

CS29003 Algorithms Laboratory

Assignment No: 9

Last date of submission: 12–October–2017

A CGPA is technically a floating-point number in the range $[5.00, 10.00]$. However, CGPAs are truncated (or rounded) two places after the decimal point. This means that if x is a CGPA, then $100x$ is an integer in the range $[500, 1000]$. Moreover, $100x - 500$ is an integer in the range $[0, 500]$. In this assignment, you deal with these integer versions of CGPAs. Your task is to sort an array A storing n integer-valued CGPAs.

Part 1: Your first task is to generate a random array A of n CGPAs. In a population of students, the floating-point CGPA x satisfies the following probability distribution: $\Pr[5 \leq x \leq 6] = 0.05$, $\Pr[6 < x \leq 7] = 0.2$, $\Pr[7 < x \leq 8] = 0.3$, $\Pr[8 < x \leq 9] = 0.3$, and $\Pr[9 < x \leq 10] = 0.15$. When we convert x to an integer $X \in [0, 500]$, we have the equivalent probabilities $\Pr[0 \leq X \leq 100] = 0.05$, $\Pr[100 < X \leq 200] = 0.2$, and so on. Assume that in each band, the different CGPA values are equally likely, that is, we have

$$\Pr[X = a] = \begin{cases} 0.05/101 & \text{if } 0 \leq a \leq 100, \\ 0.20/100 & \text{if } 101 \leq a \leq 200, \\ 0.30/100 & \text{if } 201 \leq a \leq 300, \\ 0.30/100 & \text{if } 301 \leq a \leq 400, \\ 0.15/100 & \text{if } 401 \leq a \leq 500. \end{cases}$$

You should generate the array A of n CGPAs following this probability distribution. The `rand()` library function returns a random integer in the range $[0, \text{RAND_MAX}]$. This function can be used to generate samples following other probability distributions.

In general, let X be a random variable that assumes integer values in the range $[0, m]$. Let $p_i = \Pr[X = i]$ for $i \in [0, m]$. Consider the cumulative probability $q_i = \sum_{j=0}^i p_j$. We generate a (uniformly) random floating-point number $y \in [0, 1]$ by dividing a `rand()` output by `RAND_MAX`. If $y \leq q_0$, we output the integer 0. Otherwise, we locate the $i \in [1, m]$ such that $q_{i-1} < y \leq q_i$, and output i . Notice that the cumulative-probability array $(q_0, q_1, q_2, \dots, q_m)$ is sorted in the ascending order, so the index i can be located by a binary search. The running time of this algorithm to generate the array A is $O(n \log m)$.

In the current situation, a somewhat more efficient ($O(n)$ -time) algorithm can be designed. Implement some algorithm to generate the array A of n integer-valued CGPAs following the distribution given above.

Part 2: Implement quick sort in a function `quicksort` for sorting integer arrays. Do not use the `qsort` library function defined in the standard C/C++ libraries. Partitioning should be inline, that is, no additional temporary arrays may be used.

Part 3: Write a function `countingsort1` to implement the standard stable counting sort algorithm for CGPA arrays. The function should use a count array C , and a temporary array B of size $n = |A|$ which is finally copied back to the input array A .

Part 4: In this particular application, we are sorting integer values in the range $[0, 500]$ (not records containing these integer values). Therefore the array B can be eliminated altogether. You still need the count array C . Implement this idea in a function `countingsort2`.

The `main()` function

- Read n from the user. For this assignment, n should be in the range $[10^4, 10^8]$.
- Call the function of Part 1 to generate a CGPA array A of size n with elements following the given probability distribution. Make copies of A in two separate arrays B and C .
- Run `quicksort` on A , and report the running time.
- Run `countingsort1` on B , and report the running time.
- Run `countingsort2` on C , and report the running time.

Sample output

```
n = 100000000
+++ Array generation time = 4.904727 sec
+++ Quick sort time      = 11.338450 sec
+++ Counting sort 1 time = 1.461599 sec
+++ Counting sort 2 time = 0.748255 sec
```

Submit a single C/C++ source file. Do not use global/static variables.