

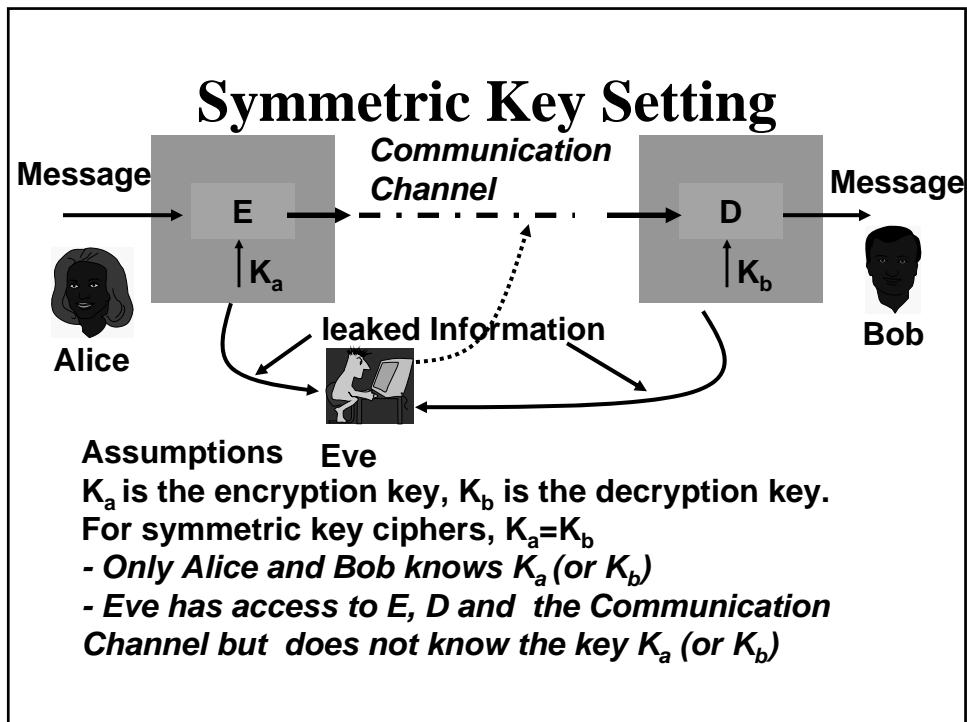
# Symmetric Key Ciphers

**Debdeep Mukhopadhyay**

**Assistant Professor  
Department of Computer Science and  
Engineering  
Indian Institute of Technology Kharagpur  
INDIA -721302**

## Objectives

- Definition of Symmetric Types of Symmetric Key ciphers
  - Modern Block Ciphers
- Components of a Modern Block Cipher
  - PBox (Permutation Box)
  - SBox (Substitution Box)
  - Swap
  - Properties of the Exclusive OR operation
- Diffusion and Confusion
- Types of Block Ciphers
- The DES Cipher



## Types of symmetric key ciphers

- Block Ciphers: Symmetric key ciphers, where a block of data is encrypted
- Stream Ciphers: Symmetric key ciphers, where block size=1

## Block Cipher

- A symmetric key modern cipher encrypts an  $n$  bit block of plaintext or decrypts an  $n$  bit block of ciphertext.
- Padding:
  - If the message has fewer than  $n$  bits, padding must be done to make it  $n$  bits.
  - If the message size is not a multiple of  $n$ , then it should be divided into  $n$  bit blocks and the last block should be padded.

## Components of a Modern Block Cipher

- Most important components:
  - PBox: It is a key-less fixed transposition cipher
  - SBox: It is a key-less fixed substitution cipher
- They are used to provide:
  - **Diffusion:** it hides the relationship between the ciphertext and the plaintext
  - **Confusion:** it hides the relationship between the ciphertext and the key

## Principle of Confusion and Diffusion

- The design principles of Block Cipher depends on these properties
- The S-Box is used to provide **confusion**, as it is dependent on the unknown key
- The P-Box is fixed, and there is no confusion due to it
- But it provides **diffusion**
- Properly combining these is necessary.

