



# **CS31001 COMPUTER ORGANIZATION AND ARCHITECTURE**

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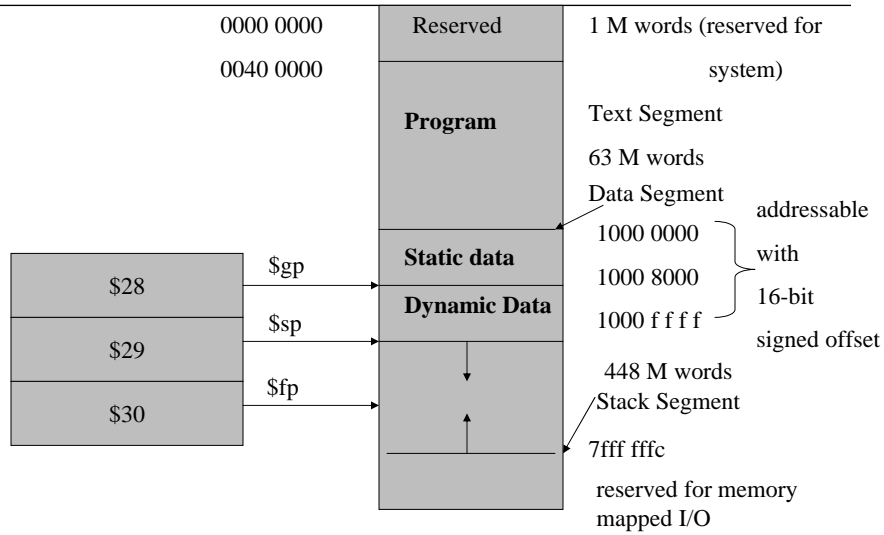
## Procedures and Data

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## The Stack for Data Storage

- We have seen that the stack can be used for procedure calls, and for temporary storage of data.
- Let us see the memory address space in the MIPS memory.

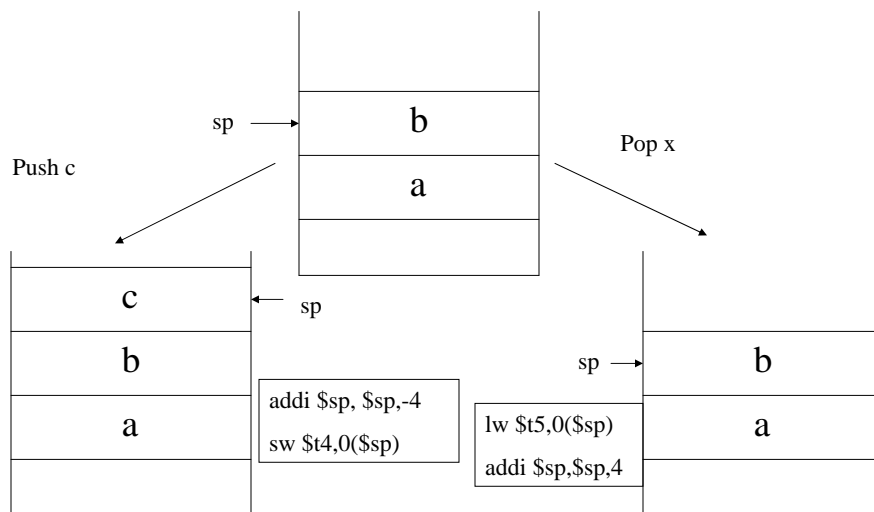
## Overview of the memory address space



## Push and Pop in Stack

- The stack pointer `$sp`, points to the top element in the stack.
- Push decrements the stack pointer and puts an element into the stack.
- Pop removes an element from the stack, and then adds the stack pointer.

## Effects of push and pop on the stack



## The pop operation

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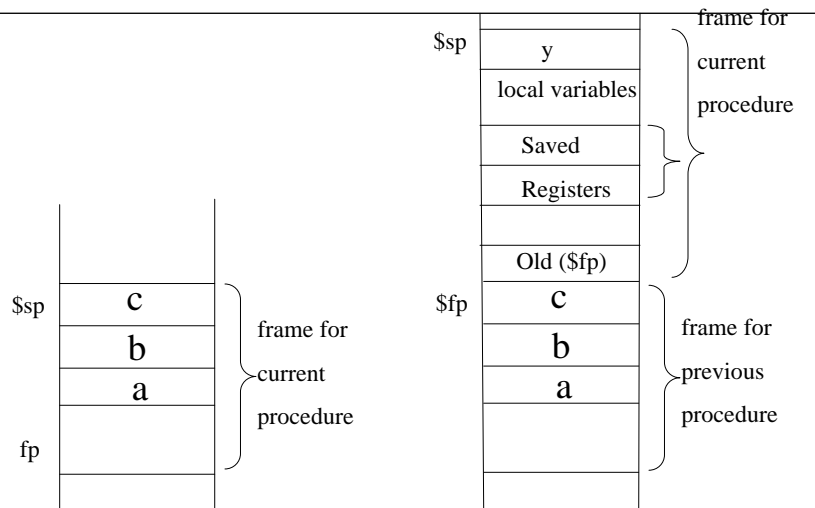
- Note that the pop operation does not erase the old top element, `c`
- `c` is still there, and would be still accessible by `-4($sp)`.
- However, it is not a part of the stack frame.
  - Thus this is a logical deletion.
- Hence to delete the top 10 elements, we can (logically) remove them by increasing the stack pointer by 40.

