stack overflow.

## Illustration of recursion

```
#include<stdio.h>
void CountNum (int n)
{
    static int i = 1;
    printf("In the function, the value of n is:%d\n",n);
    printf("\n The depth of the call is %d\n",i++);
    if(n>1)
        CountNumber(n-1);
    printf("\n After recursive call, value of i=%d at
n=%d\n", i--,n);
}
```

## Illustration of recursion

```
void main()
{
    int num = 3;
    CountNum(num);
}
```

- The if statement is encountered and a recursive call is made.
- The static variable i indicates the depth of the recursive calls.

(Note a static variable is a special data type which is initialized only once.)

## Illustrations of recursion

- When the CountNum function is called with N=1, the if statement is false.
- The function call CountNum is skipped.
- Thus the next statement is executed, which prints the argument value.
- At this stage the closing brace is encountered.
- The function does not return control to the main, but to the function body in CountNum, which gave the last call. The last call was in the if statement, hence the subsequent print statement is executed with i=3.
- This continues, till the final return to main.

## Result

- In the function, the value of n is: 3
- The depth of the call is :1
- In the function, the value of n is: 2
- The depth of the call is :2
- In the function, the value of n is: 1
- The depth of the call is :3
- After recursive call, value of i is :3
- After recursive call, value of i=4 at n=1
- After recursive call, value of i=3 at n=2
- After recursive call, value of i=2 at n=3

## Recursive Binary Search

- *Searching* and *Sorting* are inherently recursive in nature
- Many times, they follow a *divide and conquer* policy
- *Binary Search* searches a pre-sorted array of numbers
- In every function call, size of array to be sorted is halved
- Then, the half of the array in which the value might exist is checked
- If there is a possibility of finding the number in this half, this half is kept and the other half is discarded
- If there is no possibility of finding the number in this half, this half is discarded and the other half is considered

## The Binary Search Code

```
#include<stdio.h>
#define MAX 20
int binsearch(int a[], int key, int low, int high)
{
    int pos, mid;
    mid=(low+high)/2;
    if(a[mid]==key) return(mid);
    else if(a[mid]>key) pos=binsearch(a,key,low,mid);
    else pos=binsearch(a,key,mid,high);
    return(pos);
}
```

## The Binary Search Code

```
main()
{
    int a[MAX];
    int m, i, n, key;
    printf("How many elements:\n");
    scanf("%d",&n);
    for(i=0;i<n;i++) scanf("%d",&a[i]);
    printf("Enter the element to be searched?");
    scanf("%d",&key);
    m=binsearch(a,key,0,n-1);
    printf("Location is %d\n",m);
}
```

## Printing Backwards Code

```
#include<stdio.h>
#define EOLN '\n'
int main()
{
    printf("Pls enter a line of text\n");
    reverse();
    return 0;
}
void reverse(void)
{
    char c;
    if((c=getchar())!=EOLN) {
        reverse();
    }
    putchar(c);
    return;
}
```



## Recursion vs Iteration

- Recursion makes writing of a program easy.
- However recursion comes with disadvantages.

- Consider the example:

```
int F ( int n ) {  
    if ( n == 0 ) return ( 0 );  
    if ( n == 1 ) return ( 1 );  
    return ( F(n-1)+F(n-2)); }  
}
```

## Recursion vs Iteration

- Computation of  $F_n$  by the iterative version (using simple loops) requires  $n-1$  additions and some additional overheads proportional to  $n$ .
- But what about the recursive version?
- Let  $S_n$  denote the number of additions performed:
  - If  $n = 25$ , we have  $S_n = 121392$ , whereas for  $n = 50$ , we have  $S_n = 20365011073$ .
  - Compare these figures with the very small numbers (respectively 24 and 49) of additions performed by the iterative method.
  - The reason for this poor performance of the recursive algorithm is that many  $F_i$ s are computed multiple times.

## Tail Recursion

- It is, therefore, often advisable to replace recursion by iteration.
- If some function makes only one recursive call and does nothing after the recursive call returns (except perhaps forwarding the value returned by the recursive call), then one calls this recursion a **tail recursion**.
- Tail recursions are easy to replace by loops: since no additional tasks are left after the call, no book-keeping need be performed, i.e., there is no harm if we simply replace the local variables and function arguments by the new values pertaining to the recursive call.
- This leads to an iterative version with the loop continuation condition dictated by the function arguments.

- The factorial and Fibonacci routines that we presented earlier are not tail-recursive.
- The factorial routine performs a multiplication after the recursive call returns, and so it feels the necessity to store the formal parameter  $n$ .
- With the following implementation this need is eliminated.
- Here we pass to the recursive function an accumulating product.

## A Tail Recursive Factorial function

```
■ int facrec ( int n , int prod )
  {
    if (n < 0) return (-1);
    if (n == 0) return (prod);
    return (facrec(n-1,n*prod));
  }
int factorial ( int n ) {
  return (facrec(n,1));
}
```

## An Iterative Factorial Function

## An Iterative Factorial Function

```
■ int faciter ( int n ) {  
    int prod;  
    if ( n < 0 ) return (-1);    prod = 1;  
        /* Corresponds to facrec(n,1) */  
    while ( n > 0 ) {  
/* Corresponds to the sequence of recursive calls */  
        prod *= n;    /* Second argument in the recursive call */  
        n = n - 1;    /* Change the formal parameter */    }  
    return (prod); }  
  
_____
```