



CS10003: Programming & Data Structures

Dept. of Computer Science & Engineering
Indian Institute of Technology Kharagpur

Autumn 2020



Recursion

Recursion

- A process by which a function calls itself repeatedly
 - Either directly.
 - X calls X
 - Or cyclically in a chain.
 - X calls Y, and Y calls X
- Used for repetitive computations in which each action is stated in terms of a previous result
$$\text{fact}(n) = n * \text{fact}(n-1)$$

Recursion Conditions

- For a problem to be written in recursive form, two conditions are to be satisfied:
 - It should be possible to express the problem in recursive form
 - Solution of the problem in terms of solution of the same problem on smaller sized data
 - The problem statement must include a stopping condition

$$\begin{aligned} \text{fact}(n) &= 1, & \text{if } n = 0 & \leftarrow \text{Stopping condition} \\ &= n * \text{fact}(n-1), & \text{if } n > 0 & \leftarrow \text{Recursive definition} \end{aligned}$$

Examples

- Factorial:

$$\text{fact}(0) = 1$$

$$\text{fact}(n) = n * \text{fact}(n-1), \text{ if } n > 0$$

- GCD:

$$\text{gcd}(m, m) = m$$

$$\text{gcd}(m, n) = \text{gcd}(m\%n, n), \text{ if } m > n$$

$$\text{gcd}(m, n) = \text{gcd}(n, n\%m), \text{ if } m < n$$

- Fibonacci series (1,1,2,3,5,8,13,21,...)

$$\text{fib}(0) = 1$$

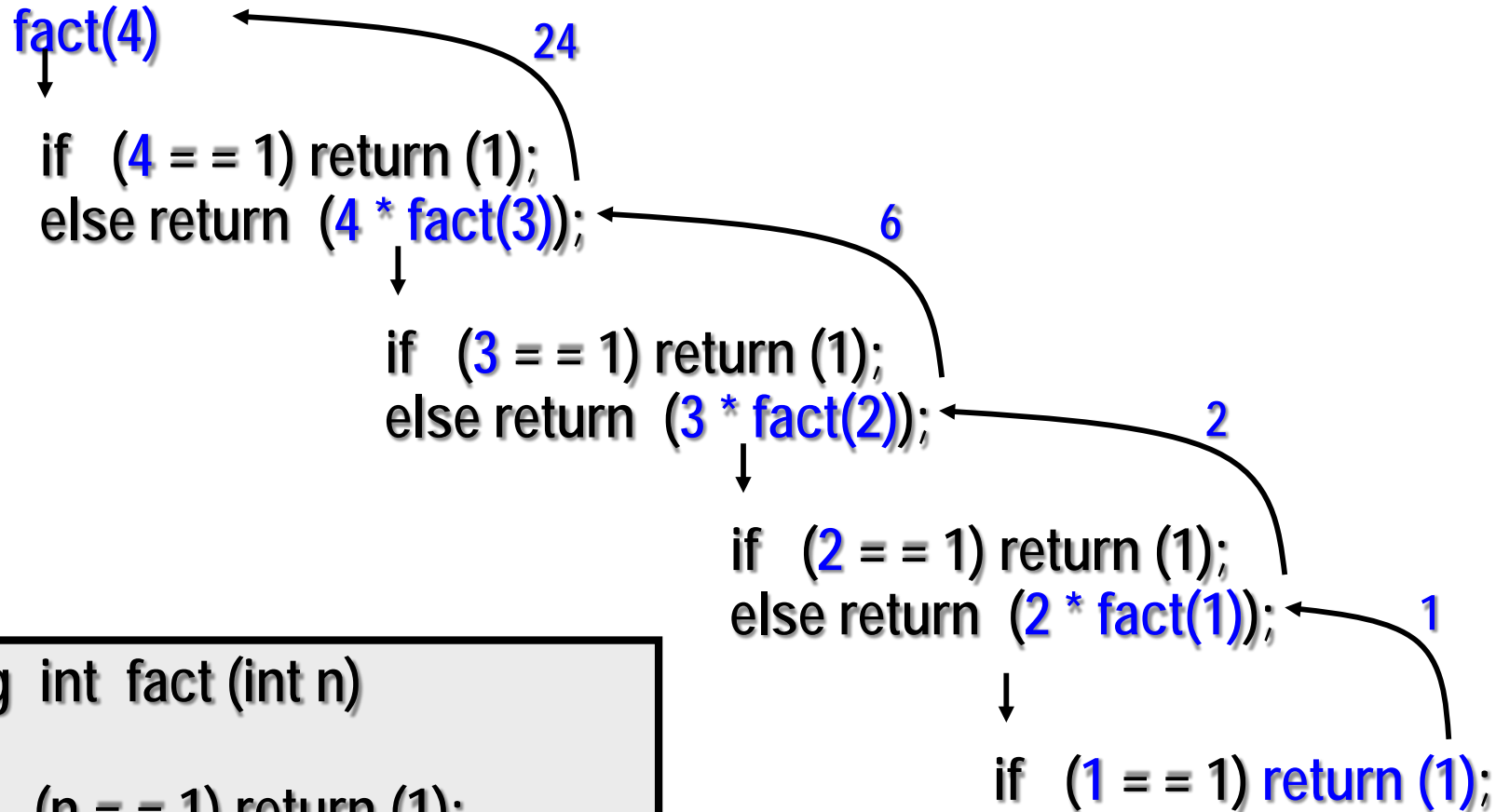
$$\text{fib}(1) = 1$$

$$\text{fib}(n) = \text{fib}(n-1) + \text{fib}(n-2), \text{ if } n > 1$$

Factorial

```
long int fact (int n)
{
    if (n == 1)
        return (1);
    else
        return (n * fact(n-1));
}
```

Factorial Execution



```
long int fact (int n)
{
  if (n == 1) return (1);
  else return (n * fact(n-1));
}
```

Look at Variable Addresses (a slightly different program) !

```
void main()
{
    int x,y;
    scanf("%d",&x);
    y = fact(x);
    printf ("M: x= %d, y = %d\n", x,y);
}

int fact(int data)
{ int val = 1;
  printf("F: data = %d, &data = %u \n
    &val = %u\n", data, &data, &val);
  if (data>1) val = data*fact(data-1);
  return val;
}
```

