

1. One of the following sets is regular, the other is not. Justify.

$$A = \{ x\$y \mid x, y \in \{a,b\}^* \text{ and } \#a(x) = \#b(y) \} \subseteq \{a,b,\$ \}^*$$

$$B = \{ xy \mid \#a(x) = \#b(y) \} \subseteq \{a,b\}^*$$

Prove that \dagger in decomposition is unique

A is not regular. — Suppose A is regular. k a PLC for A

$$u = \epsilon, v = a^k, w = \$b^k$$

B is regular

$$u = \epsilon, v = a^k, w = b^k$$

$$B \subseteq \{a,b\}^*$$

$$a^{k-l} b^k = (a^{k-l} b^l) b^{k-l}$$

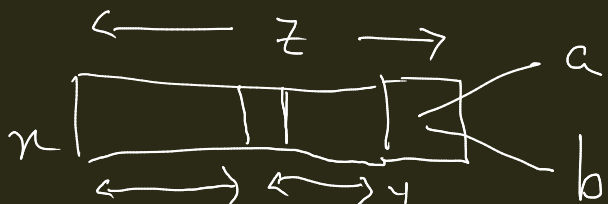
$$\supseteq a^k \quad 0, 1, 2, \dots, k$$

$$a^{k+l} b^k = a^k (a^l b^k)$$

Take $z \in \{a,b\}^*$. $\exists ST \quad z \in B$.

$$|z| = 0 \quad z = \epsilon = \epsilon \epsilon$$

Induction on $|z|$



$$\checkmark \quad y = \epsilon \Rightarrow x = b^k, k \geq 0$$

$$y \neq \epsilon$$

2. Let L be a regular language. Define $\text{dup}(L) = \{ ww \mid w \in L \}$. Which of the following statements is true?

- (a) $\text{dup}(L)$ must be regular.
- (b) $\text{dup}(L)$ must be non-regular.
- ✓ (c) $\text{dup}(L)$ may or may not be regular.

↓

L finite

$L = \mathcal{L}(a^*)$

$\text{dup}(L) = \mathcal{L}(aa)^*$

→

$$L = \mathcal{L}((a+b)^*)$$

PL proof
Take

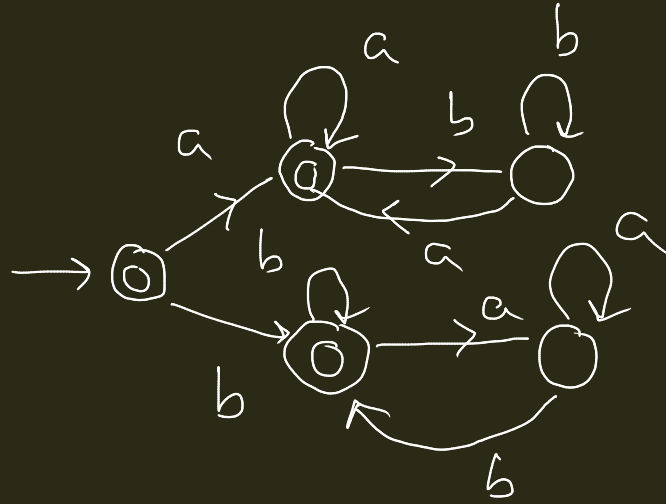
$$\begin{array}{ccc} a^k & b^k & a^k & b^k \\ \hline u & v & w & \end{array}$$

3. Regular or not? Justify.

$A = \{ w \in \{a,b\}^* \mid w \text{ contains an equal number of occurrences of the substrings } ab \text{ and } ba \}$

$B = \{ w \in \{a,b\}^* \mid w \text{ contains an equal number of occurrences of the substrings } aab \text{ and } bba \}$

A is regular



B is not regular

$$u = (aab)^{k+1} bba$$

$$v = (bba)^k \quad w = \epsilon$$

$$u = (aab)^k, \quad v = (bba)^k, \quad w = \epsilon$$

$$\underbrace{aab \ aab \ \dots \ aab}_{u} \ \underbrace{bba \ bba \ \dots \ bba}_v$$

$x = \epsilon, \quad y = b$

4. Let $A \subseteq \mathbb{N}$ be a set of positive integers. Define

$\text{binary}(A) = \{\text{binary representations of elements of } A\} \subseteq \{0,1\}^*$, and

$\text{unary}(A) = \{0^n \mid n \in A\} \subseteq \{0\}^*$.

Prove/Disprove:

(a) If $\text{unary}(A)$ is regular, then $\text{binary}(A)$ is regular.

