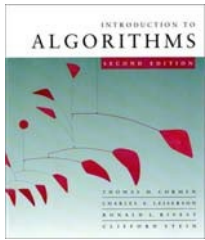


# CS60020: Foundations of Algorithm Design and Machine Learning

Sourangshu Bhattacharya



# The problem of sorting

**Input:** sequence  $\langle a_1, a_2, \dots, a_n \rangle$  of numbers.

**Output:** permutation  $\langle a'_1, a'_2, \dots, a'_n \rangle$  such that  $a'_1 \leq a'_2 \leq \dots \leq a'_n$ .

**Example:**

**Input:** 8 2 4 9 3 6

**Output:** 2 3 4 6 8 9

# Bubble Sort

# Sorting

- **Sorting takes an unordered collection and makes it an ordered one.**

|    |    |    |    |     |   |
|----|----|----|----|-----|---|
| 1  | 2  | 3  | 4  | 5   | 6 |
| 77 | 42 | 35 | 12 | 101 | 5 |



|   |    |    |    |    |     |
|---|----|----|----|----|-----|
| 1 | 2  | 3  | 4  | 5  | 6   |
| 5 | 12 | 35 | 42 | 77 | 101 |

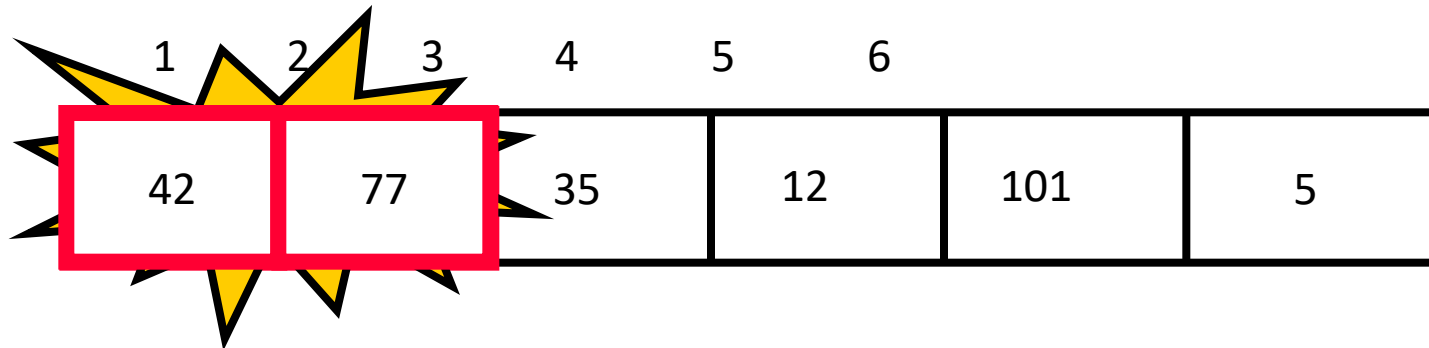
# "Bubbling Up" the Largest Element

- **Traverse a collection of elements**
  - Move from the front to the end
  - “Bubble” the **largest value** to the end using **pair-wise comparisons and swapping**

|    |    |    |    |     |   |
|----|----|----|----|-----|---|
| 1  | 2  | 3  | 4  | 5   | 6 |
| 77 | 42 | 35 | 12 | 101 | 5 |