Output

```
4
F: data = 4, &data = 3221224528
  &val = 3221224516
F: data = 3, &data = 3221224480
  &val = 3221224468
F: data = 2, &data = 3221224432
  &val = 3221224420
F: data = 1, &data = 3221224384
  &val = 3221224372
M: x= 4, y = 24
```


Fibonacci Numbers

Fibonacci recurrence:

$$\begin{aligned} \text{fib}(n) &= 1 \text{ if } n = 0 \text{ or } 1; \\ &= \text{fib}(n - 2) + \text{fib}(n - 1) \\ &\text{otherwise;} \end{aligned}$$

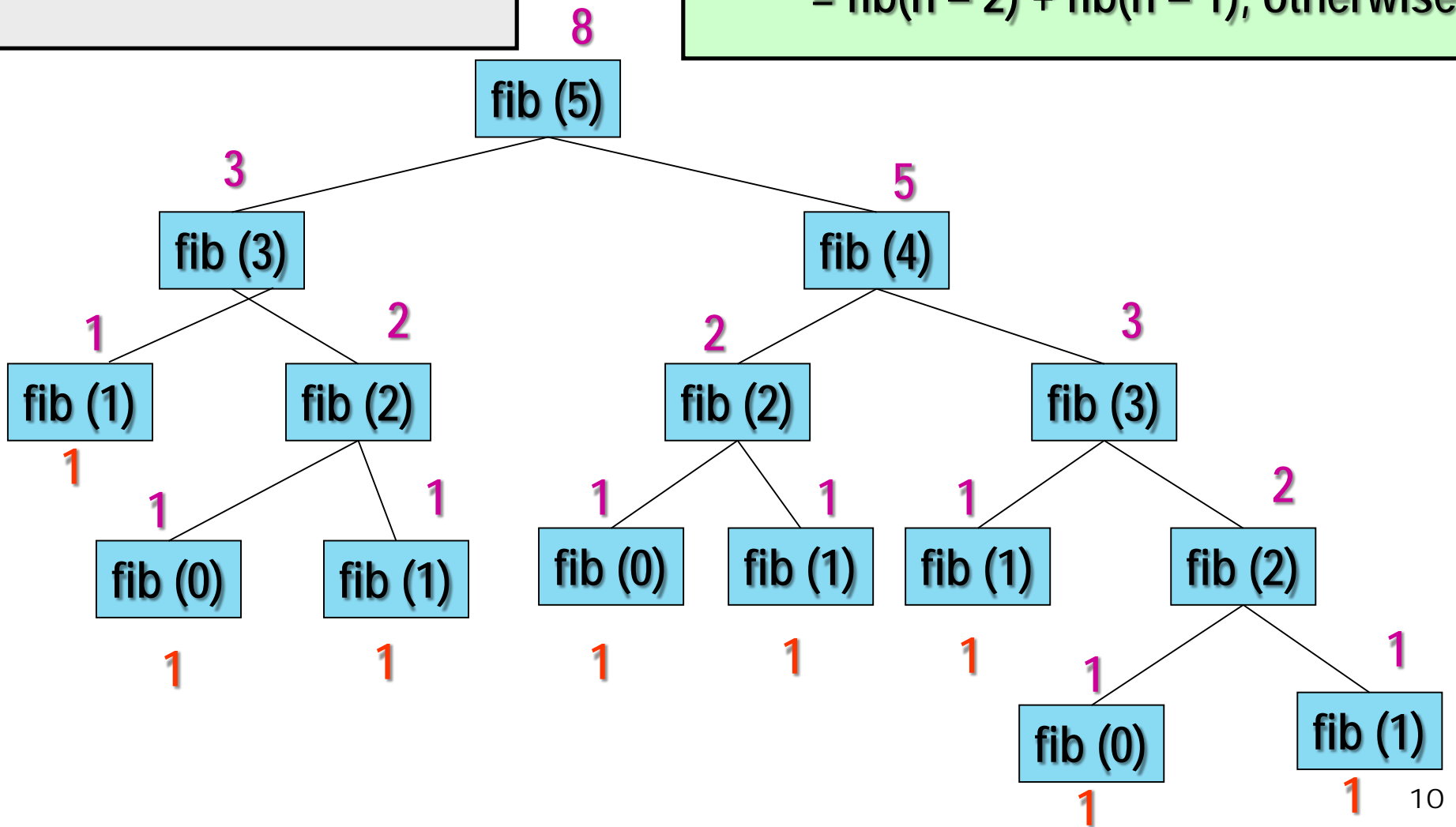
```
int fib (int n) {  
    if (n == 0 or n == 1)  
        return 1; [BASE]  
    return fib(n-2) + fib(n-1);  
                [Recursive]  
}
```

Fibonacci Numbers

```
int fib (int n) {  
    if (n == 0 || n == 1) return 1;  
    return fib(n-2) + fib(n-1);  
}
```

Fibonacci recurrence:

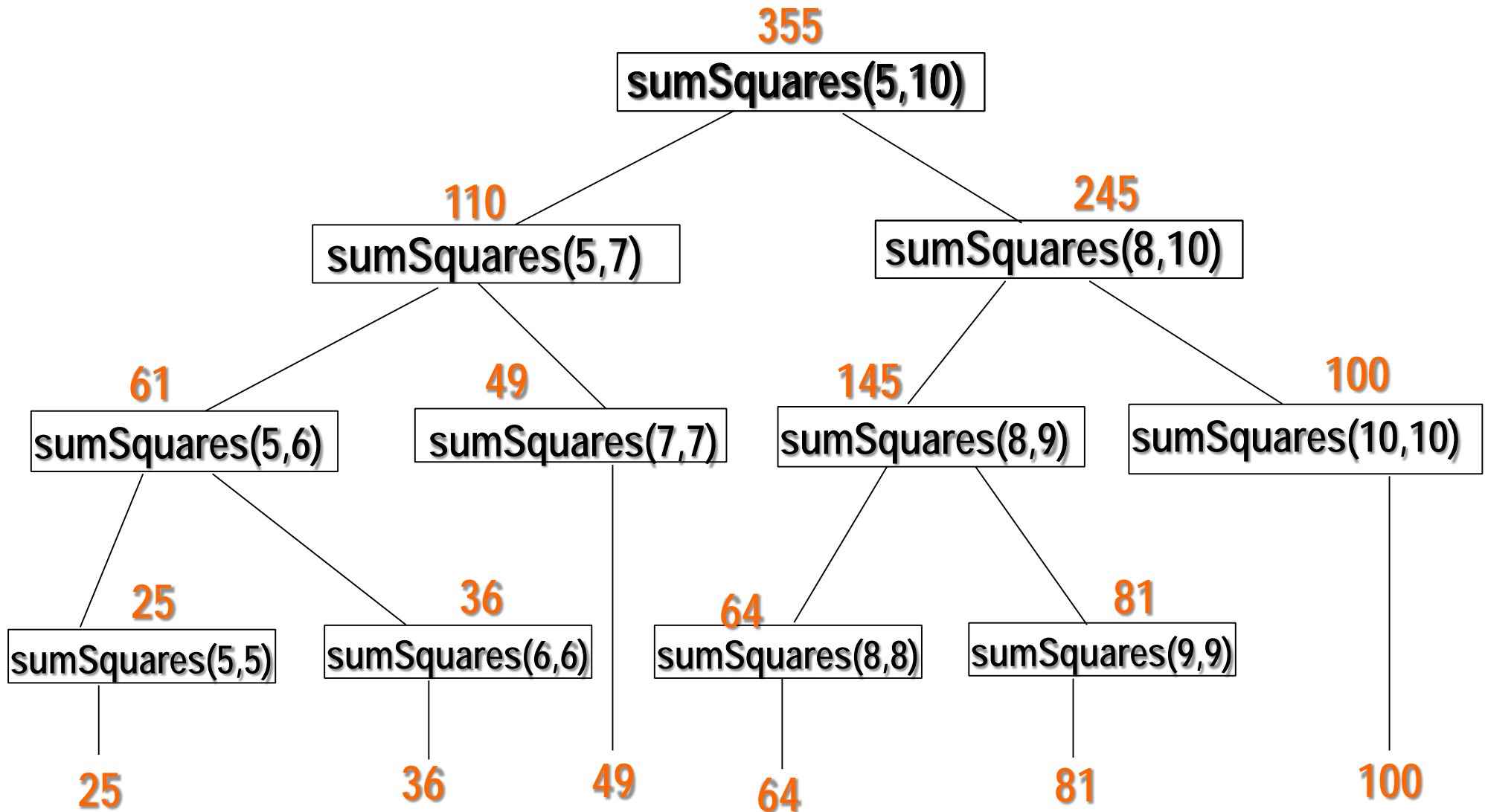
$$\text{fib}(n) = \begin{cases} 1, & \text{if } n = 0 \text{ or } 1 \\ \text{fib}(n - 2) + \text{fib}(n - 1), & \text{otherwise} \end{cases}$$



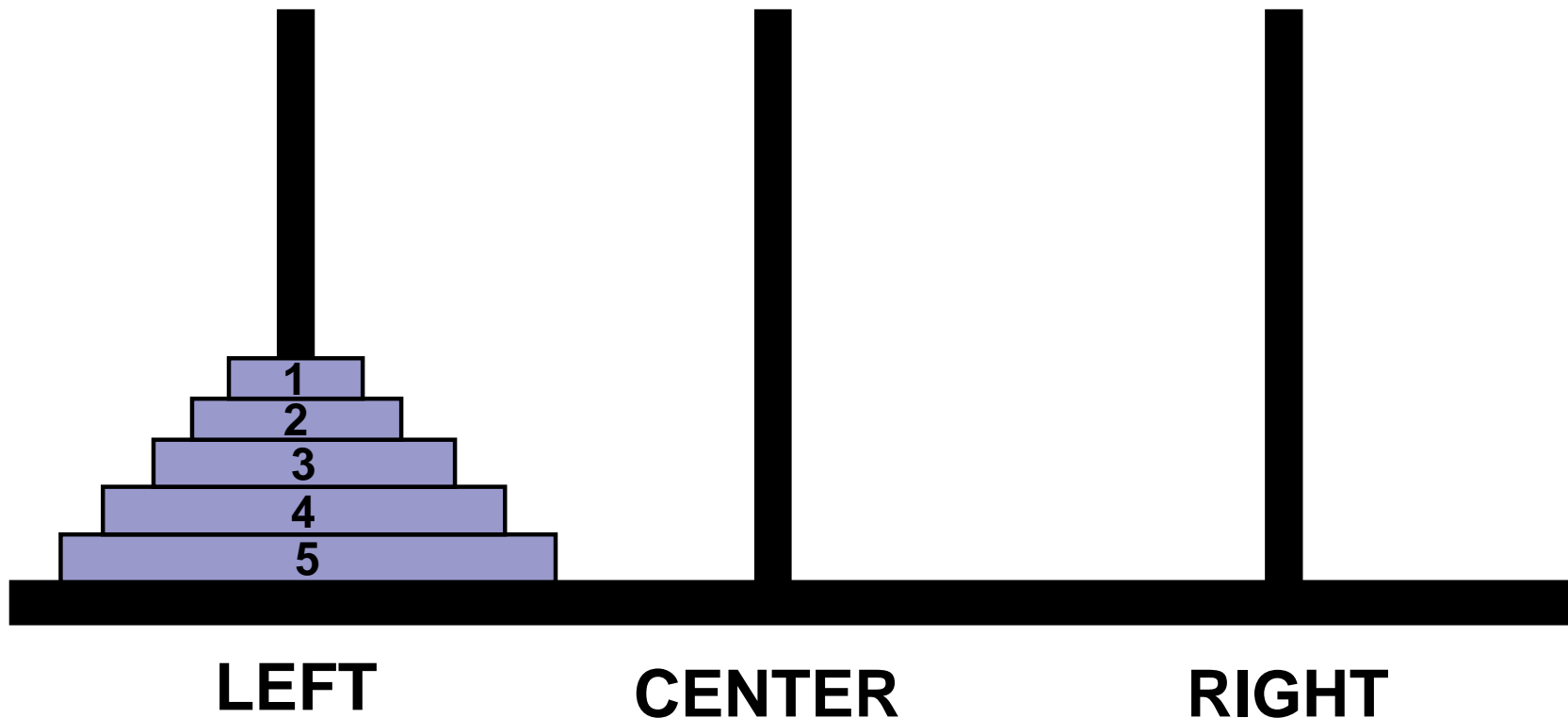
Sum of Squares

```
int sumSquares (int m, int n)
{
    int middle ;
    if (m == n) return m*m;
    else
    {
        middle = (m+n)/2;
        return sumSquares(m,middle)
            + sumSquares(middle+1,n);
    }
}
```

