

## Pumping lemma

Let  $L$  be a regular language. Then, there exists a positive constant  $k$  such that given any string  $uvw \in L$  with  $|v| \geq k$  we have the decomposition  $v = xyz$  satisfying

(1)  $y$  is non-empty

(2)  $|y| \leq k$

(3)  $u x y^i z w \in L$  for all  $i \geq 0$ .

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$w = a^{\lceil k/2 \rceil} b^{\lceil k/2 \rceil}$  — 3 cases

$w = \begin{pmatrix} a^k & b^k \\ v & w \end{pmatrix}$  — 1 case  
 $u = \epsilon$

introduction of  $u$  and  $w$   
forces pump in/out  
in  $v$   
may  $\Rightarrow$  proof simplification

Note: while using the lemma, forget about the DFA.

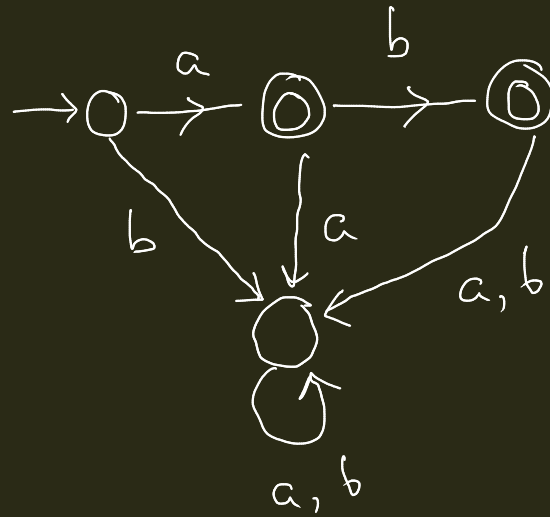
## Game with the demon

Demon

You

1. Assume  $L$  is regular.
2. Declare  $k$
3. Prepare  $uvw \in L$  with  $|v| \geq k$ .
4.  $v = xyz$ ,  
 $1 \leq |y| \leq k$   
 $uxy^i z w \in L$  for all  $i \geq 0$
5. choose  $i$  for a contradiction

$$L = \{a, ab\}$$



# of states is 4.

Demon given you  $k = 3$ .



Example 3

$$L_3 = \{ a^m b^n \mid m, n \geq 0, m \neq n \}$$

Suppose  $L_3$  is reg. Let  $k$  be a PLC for  $L_3$ .

$$u = \epsilon, \quad v = a^k, \quad w = b^{k+1}$$

$$v = xyz \quad 1 \leq |y| \leq k \quad uxy^i z w \in L \quad \forall i \geq 0$$

$$i = 2$$

$$a^{k+l} b^{k+1} \in L$$

\* No control on  $y$  or  $|y|$



Take  $l = 1$

$$u = \epsilon, \quad v = a^k, \quad w = b^{k+l}$$

Replay with demon

\* Do not think about the DFA. ~~X~~

$$v = x'y'z'$$

$$1 \leq |y'| \leq k = l'$$

Demon does not guarantee  $l' = l$ .

