

Finite Automata and Regular Languages

1. Consider the following language over the alphabet $\Sigma = \{a, b\}$:

$$L_1 = \{x \in \{a, b\}^* \mid x \text{ starts with } ab \text{ but does not end with } ab\}.$$

- (a) Write a regular expression for L_1 .
(b) Design a DFA for L_1 .

2. Construct a regular expression over the alphabet $\Sigma = \{a, b, c\}$ for:

$$L_2 = \{x \in \{a, b, c\}^* \mid x \text{ has } 4i + 1 \text{ } b\text{'s for some integer } i \geq 0\}.$$

3. The language $L_4 = \{uvv^r w \mid u, v, w \in \{a, b\}^+\}$ is regular. Here v^r is the reverse of v . Answer the following questions.

- (a) Design a regular expression whose language is L_4 .
(b) Convert the regular expression of Part-(a) to an equivalent NFA.
(c) Convert the NFA in Part-(b) into an equivalent DFA using *subset-construction procedure*.
(d) Optimize the number of states of the DFA obtained in Part-(c) to produce the minimum state DFA.

4. Use pumping lemma to prove that the following languages are non-regular:

$$\begin{aligned} \bullet L_{4a} &= \{a^{k^3} \mid k \geq 0\}. & \bullet L_{4c} &= \{a^p \mid p \text{ is prime}\}. \\ \bullet L_{4b} &= \{a^{n!} \mid n \geq 0\}. & \bullet L_{4d} &= \{a^i b^j a^{ij} \mid i, j \geq 0\}. \end{aligned}$$

5. Determine the regularity/non-regularity of the following languages:

$$\begin{aligned} \bullet L_{5a} &= \{x \in \{a, b\}^* \mid \#a(x) - \#b(x) = 2021\}. \\ \bullet L_{5b} &= \{x \in \{a, b\}^* \mid \#a(x) - \#b(x) \text{ is a multiple of } 2021\}. \end{aligned}$$

where, $\#a(x)$ and $\#b(x)$ denote the number of a 's and b 's in x .

[*Hint:* Instead of finding regular expressions or using pumping lemma, you may alternatively try using the MyhillNerode theorem and show whether the number of equivalent classes formed is finite/infinite.]

6. Let A, B be languages over an alphabet Σ , and $C = A - B$. Then, justify which of the following statements is/are true?

- (a) If A and B are regular, then C is regular.
(b) If A and C are regular, then B is regular.
(c) If B and C are regular, then A is regular.
(d) If C is regular, then A and B are regular.

7. Two regular expressions over the same alphabet are called *equivalent* if they generate the same language. Prove/Disprove the equivalence of the following pairs of regular expressions over $\{a, b\}$.

- (a) $(ab + a)^* a$ and $a(ba + a)^*$.
(b) $(ab^* a + ba^* b)^*$ and $(ab^* a)^* + (ba^* b)^*$.