Annotated Call Tree



Towers of Hanoi Problem

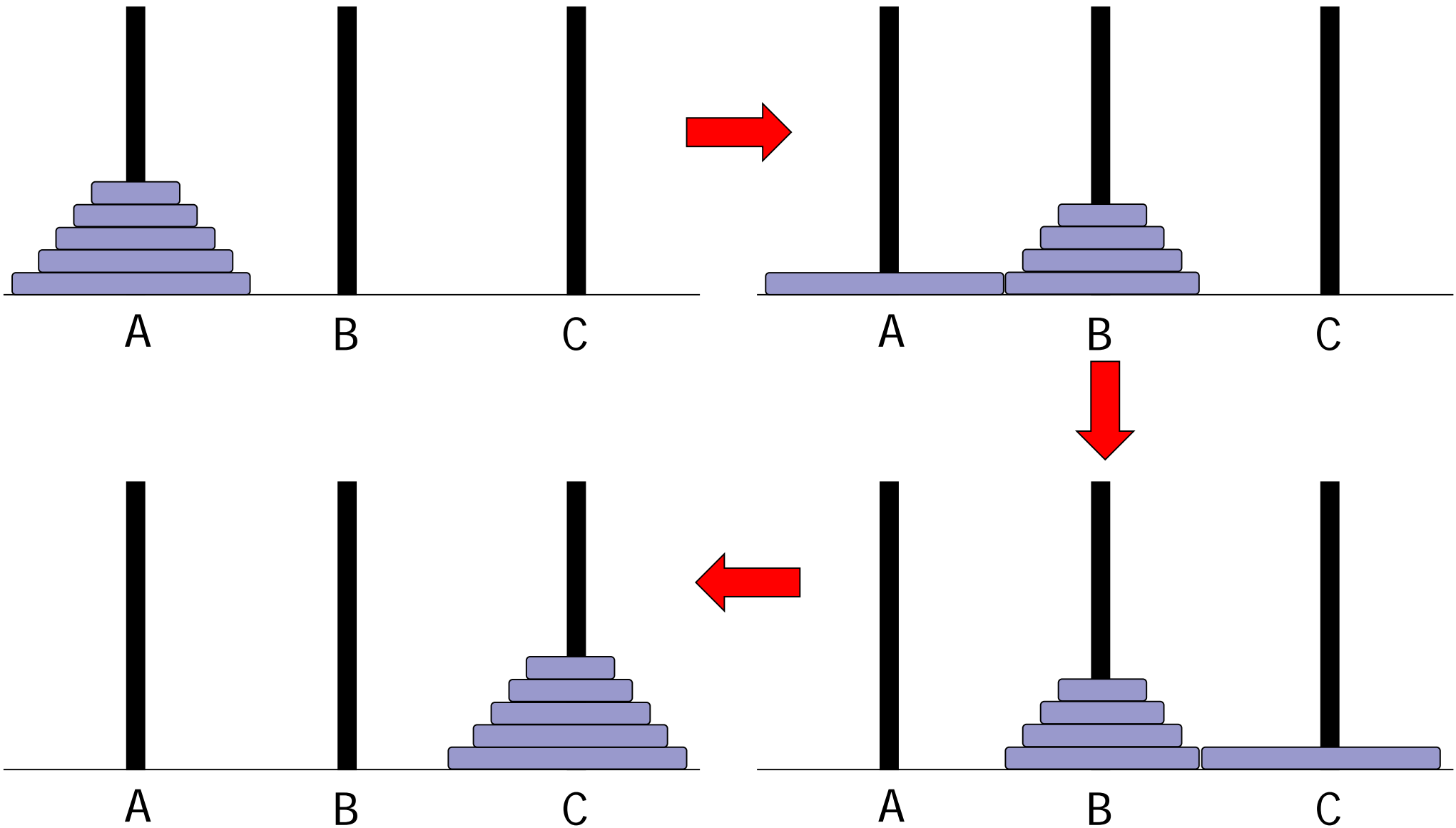


- Initially all the disks are stacked on the LEFT pole
- Required to transfer all the disks to the RIGHT pole
 - Only one disk can be moved at a time.
 - A larger disk cannot be placed on a smaller disk
- CENTER pole is used for temporary storage of disks

Towers of Hanoi Problem: Recursion

- Recursive statement of the general problem of n disks
 - Step 1:
 - Move the top $(n-1)$ disks from LEFT to CENTER
 - Step 2:
 - Move the largest disk from LEFT to RIGHT
 - Step 3:
 - Move the $(n-1)$ disks from CENTER to RIGHT

Tower of Hanoi: Recursive Steps



Towers of Hanoi function

```
void towers (int n, char from, char to, char aux)
{
    /* Base Condition */
    if (n==1) {
        printf ("Disk 1 : %c → %c \n", from, to) ;
        return ;
    }
    /* Recursive Condition */
    towers (n-1, from, aux, to) ;
    .....
    .....
}
```


Towers of Hanoi function

```
void towers (int n, char from, char to, char aux)
{
    /* Base Condition */
    if (n==1) {
        printf ("Disk 1 : %c → %c \n", from, to) ;
        return ;
    }
    /* Recursive Condition */
    towers (n-1, from, aux, to) ;
    printf ("Disk %d : %c → %c\n", n, from, to) ;
    .....
}
```

Towers of Hanoi function

```
void towers (int n, char from, char to, char aux)
{
    /* Base Condition */
    if (n==1) {
        printf ("Disk 1 : %c → %c \n", from, to) ;
        return ;
    }
    /* Recursive Condition */
    towers (n-1, from, aux, to) ;
    printf ("Disk %d : %c → %c\n", n, from, to) ;
    towers (n-1, aux, to, from) ;
}
```

TOH runs

```
void towers(int n, char from, char to, char aux)
{ if (n==1)
  { printf ("Disk 1 : %c -> %c \n", from, to) ;
    return ;
  }
  towers (n-1, from, aux, to) ;
  printf ("Disk %d : %c -> %c\n", n, from, to) ;
  towers (n-1, aux, to, from) ;
}

void main()
{ int n;
  scanf("%d", &n);
  towers(n,'A','C','B');
}
```

Output

```
3
Disk 1 : A -> C
Disk 2 : A -> B
Disk 1 : C -> B
Disk 3 : A -> C
Disk 1 : B -> A
Disk 2 : B -> C
Disk 1 : A -> C
```

