CS10003: Programming & Data Structures

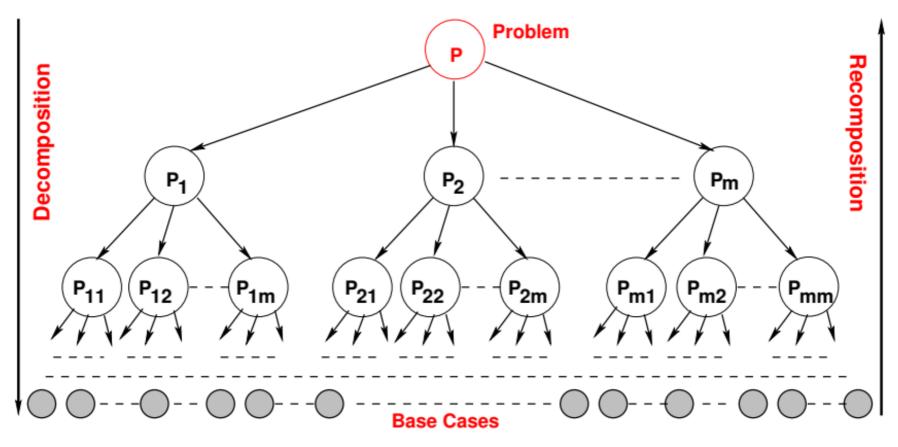
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Problem Solving using Recursion

Solving Problems Recursively

- Top-Down Approach:
 - Base Cases: elementary problem instances with known solutions
 - Decompose: split problems into smaller instance sub-problems
 - Recursive Calls: call same function recursively for each sub-problems
 - Recompose: combine solutions from sub-problems to get result

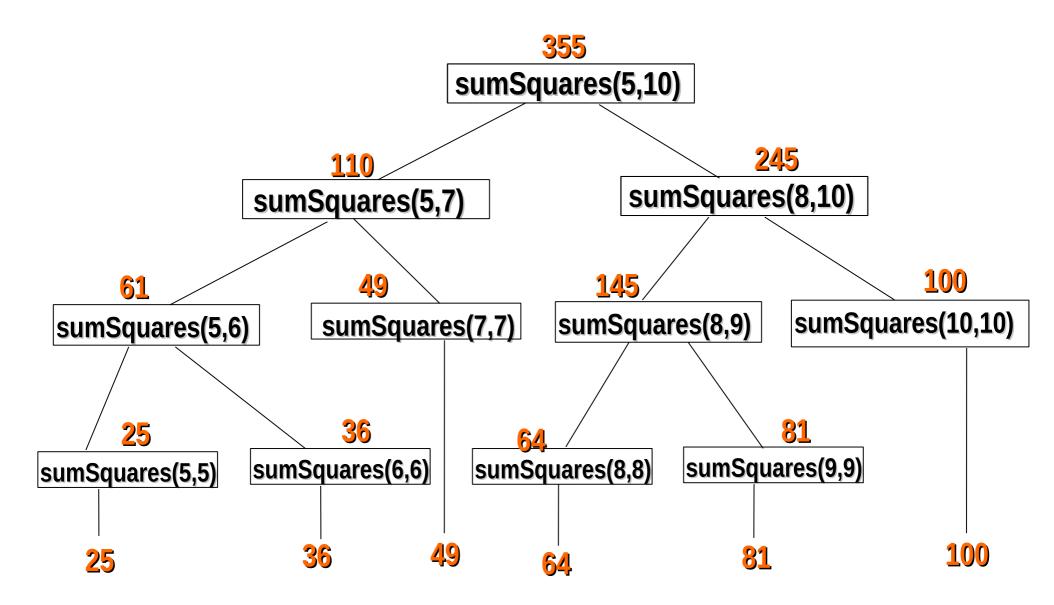


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Problem-0: Sum of Squares (Recall ...)

- Sum of squares within range [m,n], i.e.
 m² + (m+1)² + (m+2)² + ... + (n-1)² + n²
- Solution: int sumSquares (int m, int n) { Base int middle ; if (m == n) return m*m; Decompose else { middle = (m+n)/2; Recurse return sumSquares(m,middle)* Recompose + sumSquares(middle+1,n); }

Problem-0: Annotated Recursion Tree



Problem – 1 : Searching an Element

- Function, find (L, n, x), to find element x from unordered list L of n elements
 - **Base Case:** if n == 1, then compare the only element with x and return
 - Decompose: split list, L, into two parts, L_1 and L_2 , having n_1 and n_2 elements each ($n_1 + n_2 = n$), respectively