## The Stack Frame

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- The stack is used for various purposes.
- Two of the important ones are:
  - To pass more than 4 input parameters, or receive more than two results.
  - Place to store when calling other procedures (during nested call).
- Each procedure maintains an area: called as its stack frame.
  - Its delimiters are `$sp` (top of the stack frame), and `$fp`(frame-pointer, the other end).
- After the procedure terminates, the calling procedure expects to find the stack undisturbed
  - Thus it can restore the saved registers to their original values and proceed with its own computations.

## Use of the stack for procedures



## The Frame Pointer

- It provides a stable reference point for addressing memory words in the portion of the stack corresponding to the present procedure.
- The words in the current frame are  $4(\$fp)$ ,  $8(\$fp)$ , ...
- Though the stack pointer changes in the course of the procedure, the frame pointer holds a fixed address.
- **Note that the use of the frame pointer is entirely optional.**

## Example

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```
proc: sw $fp, -4($sp)
      addi $fp,$sp,0
      addi $sp,$sp,-12
      sw $ra,4($sp)
      sw $s0,0($sp)
      ...
      lw $s0,0($sp)
      lw $ra,4($sp)
      addi $sp,$fp,0
      lw $fp,8($sp)
      jr $ra
```

## Reduce unnecessary stack operations

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```
proc: sw $s0,-4($sp)
      ...
      lw $s0, -4($sp)
      jr $ra
```

This reduces the procedure call substantially.

## Some Other Instructions

- Special arithmetic/logical instructions
- mult/div
- mult \$s0, \$s1 #Hi, Lo are set to \$s0 x \$s1
- div \$s0, \$s1 #Hi is \$s0 mod \$s1
- mfhi \$t0 #set \$t0 to (Hi)
- mflo \$t0 #set \$t0 to (Lo)

## The mult and div instructions

000000	10000	10001	00000	00000	0110x0
ALU Instruction	Source register 1	Source register 2	Unused	Unused	function mult=24, div=26

000000	00000	00000	01000	00000	0100x0
ALU Instruction	Unused	Unused	Destination register	Unused	function mfhi=16 mflo=18

## The Shift Operations

000000	00000	10001	01000	00010	0000x0
ALU Instruction	Unused	Source register	Destination Register	Shift amount	function sll=0, srl=2

000000	10000	10001	01000	00000	0001x0
ALU Instruction	Amount Register	Unused	Destination register	Unused	function sllv=4 srlv=6

## Arrays and Pointers

- Often it is important to step through an array or list.
- Two basic ways:
  - **Index:** Uses a register that holds the index  $i$  and increment the register in each step to effect moving from element  $i$  of the list to element  $i+1$ .
  - **Pointer:** Uses a register that points to (holds the address of) the list element being examined, and updates it each step to point to the next element.



## Maximum sum prefix in a list of integers

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array base address A in \$a0, its length n in \$a1.

length of max-sum prefix: \$v0

associated sum: \$v1

## Program

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```
mshfx: addi $v0, $zero, 0 #initialize length
          addi $v1, $zero, 0 #initialize max sum
          addi $t0, $zero, 0 #initialize index to 0
          addi $t1, $zero, 0 #initialize running sum
loop: add $t2, $t0, $t0
          add $t2, $t2, $t2
          add $t3, $t2, $a0
          lw $t4, 0($t3)
```

## Program

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```
add $t1, $t1, $t4
slt $t5, $v1, $t1
bne $t5, $zero, mdfy
j test
mdfy: addi $v0, $v0, 1
        addi $v1, $t1, 0
test: addi $t0, $t0, 1
        slt $t5, $t0, $a1
        bne $t5, $zero, loop
done: jr $ra
```

## Selection Sort using Pointers

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```
#$a0 pointer to first element in unsorted array
#$a1 pointer to last element in unsorted array
#$t0 temporary place for value of last element
#$v0 pointer to max element in unsorted part
#$v1 value of max element in unsorted part
sort: beq $a0, $a1, done
      jal max
      lw $t0, 0($a1)
      sw $t0, 0($v0)
      sw $v1, 0($a1)
      addi $a1, $a1, -4
```

## Selection Sort using Pointers

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#\$a0 pointer to first element	loop: beq \$t0,\$a1,ret
#\$a1 pointer to last element	<b>addi \$t0,\$t0,4</b>
#\$t0 pointer to next element	lw \$t1,0(\$t0)
#\$t1 value of next element	slt \$t2,\$t1,\$v1
#\$t2 result of (next) < (max)	bne \$t2,\$zero,loop
#\$v0 pointer to max element	addi \$v0,\$t0,0 #update
#\$v1 value of max element	addi \$v1,\$t1,0 #new max
max: addi \$v0,\$a0,0	j loop
lw \$v1,0(\$v0)	ret: jr \$ra
addi \$t0,\$a0,0	

Rest of Procedures, refer to  
Tutorial provided in the Lab.

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