More TOH runs

```
void towers(int n, char from, char to, char aux)
{ if (n==1)
  { printf ("Disk 1 : %c -> %c \n", from, to) ;
    return ;
  }
  towers (n-1, from, aux, to) ;
  printf ("Disk %d : %c -> %c\n", n, from, to) ;
  towers (n-1, aux, to, from) ;
}

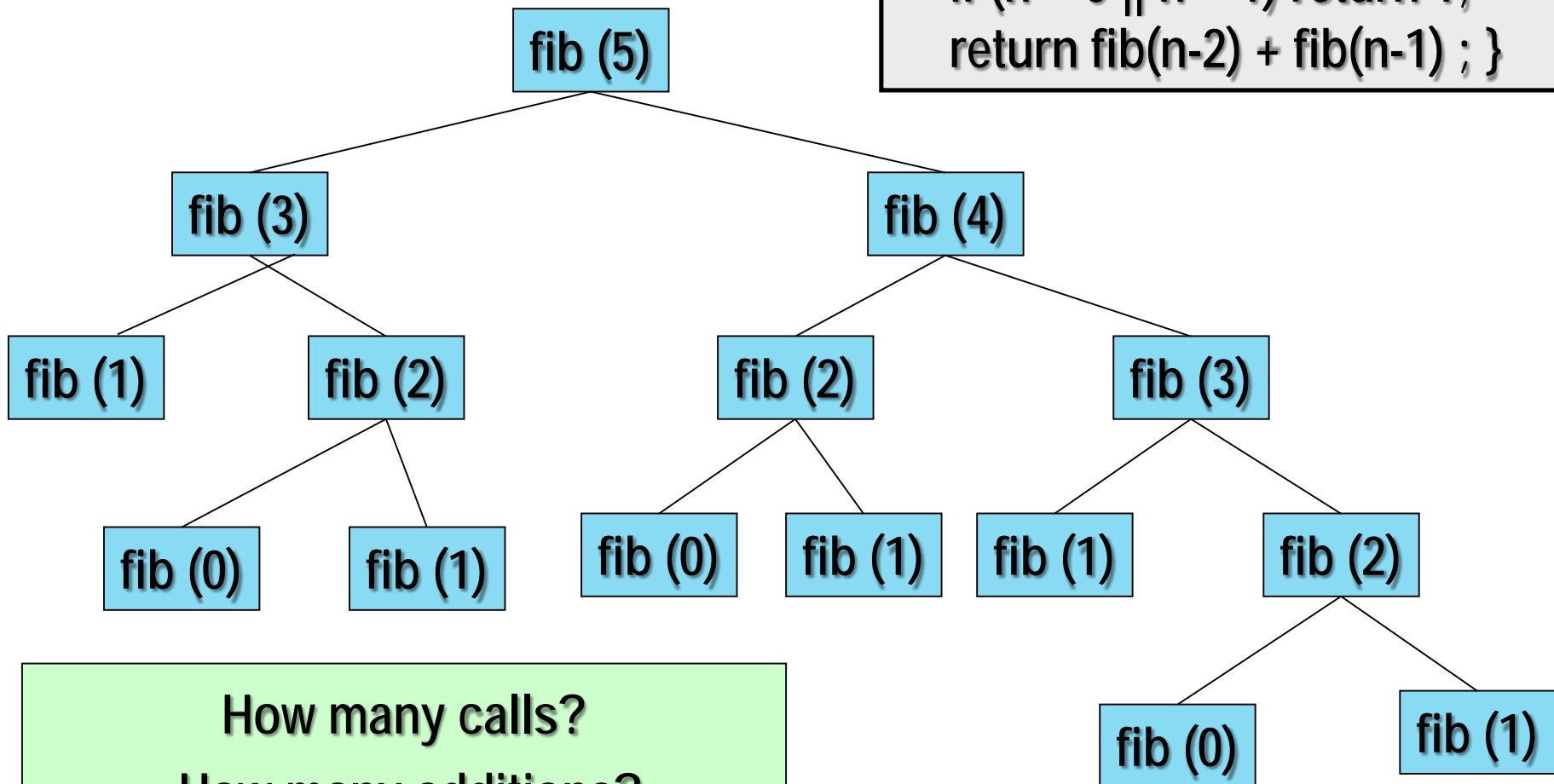
void main()
{ int n;
  scanf("%d", &n);
  towers(n,'A','C','B');
}
```

```
4
Disk 1 : A -> B
Disk 2 : A -> C
Disk 1 : B -> C
Disk 3 : A -> B
Disk 1 : C -> A
Disk 2 : C -> B
Disk 1 : A -> B
Disk 4 : A -> C
Disk 1 : B -> C
Disk 2 : B -> A
Disk 1 : C -> A
Disk 3 : B -> C
Disk 1 : A -> B
Disk 2 : A -> C
Disk 1 : B -> C
```

Relook at recursive Fibonacci:

Not efficient !! Same sub-problem solved many times.

```
int fib (int n) {  
    if (n==0 || n==1) return 1;  
    return fib(n-2) + fib(n-1); }
```



How many calls?
How many additions?

Iterative Fib

```
int fib( int n)
{ int i=2, res=1, m1=1, m2=1;
  if (n ==0 || n ==1) return res;
  for ( ; i<=n; i++) {
    res = m1 + m2;
    m2 = m1;
    m1 = res;
  }
  return res;
}

void main()
{ int n;
  scanf("%d", &n);
  printf(" Fib(%d) = %d \n", n, fib(n));
}
```

**Much Less
Computation here!
(How many additions?)**

An efficient recursive Fib

```
int Fib ( int, int, int, int);

void main()
{
    int n;
    scanf("%d", &n);
    if (n == 0 || n ==1)
        printf("F(%d) = %d \n", n, 1);
    else
        printf("F(%d) = %d \n", n, Fib(1,1,n,2));
}
```

```
int Fib(int m1, int m2, int n, int i)
{
    int res;
    if (n == i)
        res = m1+ m2;
    else
        res = Fib(m1+m2, m1, n, i+1);
    return res;
}
```

Much Less Computation here!
(How many calls/additions?)

