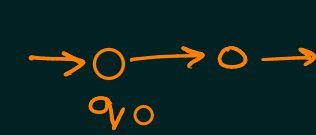


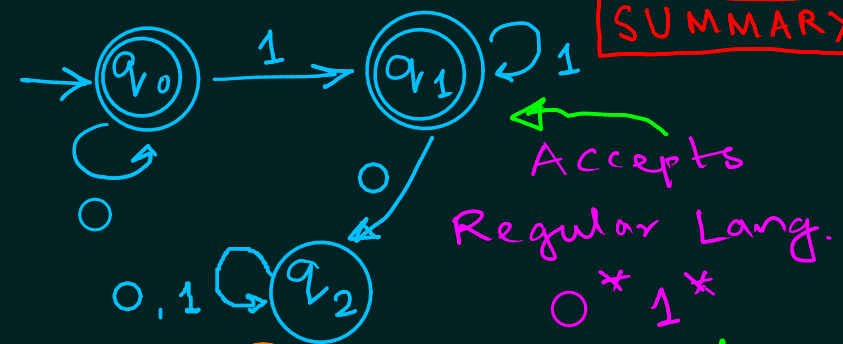
Deterministic Finite Automaton (DFA):

$\hookrightarrow \langle Q, \Sigma, q_0, \delta, F \rangle$
 Set of states $\rightarrow Q$, alphabet $\rightarrow \Sigma$, start state $\rightarrow q_0$, transition function $\rightarrow \delta$, Final states $\rightarrow F$
 $q_0 \in Q$, $F \subseteq Q$
 $\delta: Q \times \Sigma \rightarrow Q$

Accept Condition: $w \in L$ if $\hat{\delta}(q_0, w) = q_f \in F$



SUMMARY



Accepts Regular Lang.
 0^*1^*

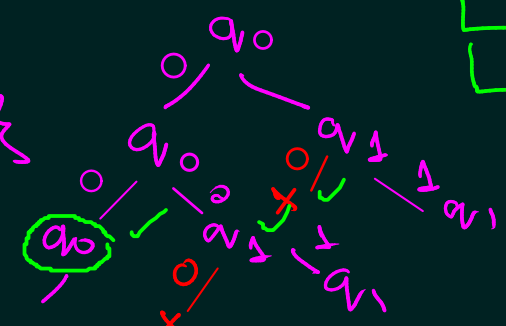
Non-deterministic Finite Automata (NFA):

$\hookrightarrow \langle Q, \Sigma, q_0, \delta, F \rangle$ where $\delta: Q \times \Sigma \rightarrow 2^Q$ (Power states of Q)

Accept condition: $w \in L$ if $\hat{\delta}(q_0, w) = \{q_1, q_2, \dots, q_n\}$ and $\exists q_i \in F$