(b) If $\text{binary}(A)$ is regular, then $\text{unary}(A)$ is regular.

$$A = \{2, 3, 5, 7, \dots\}$$

$$\text{binary}(A) = \{10, 11, 101, 111, \dots\}$$

$$\text{unary}(A) = \{010, 0010, \dots, 00, 000, 00000, 0000000, \dots\}$$

$$(b) A = \{2^n - 1 \mid n \geq 1\}$$

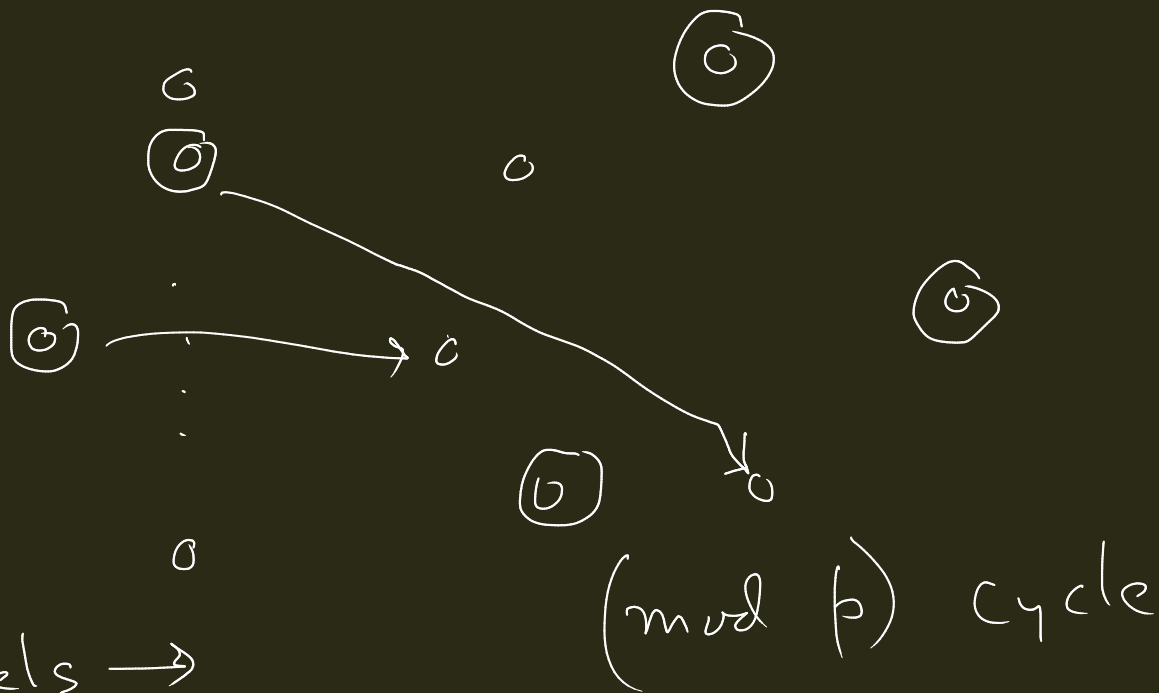
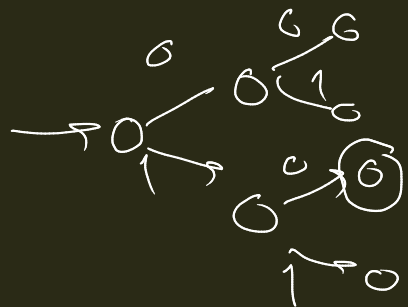
False
 $\text{binary}(A) = \{1, 11, 111, 1111, \dots\}$

$\text{lengths}(\text{unary}(A))$ is not u.p.
 $= A$

(a) True A is u.p.

$$A = \{a_1, a_2, \dots, a_k\} \cup$$

$$\{n \equiv b_1, b_2, \dots, b_\ell \pmod{\beta}, n \geq n_0\}$$



← n_0 levels →

5. Prove that no infinite subset of $\{ a^n b^n \mid n \geq 0 \}$ is regular.

$$\begin{array}{c} \downarrow \\ \frac{a^k}{v} \quad \frac{b^k}{w} \\ u = \epsilon \end{array}$$

$$\begin{array}{c} \downarrow \\ \frac{a^{k'}}{v} \quad \frac{b^{k'}}{w} \\ u = \epsilon \end{array}$$

for $k' \geq k$

6. Prove/Disprove: There exists a language L such that no infinite subset of L or its complement is regular.

$$L = \{a^n b^n \mid n \geq 0\}$$

\overline{L} ?

$\sim L$ contains $\mathcal{L}(a^+)$, $\mathcal{L}(b^+ a^+)$

L $\sim L$ $\mathcal{L}((ab)^+)$

0

1

2-3

4-7

8-15

16-31

32-63

64-127

7. Let D be a DFA with exactly k states and with $L = \mathcal{L}(D)$ infinite. Prove that L must contain a string of length in the range $[k, 2k)$. Give an example of D such that the shortest string in L of length $\geq k$ is of length $2k - 1$.

See Q4 of <http://cse.iitkgp.ac.in/~abhij/course/theory/FLAT/Spring20/ct1.pdf>