Run

```
int Fib ( int, int, int, int);
void main()
{ int n;
  scanf("%d", &n);
  if (n == 0 || n ==1) printf("F(%d) = %d \n", n, 1);
  else printf("F(%d) = %d \n", n, Fib(1,1,n,2));
}
int Fib(int m1, int m2, int n, int i)
{ int res;
  printf("F: m1=%d, m2=%d, n=%d, i=%d\n",
        m1,m2,n,i);

  if (n == i)
    res = m1+ m2;
  else
    res = Fib(m1+m2, m1, n, i+1);
  return res;
}
```

Output

3

F: m1=1, m2=1, n=3, i=2

F: m1=2, m2=1, n=3, i=3

F(3) = 3

5

F: m1=1, m2=1, n=5, i=2

F: m1=2, m2=1, n=5, i=3

F: m1=3, m2=2, n=5, i=4

F: m1=5, m2=3, n=5, i=5

F(5) = 8

Static Variables

```
int Fib (int, int);

void main()
{
    int n;
    scanf("%d", &n);
    if (n == 0 || n ==1)
        printf("F(%d) = %d \n", n, 1);
    else
        printf("F(%d) = %d \n", n,
Fib(n,2));
}
```

```
int Fib(int n, int i)
{
    static int m1, m2;
    int res, temp;
    if (i==2) {m1 =1; m2=1;}
    if (n == i) res = m1+ m2;
    else
    { temp = m1;
      m1 = m1+m2;
      m2 = temp;
      res = Fib(n, i+1);
    }
    return res;
}
```

Static variables remain in existence rather than coming and going each time a function is activated

Static Variables: See the addresses!

```
int Fib(int n, int i)
{
    static int m1, m2;
    int res, temp;
    if (i==2) {m1 =1; m2=1;}
    printf("F: m1=%d, m2=%d, n=%d,
           i=%d\n", m1,m2,n,i);
    printf("F: &m1=%u, &m2=%u\n",
           &m1,&m2);
    printf("F: &res=%u, &temp=%u\n",
           &res,&temp);
    if (n == i) res = m1+ m2;
    else { temp = m1; m1 = m1+m2;
          m2 = temp;
          res = Fib(n, i+1);  }
    return res;
}
```

Output

```
5
F: m1=1, m2=1, n=5, i=2
F: &m1=134518656, &m2=134518660
F: &res=3221224516, &temp=3221224512
F: m1=2, m2=1, n=5, i=3
F: &m1=134518656, &m2=134518660
F: &res=3221224468, &temp=3221224464
F: m1=3, m2=2, n=5, i=4
F: &m1=134518656, &m2=134518660
F: &res=3221224420, &temp=3221224416
F: m1=5, m2=3, n=5, i=5
F: &m1=134518656, &m2=134518660
F: &res=3221224372, &temp=3221224368
F(5) = 8
```

Recursion vs. Iteration

- Repetition
 - Iteration: explicit loop
 - Recursion: repeated function calls
- Termination
 - Iteration: loop condition fails
 - Recursion: base case recognized
- Both can have infinite loops
- Balance
 - Choice between performance (iteration) and good software engineering (recursion).

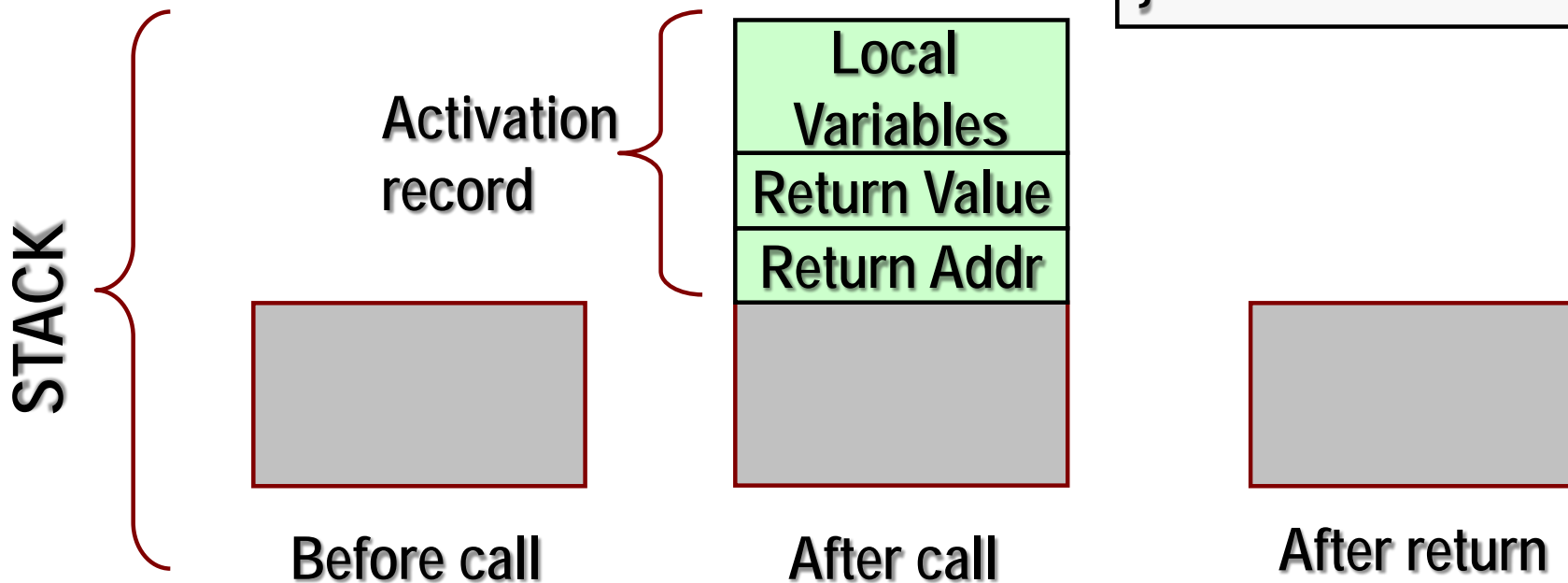
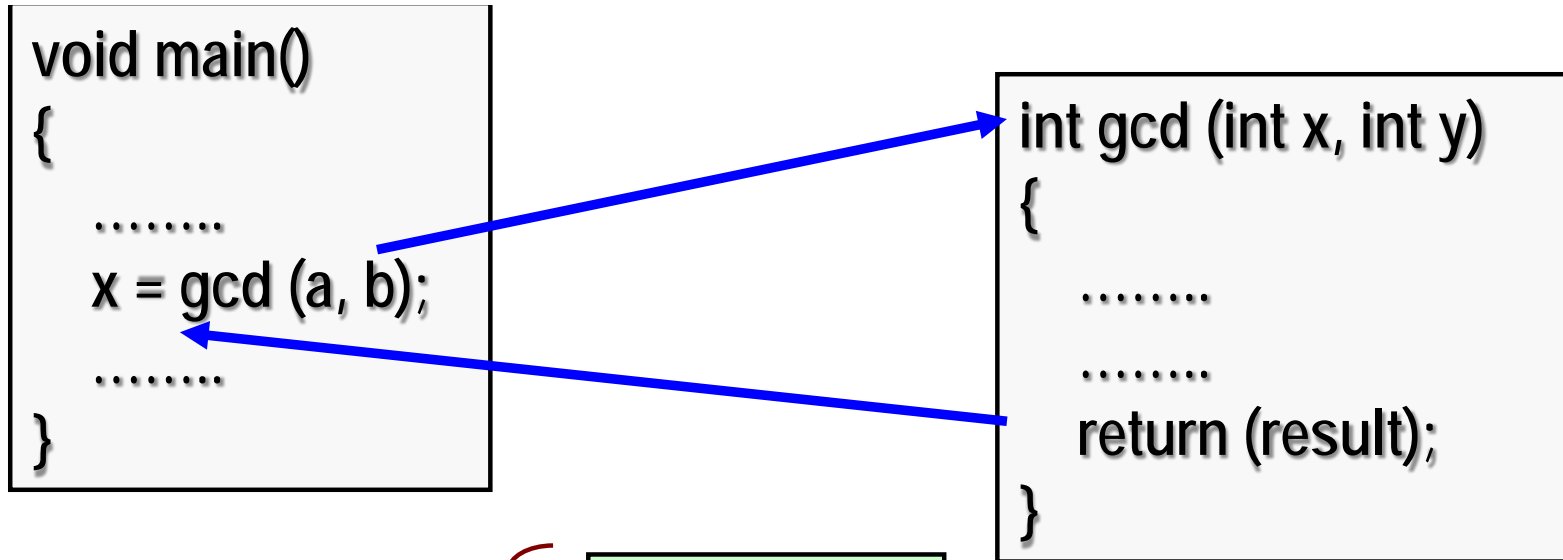
Characteristics of Recursion