at least one path leads to final states

\hookrightarrow no transition
 \hookrightarrow multiple transitions

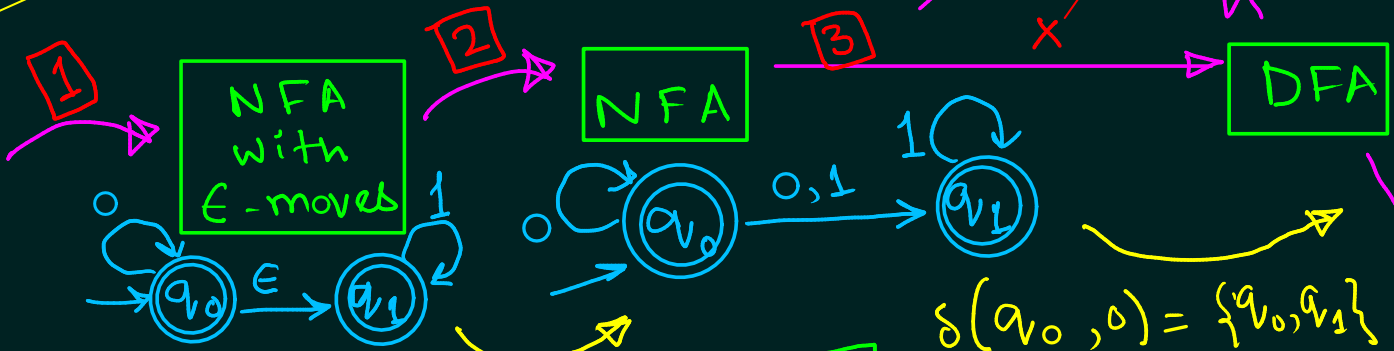


Regular Expression

0^*1^*

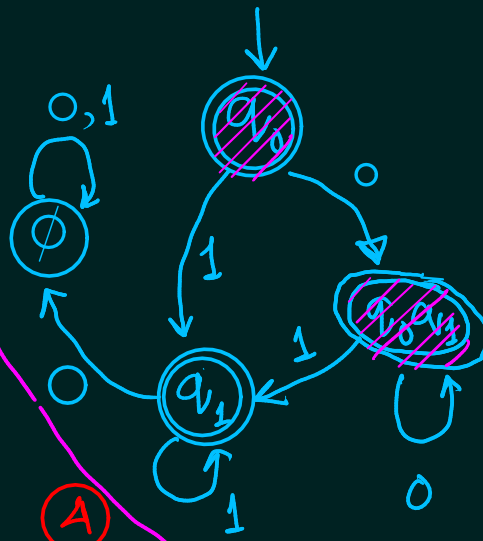
$\rightarrow \epsilon$ -closure(q_0) = $\{q_0, q_1\}$

$[0^*\epsilon = 0^*, \epsilon 1^* = 1^*]$

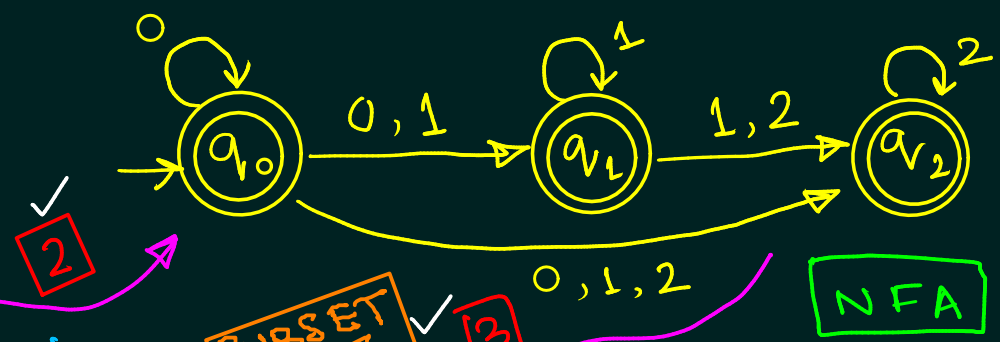
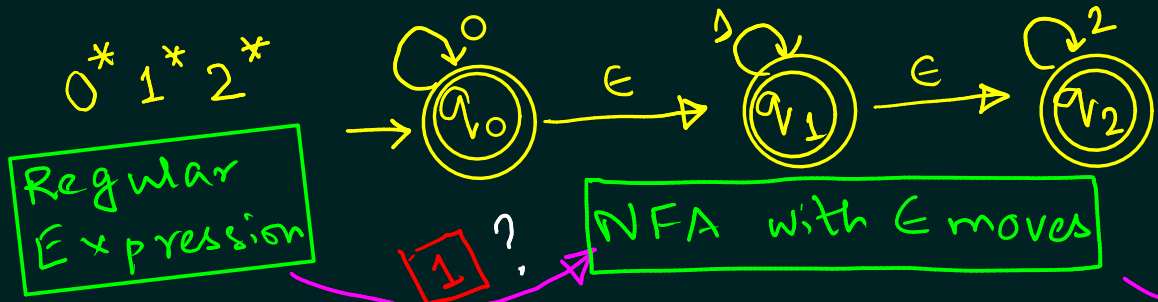


$$\delta(q, a) = \bigcup_{q' \in Q} \delta(q, a)$$

$\delta(q_0, 0) = \{q_0, q_1\}$
 $\delta(q_0, 1) = \{q_1\}$
 $\delta(\{q_0, q_1\}, 0) = \{q_0, q_1\}$
 $\delta(\{q_0, q_1\}, 1) = \{q_1\}$