 $\exists \text{ Recursive Calls: } f_1 \leftarrow \text{find } (L_1, n_1, x) \text{ and } f_2 \leftarrow \text{find } (L_2, n_2, x)$

Recompose: if $(f_1 || f_2)$, then return found, else return not-found

- Points to Ponder:
 - Does the split position really matters?
 - Middle, Arbitrary, Always 1 and (n-1) sized-division ... etc.
 - The naive approach is a special case of this recursive solution
 - Always split into 1 and (n-1) sized lists (think!)
 - □ Need to compare n times to find x among n elements

Problem – 2 : Finding MAX

- Function, max (L, n), to find maximum from a list L of n elements:
 - Base Case: if n == 1, then return the only element
 - Decompose: split list, L, into two parts, L_1 and L_2 , having n_1 and n_2 elements each $(n_1 + n_2 = n)$, respectively

 $\Box \text{ Recursive Calls: } m_1 \leftarrow \max(L_1, n_1) \text{ and } m_2 \leftarrow \max(L_2, n_2)$

 \Box **Recompose:** if (m₁ > m₂), then return m₁, else return m₂

- Points to Ponder:
 - Does the split position really matters?
 - Middle, Arbitrary, Always 1 and (n-1) sized-division ... etc.
 - \Box The naive approach is a special case of this recursive solution
 - Always split into (n-1) and 1 sized lists (think!)
 - Need to compare (n-1) times to get maximum among n elements

Problem – 3: Finding MAX+MIN

Function, maxmin (L, n), to find max+min from list L of n elements:

Base Cases:

- if n == 1, then return the only element as maximum and minimum
- If n == 2, then compare between two elements and return maximum and minimum

Decompose: split list, L, into two parts, L_1 and L_2 , having n_1 and n_2 elements each $(n_1 + n_2 = n)$, respectively

Recursive Calls:

- $h_1 \leftarrow \max(L_1, n_1)$ and $h_2 \leftarrow \max(L_2, n_2)$
- $I_1 \leftarrow \min(L_1, n_1) \text{ and } I_2 \leftarrow \min(L_2, n_2)$

Recompose:

- if $(h_1 > h_2)$, then return h_1 as maximum, else return h_2 as maximum
- if $(I_1 < I_2)$, then return I_1 as minimum, else return I_2 as minimum

Problem – 3: Finding MAX+MIN

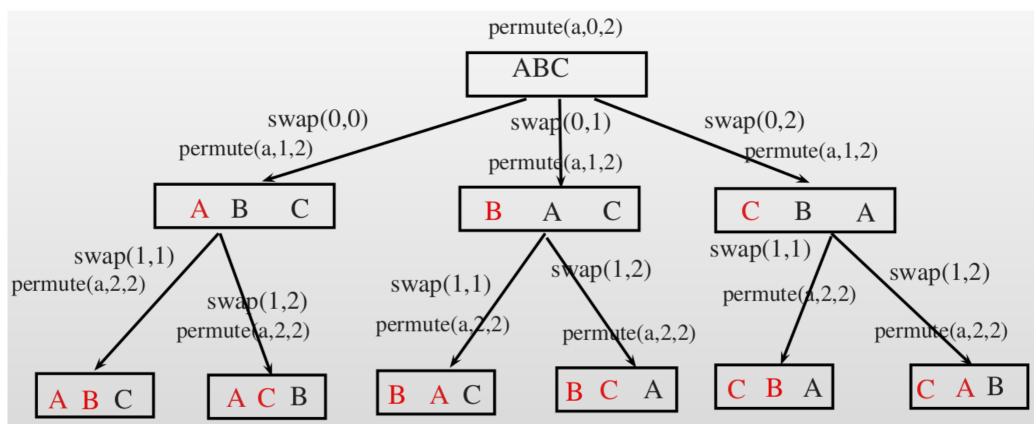
- Points to Ponder:
 - The naive approach is to traverse the elements twice once for finding max and again for finding min

Total comparisons = (n - 1) + (n - 1) = 2n - 2

- Smart Observations leads to better approach!
 - Take two elements at a time when finding max
 - Total comparisons for max finding = n/2 + n/4 + ... + 1 = (n 1)
 - Then, the first round losers (i.e. n/2 elements) are candidates for min
 - Total comparisons for min finding now = (n/2 1)
- Does the split position really matters to reduce number of comparison?
 - Check for comparisons with 1 and n-1 split \rightarrow 2(n 1)
 - **Check for comparisons with 2 and n-2 split** \rightarrow **3n/2 2**
 - **Check for comparisons with middle split** \rightarrow **3n/2 2**

Problem – 4: Generate all Permutations of String

Strategy:



Recursive Formulation and Solution:

Do Yourself ! :-)

Thank You!