## Diffusion (P) Boxes

- Straight Boxes

**Example**  
**24x24 Box**

01	15	02	13	06	17	03	19	09	04	21	11
14	05	12	16	18	07	24	10	23	08	22	20

- Expansion Boxes

**Example**  
**12x24 Box**

01	03	02	01	06	17	03	07	09	04	09	11
02	05	12	04	06	07	12	10	11	08	10	08

- Compression Boxes

**Example**  
**24x12 Box**

01	15	02	13	06	17	03	19	09	04	21	11
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## SBox

An SBox (substitution box) is an  $m \times n$  substitution box, where  $m$  and  $n$  are not necessarily same.

Each output bit is a Boolean function of the inputs.

$$y_1 = f_1(x_1, x_2, \dots, x_n)$$

$$y_2 = f_2(x_1, x_2, \dots, x_n)$$

...

$$y_m = f_m(x_1, x_2, \dots, x_n)$$

## Non-linear SBox

$$y_1 = a_{11}x_1 \oplus a_{12}x_2 \oplus \dots \oplus a_{1n}x_n$$

$$y_2 = a_{21}x_1 \oplus a_{22}x_2 \oplus \dots \oplus a_{2n}x_n$$

...

$$y_m = a_{m1}x_1 \oplus a_{m2}x_2 \oplus \dots \oplus a_{mn}x_n$$

In a non-linear S-Box, each of the elements cannot be expressed as above.

Eg.

$$y_1 = (x_1)^3 + x_2, y_2 = (x_1)^2 + x_1x_2 + x_3$$

## Other Components

- Circular Shift:
  - It shifts each bit in an n-bit word k positions to the left. The leftmost k bits become the rightmost bits.
  - Invertible Transformation
- Swap:
  - A special type of shift operation where  $k=n/2$
- Other operations involve split and combine.
- An important component is exclusive-or operation

## Properties of Exor

Ex-or is a binary operator, which results in 1 when both the inputs have a different logic. Otherwise, it computes 0.

Symbol:  $\oplus$

Closure: Result of exoring two n bit numbers is also n bits.

Associativity: Allows to use more than one ' $\oplus$ 's in any order:

$$x \oplus (y \oplus z) = (x \oplus y) \oplus z$$

Commutativity:  $x \oplus y = y \oplus x$

Identity: The identity element is the n bit 0, represented by

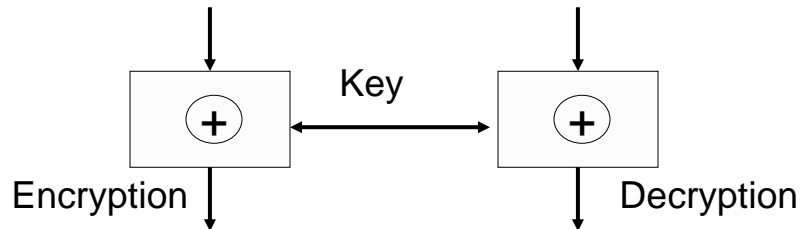
$$(00\dots 0) = 0^n$$

$$\text{Thus, } x \oplus 0^n = x$$

Inverse: Each word is the additive inverse of itself.