Minimized DFA
 $q_0 = q_0, q_1$



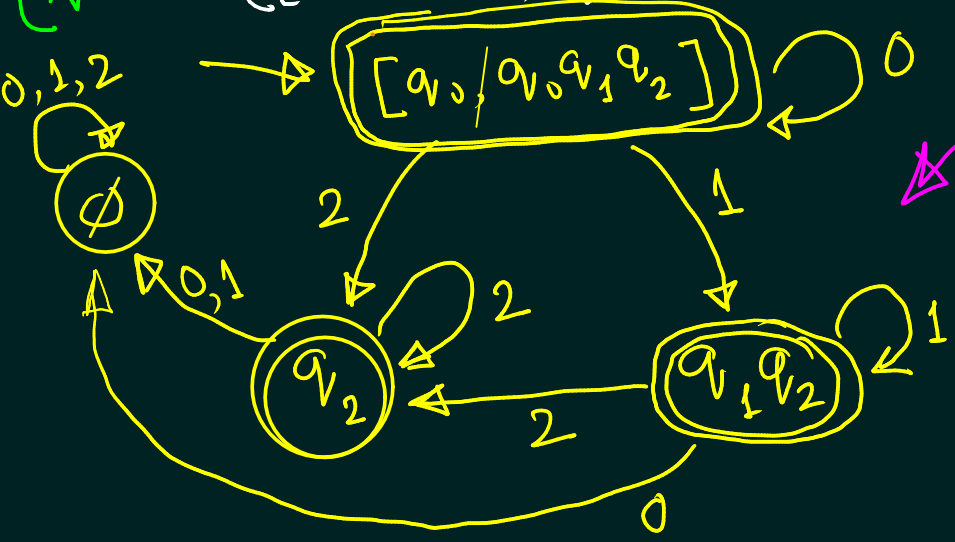
Redundant States $\{ [q_0q_2], [q_0q_1], [q_1] \}$
 [unreachable from q_0]

Equivalent states:

$$[q_0] \equiv [q_0q_1q_2]$$

Why: $\delta([q_0], w) = \delta([q_0q_1q_2], w)$

$$\text{OR } \delta([q_0], a) = [p] \\ \text{AND } \delta([q_0q_1q_2], a) = [p] \Rightarrow [p] \equiv [q]$$



Unique (Myhill-Narode Thm)

Regular Expression

Other Closure Properties.

TODAY'S LECTURE

① Regular Expression

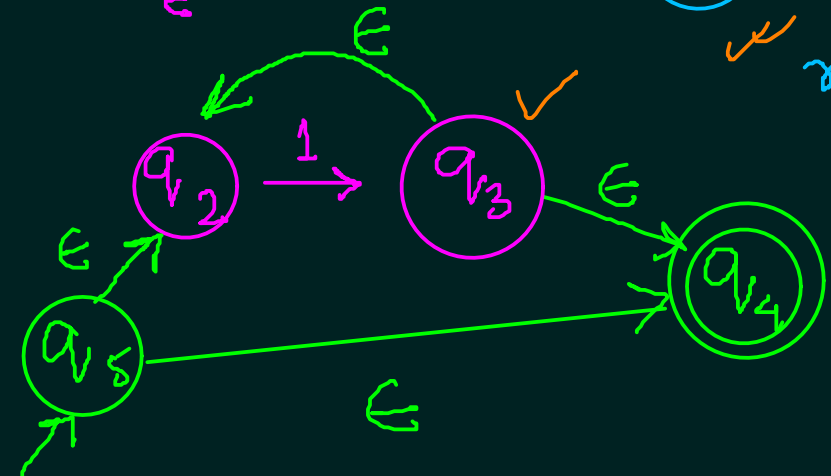
NFA with ϵ -Moves \rightarrow NFA

$(01^*+1) = \{0\underbrace{1}_{RL}, 01, 011, 0111, 01111, \dots\}$
 RE ✓
 $w \in L?$

