- 1. Write a program to recognize balanced strings with the delimiters () { } []. For example, the string { (() []) } ([] [] ()) { } is balanced. Two strings that are not balanced are { (() []) } ([] [] [] () ] { } and { (() []) } ([] [] [] () { }. Use a stack.
- 2. You are given integers n, k with 0 ≤ k ≤ n. Your need to print all the subsets of size k of the set {1,2,3,...,n}. Use a FIFO queue in the following fashion. Insert the empty set Ø to an initially empty queue Q. Repeat so long as Q is not empty: Let S be the element (a set) at the front of Q. Dequeue S from Q. Let the size of S be l, and let the maximum element of S be m (take m = 0 if S = Ø). If l = k, print S. Otherwise (that is, if l < k), enqueue the (l+1)-element sets S ∪ {i} in Q for i = m+1, m+2, ..., n.</li>
- **3.** You are given a min-heap *H* of integers in the contiguous representation. Write a function that takes *H* and an integer *B* as input, and prints all the elements of the heap that are  $\leq B$ . The function must not modify *H*. The running time of your function should be O(k), where *k* is the number of integers printed.
- **4.** You are given *k* sorted arrays  $A_1, A_2, \ldots, A_k$  with a total of *n* elements. You are required to build an array *B* of size *n* by merging the *k* input arrays. You are allowed to use only O(k) additional space (in addition to the arrays  $A_i$  and *B*). Your program should run in  $O(n \log k)$  time. Use a min-priority queue to implement your algorithm. (**Hint:** At any point of time, the priority queue should store at most one element from each  $A_i$ .)
- 5. (a) A random binary tree *T* can be constructed as follows. Let *n* be the number of nodes in *T*. Randomly generate the numbers  $n_l$  and  $n_r$  of nodes in the left and right subtrees (so that  $n = 1 + n_l + n_r$ ). Recursively build the left and the right subtrees, whichever is/are non-empty. Store random keys at the nodes.

(b) Let *r* be the root and *v* a leaf node in *T*. There is a unique *r*, *v* path  $r = u_0, u_1, u_2, ..., u_l = v$  in *T*. Let the key stored at node  $u_i$  be  $k_i$ . The alternating sum of these key values is

$$altsum(v) = k_0 - k_1 + k_2 - \dots + (-1)^l k_l.$$

Write a function to print the alternating sums at all the leaf nodes in T.

**6.** Let *T* be a general rooted tree, in which each node can have any number of children.

(a) Let v be a node in T with c child nodes. In addition to a key, v stores an array of c child pointers (the count c should also be stored at v). The child pointers point to the c subtrees of v. Build a random general rooted tree using a constructor similar to that for binary trees in the last exercise.

(b) Write a function to compute the height of T.

(c) Write a function that, given the general tree T, prepares and returns a binary tree B which is the first-child-next-sibling representation of T.

- (d) Write a function to compute the height of T if B is supplied as the only input.
- 7. Write an O(h(T))-time function to compute the immediate successor of a node in a BST *T*.
- 8. A three-way search tree (3ST) is a rooted tree with each node storing two keys and three child pointers L, M, R. Let v be a node in the 3ST storing the keys  $k_1, k_2$ . Let l (resp. m, r) be a key value stored in the left (resp. middle, right) subtree. We must have  $l < k_1 < m < k_2 < r$ .
  - (a) Write a function to search for a key in a 3ST.
  - (**b**) Write a function to insert a key in a 3ST.
- **9.** (a) Write the insert function for binary search trees. Prepare a BST *T* by inserting random keys to an initially empty tree. Let *n* be the number of nodes in *T* after all these insertions.

(b) Convert *T* to a BST of height n-1 as follows. So long as the root has a left child, make a right rotation at the root. When the left subtree of the root becomes empty, move to the right child of the root, and repeat the process. (**Remark:** This is first stage of the Day–Stout–Warren (DSW) BST rebalancing algorithm.)

**10.** Write a function that, upon input *h*, returns an AVL tree of height *h* and with the minimum possible number of nodes. (These are called Fibonacci trees.)