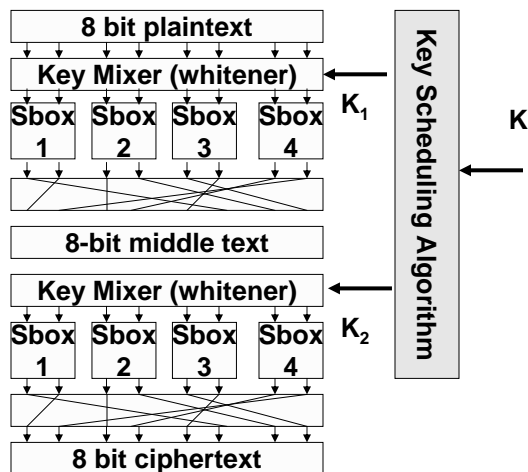
$$\text{Thus, } x \oplus x = 0^n$$

## Application of Ex-or

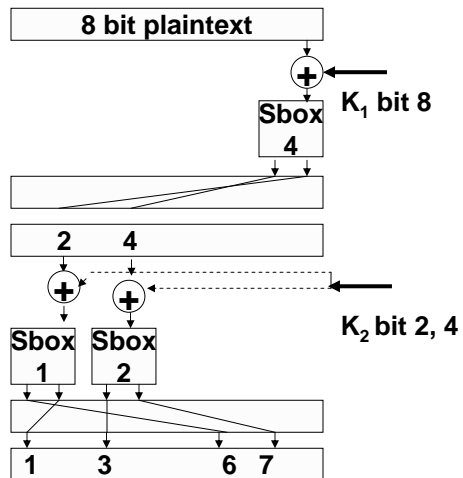


- The key is known to both the encryptor and decryptor and helps to recover the plaintext.

## A product cipher made of 2 rounds



## Diffusion and Confusion



## Practical Ciphers

- Large data blocks
- More S-Boxes
- More rounds
- These help to improve the diffusion and confusion in the cipher.



## Two classes of product ciphers

- Feistel Ciphers, example DES (Data Encryption Standard)
- Non-Feistel Ciphers (Substitution Permutation Networks), example AES (Advanced Encryption System)

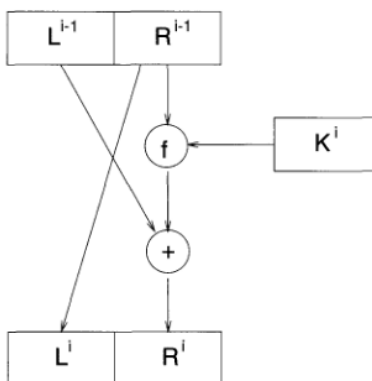
## Feistel Cipher

- **Feistel cipher** refers to a type of block cipher design, not a specific cipher
- Split plaintext block into left and right halves:  
Plaintext =  $(L_0, R_0)$
- For each round  $i=1,2,\dots,n$ , compute  
 $L_i = R_{i-1}$   
 $R_i = L_{i-1} \oplus F(R_{i-1}, K_i)$   
where  $f$  is **round function** and  $K_i$  is **subkey**
- Ciphertext =  $(L_n, R_n)$

## Feistel Permutation

- Decryption: Ciphertext =  $(L_n, R_n)$
- For each round  $i=n, n-1, \dots, 1$ , compute
$$R_{i-1} = L_i$$
$$L_{i-1} = R_i \oplus F(R_{i-1}, K_i)$$
where  $f$  is round function and  $K_i$  is subkey
- Plaintext =  $(L_0, R_0)$
- Formula “works” for any function  $F$
- But only secure for certain functions  $F$

## Encryption



Repeating/ Iterating this transformation we obtain the Feistel Cipher

## Non-Feistel Ciphers

- Composed of only invertible components.
- Input to round function consists of key and the output of previous round
- These functions are obtained by the repeated application of Substitution (invertible SBoxes) and Permutation.
- Thus they are called Substitution Permutation Networks (SPN).

## Further Reading

- C. E. Shannon, *Communication Theory of Secrecy Systems*. Bell Systems Technical Journal, 28(1949), 656-715
- B. A Forouzan and D. Mukhopadhyay, *Cryptography & Network Security, Tata Mc Graw Hills, Chapter 5*
- Douglas Stinson, *Cryptography Theory and Practice, 2<sup>nd</sup> Edition*, Chapman & Hall/CRC