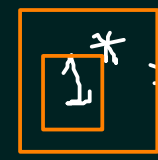
(01^*+1)



$\gamma = 0$



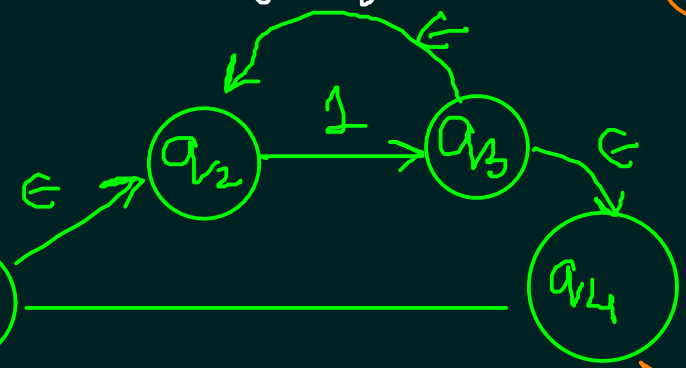
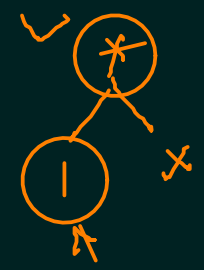
$\gamma = 1^*$



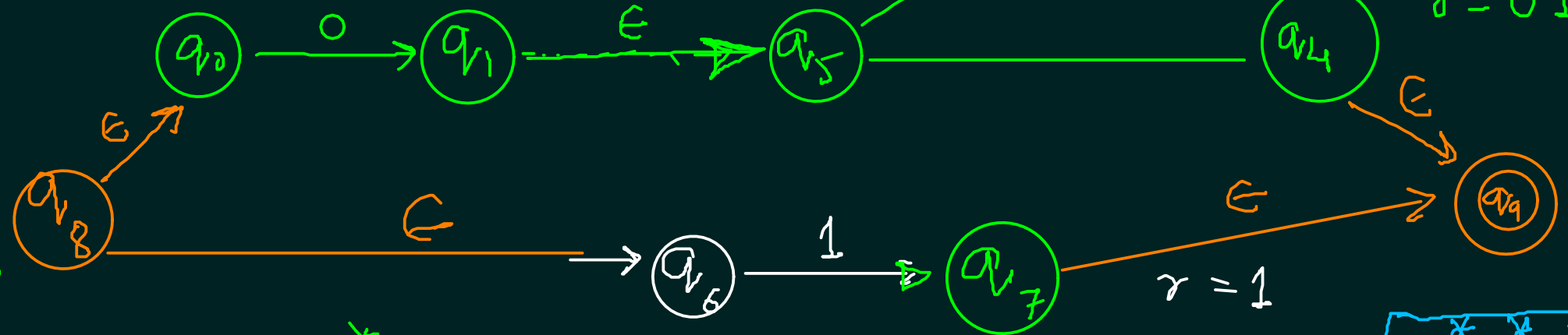
$= \epsilon + 1^+$
 $= \epsilon + \epsilon 1 \epsilon$

