

The complexity class DP is defined as:

$$DP := \{L_1 \setminus L_2 \mid L_1, L_2 \in NP\} = \{L_1 \cap L_2 \mid L_1 \in NP, L_2 \in \text{coNP}\},$$

i.e., $L \in DP$, if and only if L is the *difference* of two languages in NP, or equivalently, if and only if L is the intersection of a language in NP and a language in coNP.

The *difference hierarchy* is defined as

$$\begin{aligned} \Delta_0 P &:= P, \\ \Delta_i P &:= P^{\Sigma_{i-1} P} \text{ for } i \geq 1. \end{aligned}$$

1. Argue that the language

$$\text{HALFCYCLE} := \{\langle G \rangle \mid \text{The longest cycle in the directed graph } G \text{ is of length } \lfloor n(G)/2 \rfloor\}$$

is in DP. (Here $n(G)$ denotes the number of vertices of the graph G , and $\lfloor x \rfloor$ the floor of the real number x , i.e., the largest integer $\leq x$.) (4)

2. Demonstrate that $P = \Delta_1 P$ and $NP \cup \text{coNP} \subseteq DP \subseteq \Delta_2 P$. (4)

3. Prove that $DP = NP \cup \text{coNP}$, if and only if $NP = \text{coNP}$. (**Hint:** For proving the ‘only if’ part you may first show that HALFCYCLE is both NP-hard and coNP-hard — use poly-time reductions from HAMCYCLE and $\overline{\text{HAMCYCLE}}$.) (2+5)