

1. Which of the following languages is/are context-free? (Give sufficient justifications.) (5×5)

(a) $\{a^i b^j c^k \mid i, j, k \geq 0 \text{ and } i + j = k\}$.

(b) The set of all concatenations of strings of the form $a^i b^j c^k$ with $i, j, k \geq 0$ and $i + j = k$, i.e., the language $\{a^{i_1} b^{j_1} c^{k_1} a^{i_2} b^{j_2} c^{k_2} \dots a^{i_l} b^{j_l} c^{k_l} \mid l, i_r, j_r, k_r \geq 0 \text{ and } i_r + j_r = k_r \text{ for all } r = 1, 2, \dots, l\}$.

(c) $\{(a^i b^j c^k)^l \mid i, j, k, l \geq 0 \text{ and } i + j = k\}$.

(d) $\{a^{il} b^{jl} c^{kl} \mid i, j, k, l \geq 0 \text{ and } i + j = k\}$.

(e) $\{a^{i+l} b^{j+l} c^{k+l} \mid i, j, k, l \geq 0 \text{ and } i + j = k\}$.

2. Consider the CFL $L := \{a^k b^l \mid l \geq 0 \text{ and } 0 \leq k \leq 2l\}$.

(a) We have discussed the following CFG for L :

$$S \rightarrow \epsilon \mid Sb \mid aSb \mid aaSb.$$

Show that this grammar is ambiguous. (5)

(b) Design an unambiguous CFG for L . You must supply a proof that your grammar is unambiguous. (15)

3. An arithmetic expression like $((a+b) * (c-d) / e)$ is said to be in the infix notation, since each operator is placed between its operands. In the postfix notation the operator follows its two operands. Thus $a+b$ and $c-d$ in the postfix notation become $ab+$ and $cd-$ and $((a+b) * (c-d) / e)$ becomes $ab+cd-*e/$. It can be proved that the postfix notation does not require parentheses to resolve ambiguities.

(a) Write the postfix expression for $a + \frac{1}{b + \frac{1}{a + \frac{1}{b}}}$ (i.e., for $(a + (1 / (b + (1 / (a + (1 / b))))))$). Also convert the postfix expression $abab+/ab-*a/*b-$ to the equivalent infix expression. (5)

(b) Design a PDA whose language is the set of all valid postfix expressions involving two variables a and b and the binary operators $+$, $-$, $*$ and $/$. (10)

4. (a) Describe a one-tape one-head TM whose language is the set of all strings over $\{a, b\}$, that contain equal number of a and b . What is the maximum number of moves this TM requires for a string of length n ? (10)

(b) Show that using a scratch tape one can reduce this maximum number of moves. (10)

5. (a) Let L be an R.E. language over some alphabet Σ . Show that L is the language of some TM that does not halt on some input(s), if and only if $L \neq \Sigma^*$. (10)

(b) Show that R.E. languages are closed under union and intersection. (10)