Stacks, Queues, and Trees

In this assignment, you solve two unrelated problems. The first one is on the realization of a queue, and the second one is on binary trees. In order to reduce your programming overhead, a blackbox BB5 is provided which provides the following utilities.

- A complete implementation of the stack ADT.
- A constructor of random binary trees.

The details of the blackbox features will be explained in appropriate places. In order use the blackbox, write following line after your usual **#include** directives.

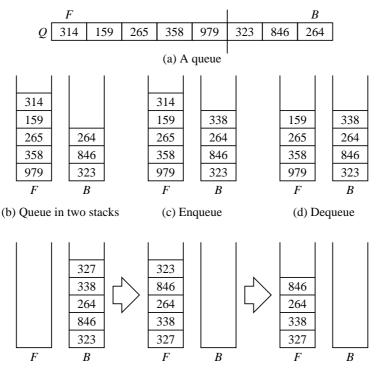
```
#include "BB5.h"
```

At the beginning of your main() function, call registerme() (this step is optional for this black box). Moreover, compile your code as:

gcc/g++ mycode.c/mycode.cpp BB5.0

Part 1: Implementation of a queue using two stacks

As the following figure illustrates, a queue Q (see Part (a) of the figure) can be realized from two stacks F and B (see Part (b)). An arbitrary break-point is chosen (between 979 and 323 in the figure). The part of Q before this break-point resides in the front stack F, and the part of Q after the break-point resides in the back stack B. Notice the order in which the elements of Q appear in F and B.



(e) Dequeue from an empty front stack

An enqueue operation involves pushing the new item to the back stack B (see Part (c)). A dequeue operation is the same as pop from the front stack F. If F was not empty before the pop, this is straightforward (see Part (d)). If both F and B are empty, then Q is empty too, and a dequeue from Q is not permitted. If F is empty but B contains one or more elements (see Part (e)), the element to dequeue lies at the bottom of B, and cannot be directly accessed. Make a sequence of pop operations from B and push operations of those elements to F, until B becomes empty. Now, a normal dequeue (pop from F) can be performed. In order to implement a queue this way, use the implementation of the **STACK** ADT from BB5. You may now define your queue as:

```
typedef struct {
   STACK F, B;
} QUEUE;
```

For a stack s, the ADT calls supplied by BB5 are tabulated below. Only stacks of integers are supported.

s = SINIT() Create an empty stack.	
ISEMPTY (S) Returns 1 or 0 depending or	n whether <i>S</i> is empty or not.
TOP (S) Returns the element (an inte	eger) at the top of S.
s = PUSH(s, x) Push an integer x to the stac	k <i>S</i> .
S = POP (S) Perform a pop operation fro	om <i>S</i> .
SPRNT2B(S) Print the elements of S from	top to bottom.
SPRNB2T (S) Print the elements of S from	bottom to top.

You do not need to understand the implementation of the **STACK** ADT (since your teachers did it, the operations are efficient). Use the above calls to implement the queue ADT as follows.

Q = QINIT()	Create an empty queue.
Q = ENQUEUE(Q, x)	Enqueue an integer x to Q .
Q = DEQUEUE(Q)	Perform a dequeue operation on Q.
QPRN (Q)	Print the elements of Q from front to back.

In your main () function, read a small integer n from the user. Start with an empty queue Q. Make n enqueue and n dequeue operations on Q. Enqueue randomly generated integers to Q. Never make an attempt to dequeue from an empty queue. Print the queue after each operation.

Part 2: Level-by-level listing of keys in a binary tree

The blackbox BB5 defines a binary-tree data type as follows. This is declared in BB5.h.

```
typedef struct _treenode {
    int key;
    struct _treenode *L, *R, *N;
} treenode;
typedef treenode *TREE;
```

By calling

T = TGEN(n);

you can construct a random binary tree on n nodes with random key values. A pointer to the root node is returned by this call. In order to print the tree in a human-readable format, use the following call.

TPRN(T);

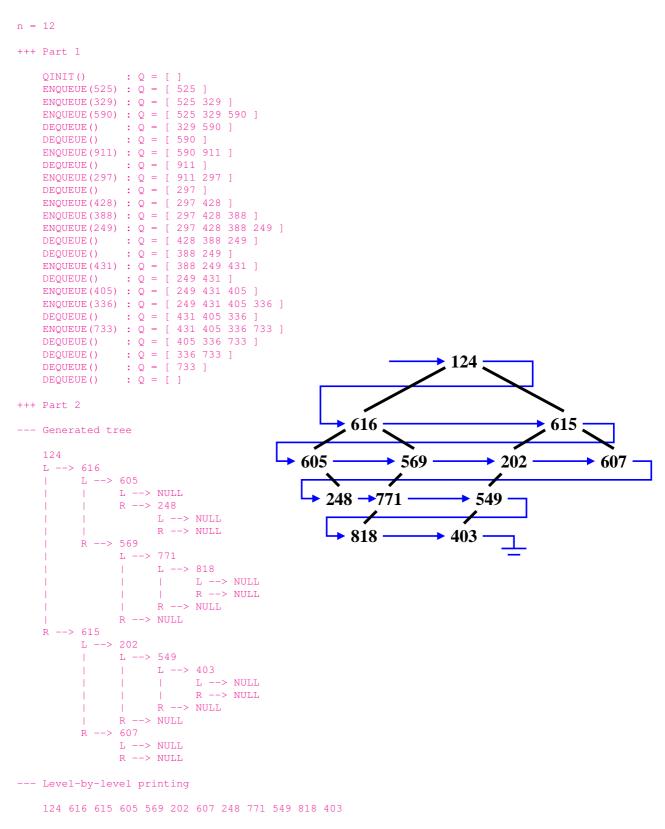
The tree constructor of BB5 keeps all the \mathbf{N} pointers NULL. Write a function

SETN(T)

to set these pointers as explained now. Let u be a node at level l in T. Impose a left-to-right ordering of the nodes in each level. If u is not the last node in T at level l, then the \mathbf{N} pointer of u should point to the next node at the same level l. If u is the last node at level l, then the \mathbf{N} pointer of u should point to the first node of the next level (or to NULL if l is the last level in T). Once the \mathbf{N} pointers are so adjusted, T becomes a linked list with respect to these pointers starting at the root and storing the nodes in a level-by-level and left-to-right-in-each-level order. To implement this function, make a recursive traversal of T, and build a linked list of nodes at *each* level. After the traversal completes, join the level-wise lists into a single list. Write another function **TPRNL (T)** that makes a level-wise printing of the nodes of T. Use a linked-list print procedure following the \mathbf{N} pointers.

In the main () function, read n (may be the same as in Part 1), build a tree on n nodes by **TGEN**, print the tree by **TPRN**, call **SETN** to set the **N** pointers, and then print the resulting linked list by calling **TPRNL**.

Sample output



Submit a single C/C++ source file. Do not use global/static variables.