- Every recursive program can also be written without recursion
- Recursion is used for programming convenience, not for performance enhancement
- Sometimes, if the function being computed has a nice recurrence form, then a recursive code may be more readable

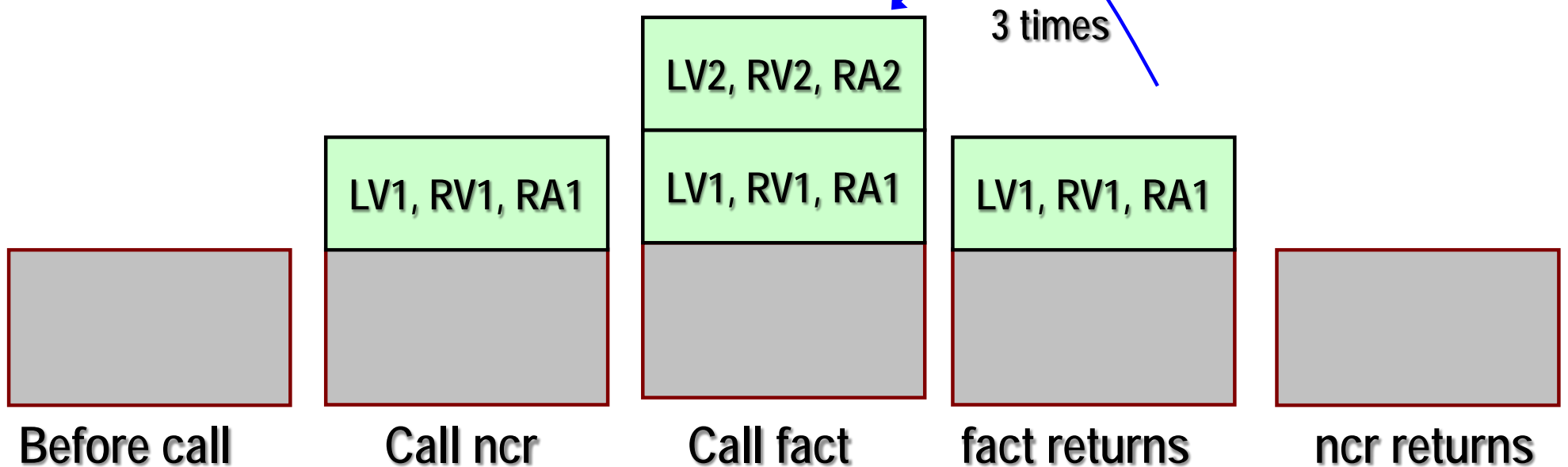
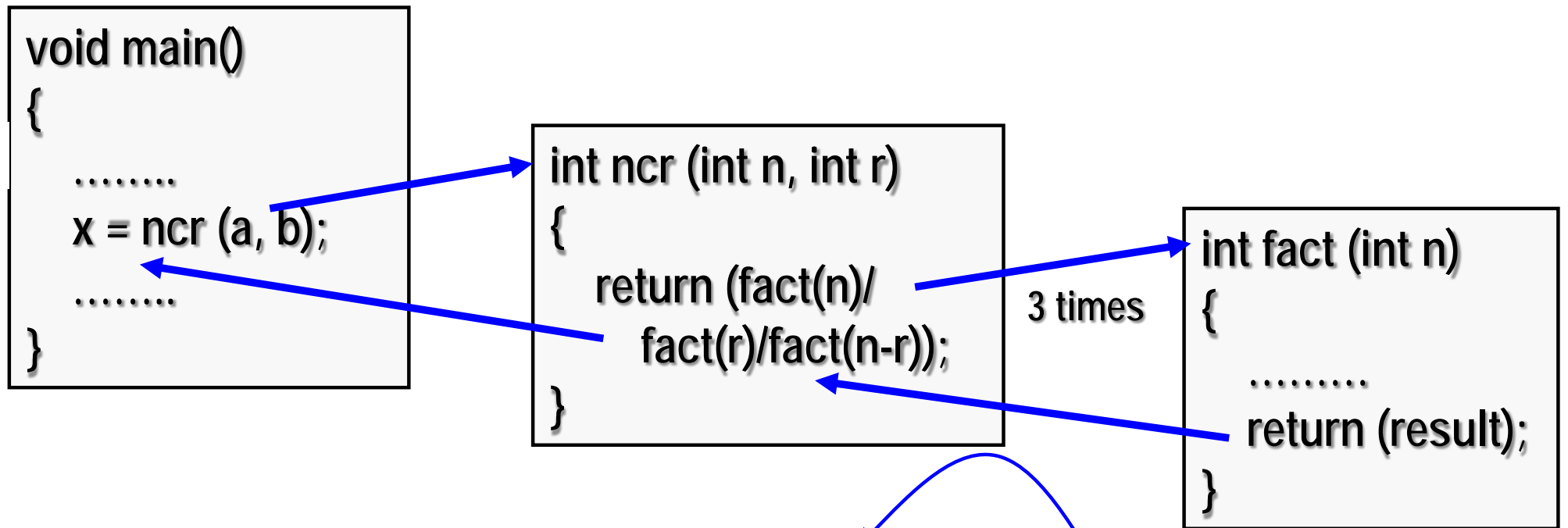
How are function calls implemented?