$+ \epsilon (1\epsilon)^* \epsilon$

$(01^*+1) + 0(11)^*$

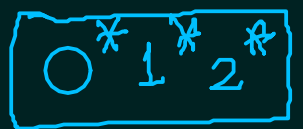


$\gamma = 01^+$



$\gamma = 01^*+1$

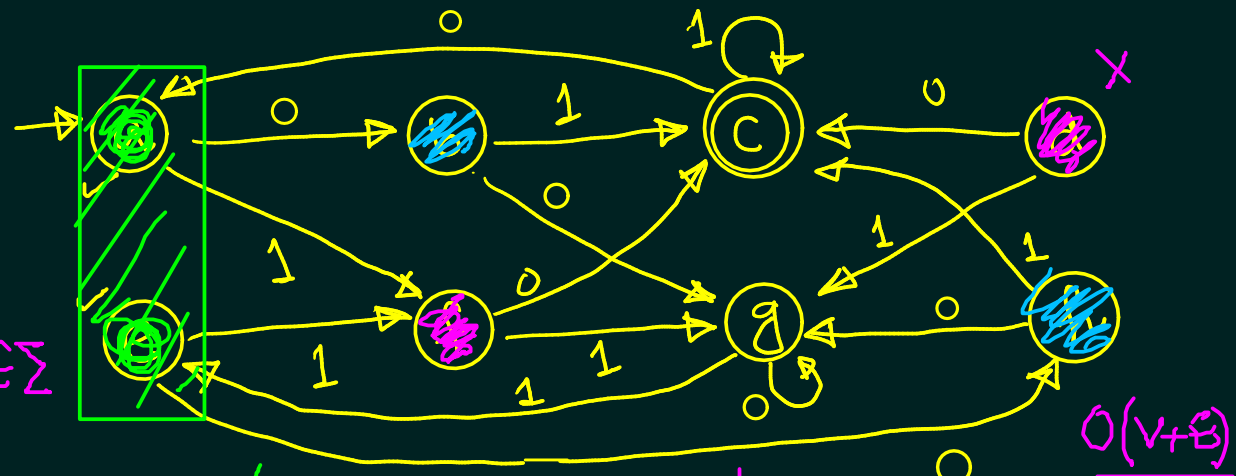
$\gamma = 1$



DFA → Minimized DFA

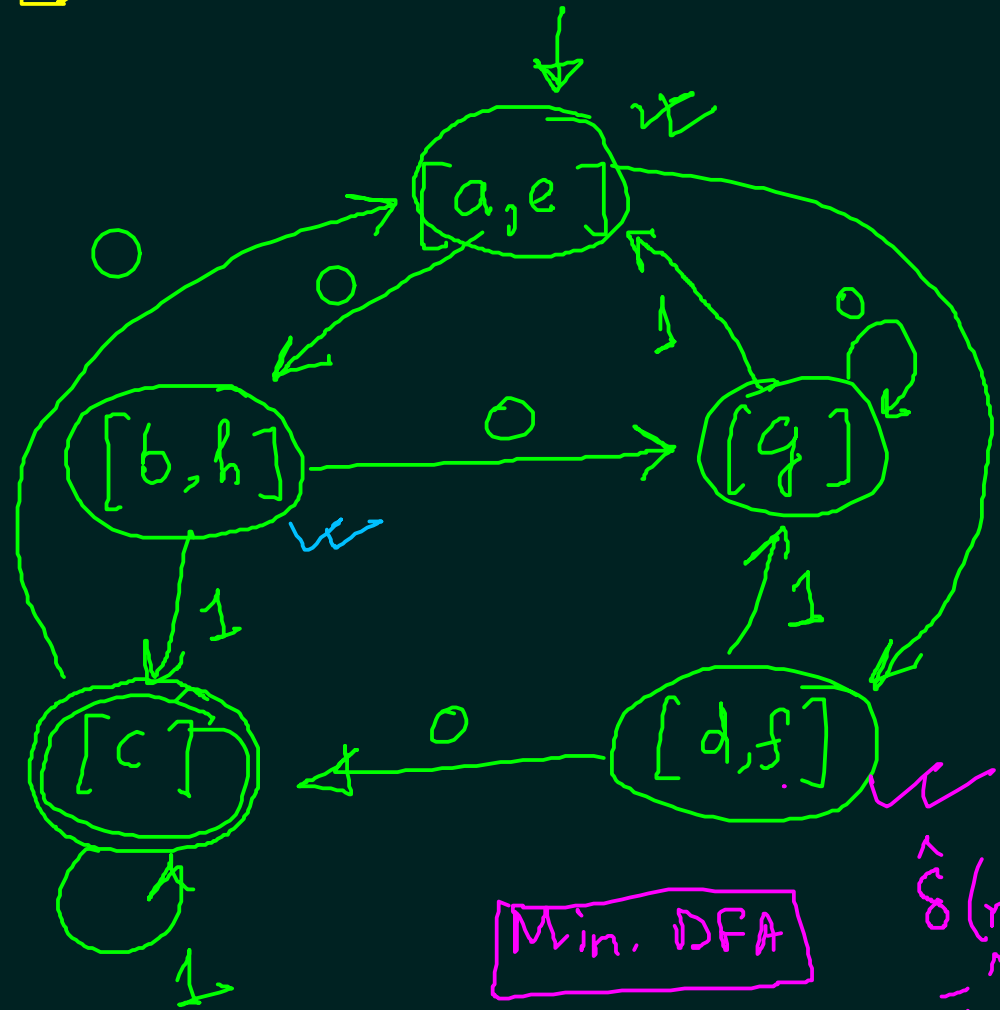
$[p] \equiv [q]$ iff

if $\delta([p], a) = \delta([q], a), \forall a \in \Sigma$



$O(V+E)$

