

Read an integer $n < 100$ and a string S of length n from the user. Write two functions `norepdecomp()` and `longestpalsubstr()` to solve the following parts.

Part I

Decompose S as a concatenation of strings S_1, S_2, \dots, S_k such that no S_i contains the same symbol in consecutive positions, and k is as small as possible. For example, for the input string `bbbcaacacabaccacbb`, your function should report $k = 8$ and the decomposition `b, b, bca, acacabac, c, cac, cb, b`.

Part II

Find longest palindromic substrings of S . If there are multiple palindromic substrings of the same maximum length, report all of them. For example, for the input $S = \text{bbbcaacacabaccacbb}$, you should report the four substrings `acaca`, `cabac`, `accca`, and `ccacc`. Your program must run in $O(n^2)$ time and use only $O(n)$ additional space. The obvious $O(n^3)$ -time algorithm of checking whether each of the $\Theta(n^2)$ substrings of S is palindromic will deserve no credit.

Hint: You need to efficiently compute which of the substrings of S are palindromic.

Sample Output

```
n = 50
S = bbbabbcbccbaaaaaccbcababccabccabcccaacbacabcbbabc
Part I: b b bab bcbc cba a a a ac cbcababc cab bcabc c ca acbacabcbbabc
Part II: Length = 7. Substrings: cbacabc abcbbca
```

Read an integer $n < 100$ and a string S of length n from the user. Write two functions `norepsubseq()` and `paldecomp()` to solve the following parts.

Part I

Find a longest subsequence T of S such that T does not contain the same symbol in consecutive positions. For example, for the input string `bcbcabacbacbbbaabb`, T should be `bcbcabacbacbbab`.

Part II

Decompose S as a concatenation of strings S_1, S_2, \dots, S_k such that each S_i is a palindrome, and k is as small as possible. For example, for the input $S = \text{bcbcabacbacbbbaabb}$, your function would report $k = 8$ and the possible decomposition `b, c, bcbacab, a, cc, bcb, aa, bbb`. There may be multiple decompositions having the same optimal k . Reporting any of the optimal decompositions suffices. Your function must run in $O(n^2)$ time and use only $O(n^2)$ additional space.

Hint: You need to efficiently compute which of the substrings of S are palindromic.

Sample Output

```
n = 50
S = cabacabaacbbcaaaabccbcabcccccbaabbbabccca
Part I: cabacabacbcababcbbcabcbacbabacba
Part II: 12 substrings: c abacaba acbbca aa baaaab c bcbcb abccccccba abbbba cbc c a
```