- The following applies in general, with minor variations that are implementation dependent
 - The system maintains a stack in memory
 - Stack is a last-in first-out structure
 - Two operations on stack, push and pop
 - Whenever there is a function call, the activation record gets pushed into the stack
 - Activation record consists of the return address in the calling program, the return value from the function, and the local variables inside the function

Memory Organization for Function Calls



Memory Organization for Recursive Calls



What happens for recursive calls?

- What we have seen
 - Activation record gets pushed into the stack when a function call is made
 - Activation record is popped off the stack when the function returns
- In recursion, a function calls itself
 - Several function calls going on, with none of the function calls returning back
 - Activation records are pushed onto the stack continuously
 - Large stack space required
 - Activation records keep popping off, when the termination condition of recursion is reached
- We shall illustrate the process by an example of computing factorial
 - Activation record looks like:

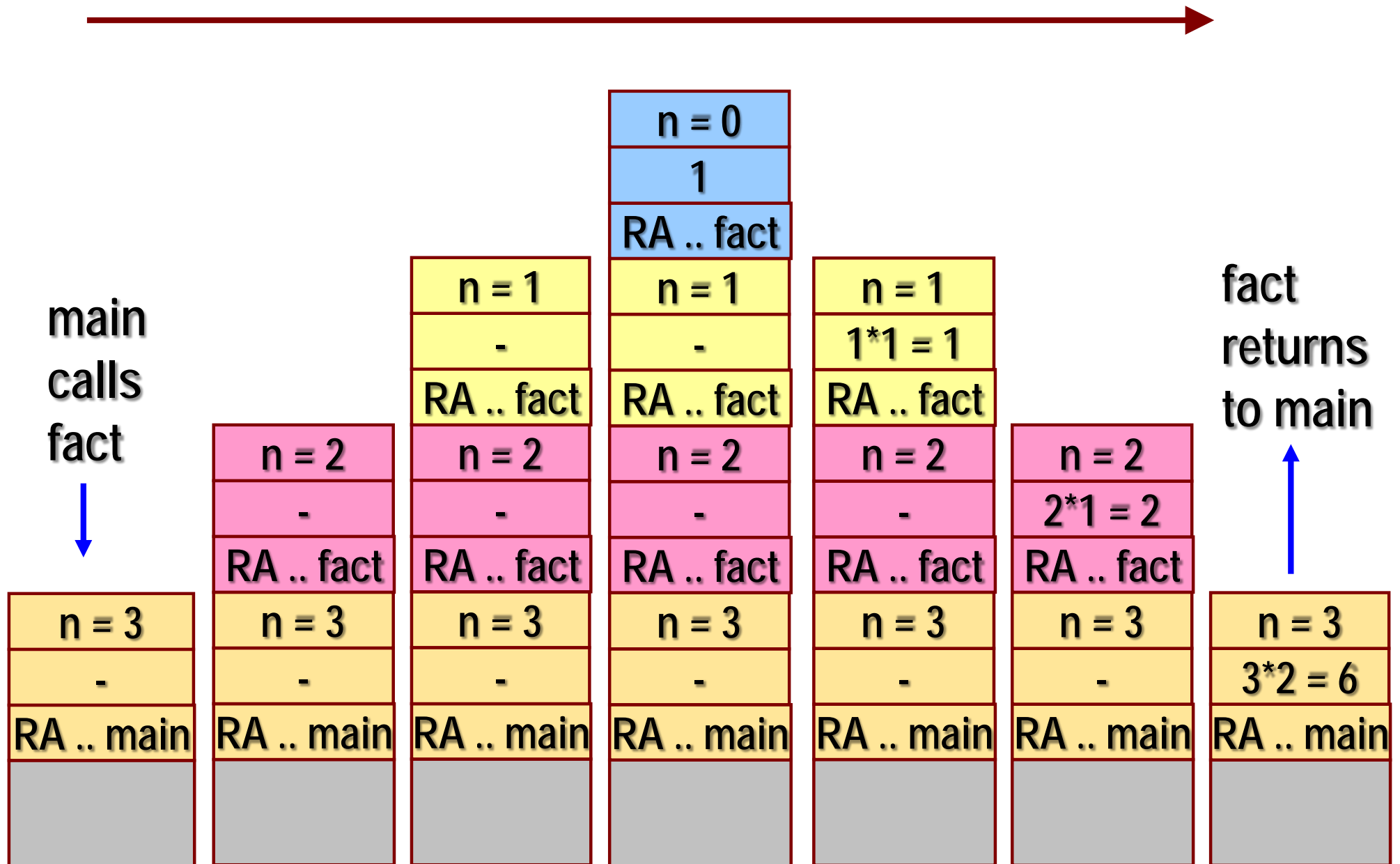
Local Variables
Return Value
Return Addr

Example:: main() calls fact(3)

```
void main()
{
    int n;
    n = 3;
    printf ("%d \n", fact(n) );
}
```

```
int fact (n)
int n;
{
    if (n == 0)
        return (1);
    else
        return (n * fact(n-1));
}
```

TRACE OF THE STACK DURING EXECUTION



Do Yourself

- Trace the activation records for the following version of Fibonacci sequence

```
int f (int n)
{
    int a, b;
    if (n < 2) return (n);
    else {
        X → a = f(n-1);
           b = f(n-2);
        Y → return (a+b);
    }
}

void main() {
    main → printf("Fib(4) is: %d \n", f(4));
}
```

Local Variables (n, a, b)
Return Value
Return Addr (either main, or X, or Y)



Thank You!