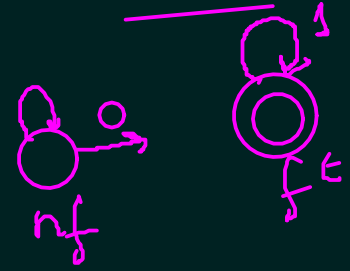
$+ k'(n+2n)$



- * Remove Red -
- * Merge Equiv.

a							
b	x						
c	x	x					
d	x	x	x				
e	hatched	x	x	x			
f	x	x	x	hatched	x		
g	x	x	x	x	x	x	
h	x	hatched	x	x	x	x	x
	a	b	c	d	e	f	g

$O(n^2)$



$\hat{\delta}(nf, w) = \hat{\delta}(f, w) \circ [kn^2 - k(n-1)^2]$

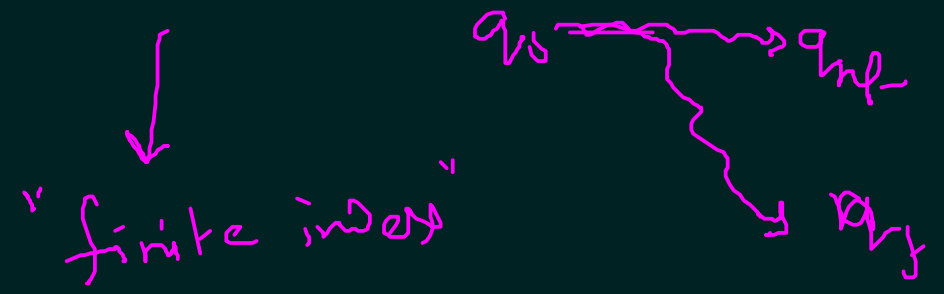
Myhill-Nasode Th. :

