



CS10003

PROGRAMMING AND DATA STRUCTURES

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Computer Science & Engineering Department

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Recursion – Advanced Examples

Sudeshna Sarkar

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Paying with fewest coins

- A country has coins of denomination 3, 5 and 10, respectively.
- We are to write a function `canchange(k)` that returns `-1` if it is not possible to pay a value of `k` using these coins.
 - Otherwise it returns the minimum number of coins needed to make the payment.
- For example, `canchange(7)` will return `-1`.
- On the other hand, `canchange(14)` will return 4 because 14 can be paid as `3+3+3+5` and there is no other way to pay with fewer coins.
- Finally, 15 can be changed as `3+3+3+3+3`, `5+5+5`, `5+10`, so `canchange (15)` will return 2.

Paying with fewest coins

```
int canchange( int V)
{
    int a;
    if (V==0) return 0;
    if ( (V ==3) || (V == 5) || (V == 10) ) return 1;
    if (V< 3) return -1 ;

    a = canchange( V- 10 ); if (a > 0) return a+1 ;
    a = canchange( V - 5 ); if (a > 0) return a+1 ;
    a = canchange( V - 3 ); if (a > 0) return a+1 ;
    return -1;
}
```

Exercise: Rewrite this code if the denominations are 3, 8, and 10. Do you see a problem? Repair it.

Find minimum number of coins to make a given value using Recursion

The minimum number of coins for a value V can be computed using the recursive formula below.

- If $V == 0$:

0 coins required

- If $V > 0$:

$\text{minCoins}(\text{coins}[0..m-1], V) = \min \{ 1 + \text{minCoins}(V - \text{coin}[i]) \}$

where, $0 \leq i \leq m-1$ and $\text{coins}[i] \leq V$.

Paying with fewest coins

```
int canchange( int V)
{
    int a, min;
    if (V==0) return 0;
    if ( (V ==3) || (V == 8) || (V == 10) ) return 1;
    if (V< 3) return -1 ;
    min = 0;
    a = canchange( V- 10 ); if (a> 0) min = a; ;
    a = canchange( V - 8 ); if ((a > 0) && (a < min)) min = a;
    a = canchange( V - 3 ); if ((a > 0) && (a < min)) min = a;
    if (min > 0)
        return a;
    return -1
}
```

Paying with fewest coins

```
int canchange( int V, int coins[], int numcoins) {
    int a, i, min=-1;
    if (V==0) return 0;
    if (V<0) return -1;
    for (i=0; i<numcoins; i+=) {
        a= canchange( V- coins[i], coins, numcoins);
        if (a==-1) continue;
        if ((a+1 < min) || (min==-1)) min = a+1;
    }
    return min;
}
```

```
int main() {
    int coins[3]={10,8,3};
    int V;
    scanf ("%d", &V);
    printf ("min coins = %d\n",
           canchange (V,coins, 3));
    return 0;
}
```

Count number of ways to make a change of V

- **Note:** Assume that you have an infinite supply of each type of coin.

Examples:

- **Input:** sum = 4, coins[] = {1,2,3},
- **Output:** 4
- **Explanation:** there are four solutions:
{1, 1, 1, 1}, {1, 1, 2}, {2, 2}, {1, 3}.

Count number of ways to make a change of V

- For each coin, there are 2 options.
- **Include the current coin:** Subtract the current coin's denomination from the target sum and call the count function recursively with the updated sum and the same set of coins i.e., $\text{count}(\text{coins}, n, \text{sum} - \text{coins}[n-1])$
- **Exclude the current coin:** Call the count function recursively with the same sum and the remaining coins. i.e., $\text{count}(\text{coins}, n-1, \text{sum})$.

- **Recurrence Relation**

$$\text{Count}(\text{coins}, n, \text{sum}) = \text{count}(\text{coins}, n, \text{sum} - \text{coins}[n-1]) + \text{count}(\text{coins}, n-1, \text{sum})$$


```
// Returns the count of ways we can sum coins[0...n-1]
// coins to get "sum"
int count(int coins[], int n, int sum) {
    // If sum is 0 then there is 1 solution (do not include any coin)
    if (sum == 0)
        return 1;
    // If sum is less than 0 then no solution exists
    if (sum < 0)
        return 0;
    // If there are no coins and sum is greater than 0,
    //then no solution exist
    if (n <= 0)
        return 0;
    return count(coins, n - 1, sum)+ count(coins, n, sum - coins[n - 1]);
}
```

```
// Driver program
int main()
{
    int i, j;
    int coins[] = { 1, 2, 3 };
    int n = 3;
    printf("%d ", count(coins, n, 5));
    return 0;
}
```

Generate all unique partitions of an integer

- Given a positive integer n , the task is to generate all possible unique ways to represent n as sum of positive integers.
- Examples:

Input: 4

Output:

4
3 1
2 2
2 1 1
1 1 1 1

Input: 3

Output:

3
2 1
1 1 1

Generate all **unique** partitions of an integer

```
void partition (int arr[], int size, int n) {
    int i;
    if (n==0) {
        for (i=0; i<size; i++)
            printf ("%3d", arr[i]) ;
        printf ("\n") ;
        return;
    }
    for (i=n; i>0; i--) {
        arr[size] = i;
        partition (arr, size+1, n-i) ;
    }
}
```

```
int main(){
    int n;
    int A[100] ;
    scanf ("%d", &n);
    printf ("n = %d\n", n) ;
    partition (A, 0, n) ;
    return 0;
}
```

```
n=4
4
3 1
2 2
2 1 1
1 3
1 2 1
1 1 2
1 1 1 1
```

```

#include <stdio.h>
void partition (int arr[], int size, int n, int max) {
    int i;    if (n<0) return;
    if (n==0) {
        for (i=0; i<size; i++)
            printf ("%3d", arr[i]) ;
        printf ("\n") ;
        return;
    }
    for (i=max; i>0; i--) {
        arr[size] = i;
        partition (arr, size+1, n-i, i) ;
    }
}

```

```

int main(){
    int n, size = 0, max;
    int A[100] ;
    scanf ("%d", &n);
    printf ("n = %d\n", n) ;
    max = n;
    partition (A, 0, n, max) ;
    return 0;
}

```

```

n = 7
7
6 1
5 2
5 1 1
4 3
4 2 1
4 1 1 1
3 3 1
3 2 2
3 2 1 1
3 1 1 1 1
2 2 2 1
2 2 1 1 1
2 1 1 1 1 1
1 1 1 1 1 1 1

```