

TUTORIAL – 5

(CONTEXT-FREE LANGUAGES & PUSH-DOWN AUTOMATA)

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Problem-1

What context-free languages will be generated by the following two (separate) CFGs,

$G = \langle \{S, A, B\}, \{a, b, c\}, P, S \rangle,$

where P consists of the following production rules?

(a) $S \rightarrow ASB \mid \varepsilon, \quad A \rightarrow a, \quad B \rightarrow bb \mid b$

(b) $S \rightarrow abScB \mid \varepsilon, \quad B \rightarrow bB \mid b$

Problem-2

Define the context-free grammars for the following context-free languages?

Are your defined CFGs ambiguous / non-ambiguous?

- $L = \{ a^i b^j c^k \mid i, j, k \geq 0, \text{ and } i=j \text{ or } i=k \}$
- $L' = \{ a^i b^j c^k \mid i, j, k \geq 0, \text{ and } i+j = k \}$

Problem-3

Which of the following language(s) is/are context-free?

Give justifications.

- $L_1 = \{ a^m b^n \mid m, n \geq 0, m = 2n \}$
- $L_2 = \{ a^m b^n \mid m, n \geq 0, m \neq 2n \}$
- $L_3 = \{ a^m b^m c^{m+n} \mid m, n \geq 1 \}$
- $L_4 = \{ a^m b^n c^{m+n} \mid m, n \geq 1 \}$
- $L_5 = \{ a^l b^m c^n \mid l \geq 0, l < m \text{ and } l < n \}$

Problem-4

Consider the two CFGs G and G' with the start symbols S and S' and with the only productions:

Productions of G : $S \rightarrow aS \mid B, \quad B \rightarrow bB \mid b$

Productions of G' : $S' \rightarrow aA' \mid bB', \quad A' \rightarrow aA' \mid B', \quad B' \rightarrow bB' \mid \varepsilon$

Prove that, $L(G) \subset L(G')$, i.e., $L(G)$ is strictly contained in $L(G')$.

Problem-5

Prove that, the following context-free grammar,

$G = (\{S\}, \{a, b, c\}, P, S)$, is ambiguous.

Here, the production rules (P) are given as:

$$S \rightarrow aS \mid aSbS \mid c.$$

Construct a non-ambiguous grammar, G' , that derives the same language.

Also prove $L(G) = L(G')$.

Problem-6

Given the following languages over the alphabet $\{a, b\}$, design PDAs that accepts by empty stack (separate PDA for each one).

- $L = (a+b)^*b$
- $L' =$ all palindromes over $\{a,b\}$

Since L is a regular languages, can you directly present the left linear and the right linear grammar for L and then formally derive the NFAs from these grammars?



THANK YOU !

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