$$\boxed{x R y \text{ iff } \delta(q_0, x) = \delta(q_0, y)}$$

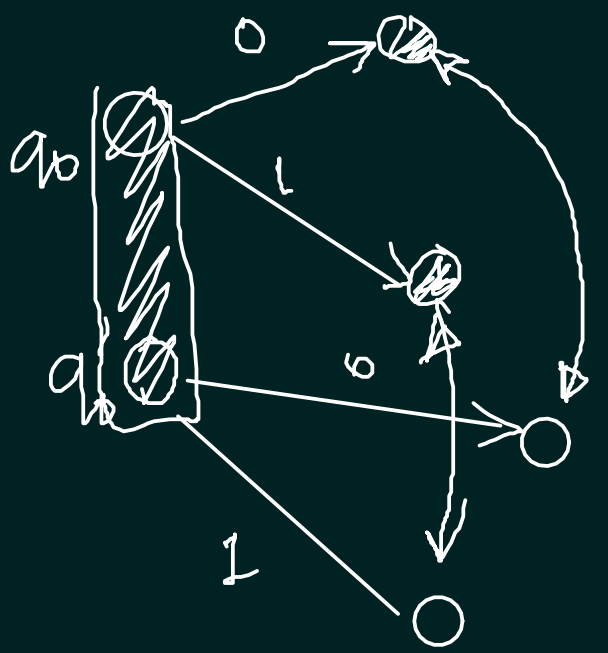
$\{x_1, x_2, \dots\} \rightarrow \text{Equiv. } \checkmark$

$\{[x], [x'], [x''] \dots\}$

algeb. mb of min.



min
0 1



$$[0w] \equiv [0w]$$

quotient
Set
Aut.

$$\left[\begin{array}{l} z \in \Sigma^* \quad x R y \\ xz \in L \Rightarrow yz \in L \end{array} \right]$$

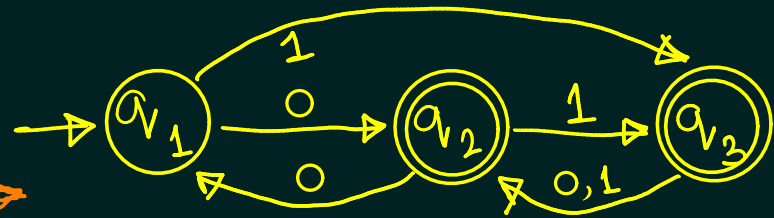
DFA \rightarrow Regular Expression



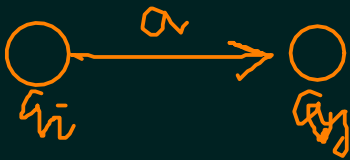
$$R = \underbrace{\gamma_{12}^3} + \underbrace{\gamma_{13}^3}$$



$$\gamma_{ij}^0 = \begin{cases} a & \delta(q_i, a) = q_j, i \neq j \\ \{a \mid \delta(q_i, a) = q_i\} & i = j \\ \cup \{ \epsilon \} \end{cases}$$



Min DFA



$$\gamma_{ij}^k = \underbrace{\gamma_{ik}^{k-1} (\gamma_{kk}^{k-1})^* \gamma_{kj}^{k-1}}_{k \text{ th}} + \underbrace{\gamma_{ij}^{k-1}}_{1 \text{ th}}$$

$$\begin{aligned} \gamma_{22}^{k=1} &= \gamma_{21}^0 (\gamma_{11}^0)^* (\gamma_{12}^0) + \gamma_{22}^0 \\ &= \emptyset \underline{\epsilon^*} \emptyset + \underline{\epsilon} \end{aligned}$$

$$\gamma_{32}^0 = \emptyset + 1 \quad \gamma_{11}^0 = \epsilon$$

	k=0	k=1	k=2
γ_{11}^k	ϵ	ϵ	$(00)^*$
γ_{12}^k	0	0	$0(00)^*$
γ_{13}^k	1	1	0^*1^*
γ_{21}^k	0	0	$0(00)^*$
γ_{22}^k	ϵ	$\epsilon + 00$	$(00)^*$
γ_{23}^k	1	$1 + 01$	0^*1
γ_{31}^k	\emptyset	\emptyset	$(0+1)(00)^*0$
γ_{32}^k	$0+1$	$0+1$	$(0+1)(00)^*$
γ_{33}^k	ϵ	ϵ	$\epsilon + (0+1)0^*1$