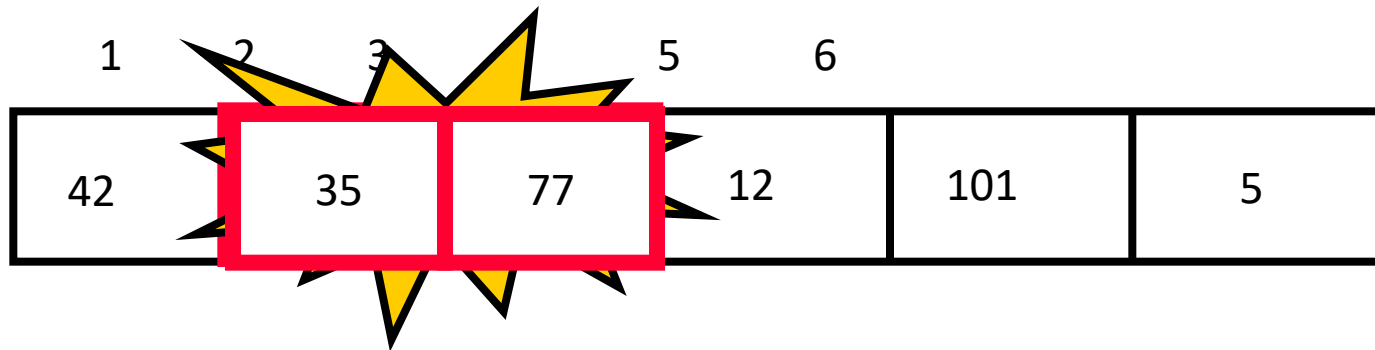
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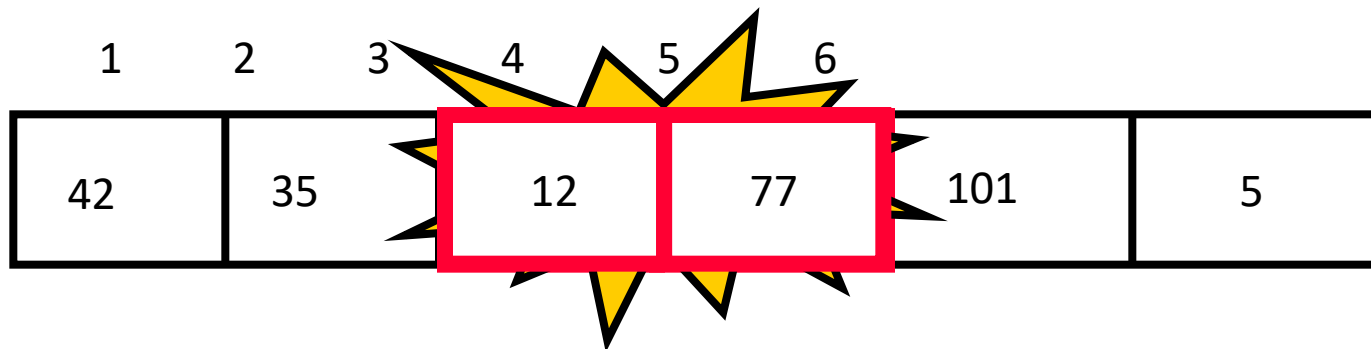
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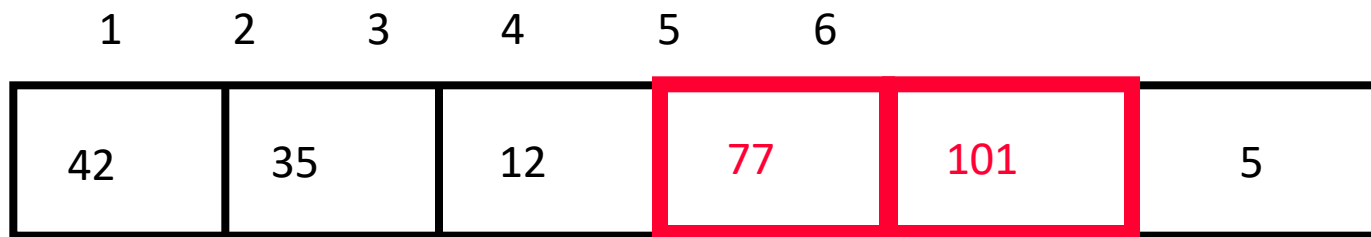
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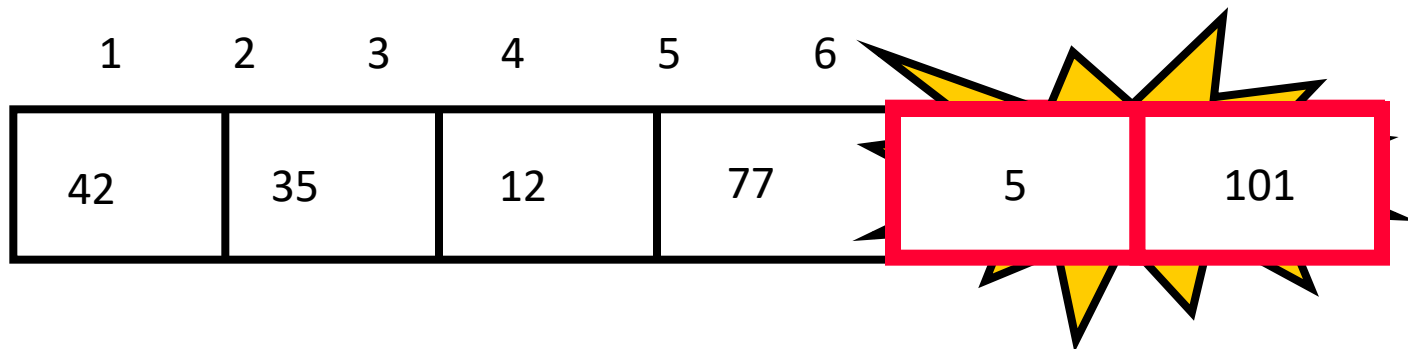
- **Traverse a collection of elements**
  - Move from the front to the end
  - “Bubble” the largest value to the end using pairwise comparisons and swapping



No need to swap

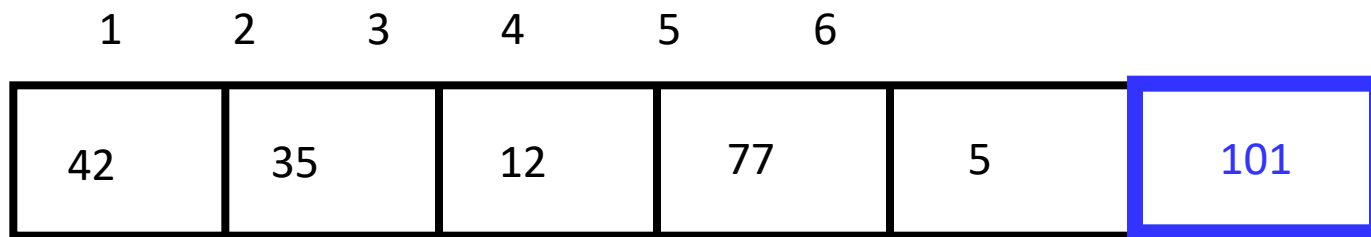
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# "Bubbling Up" the Largest Element

- **Traverse a collection of elements**
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Largest value correctly placed

# The “Bubble Up” Algorithm

```
index <- 1
```

```
last_compare_at <- n - 1
```

```
loop
```

```
  exitif(index > last_compare_at)
```

```
  if(A[index] > A[index + 1]) then
```

```
    Swap(A[index], A[index + 1])
```

```
  endif
```

```
  index <- index + 1
```

```
endloop
```

# Items of Interest

- Notice that only the largest value is correctly placed
- All other values are still out of order
- So we need to **repeat this process**

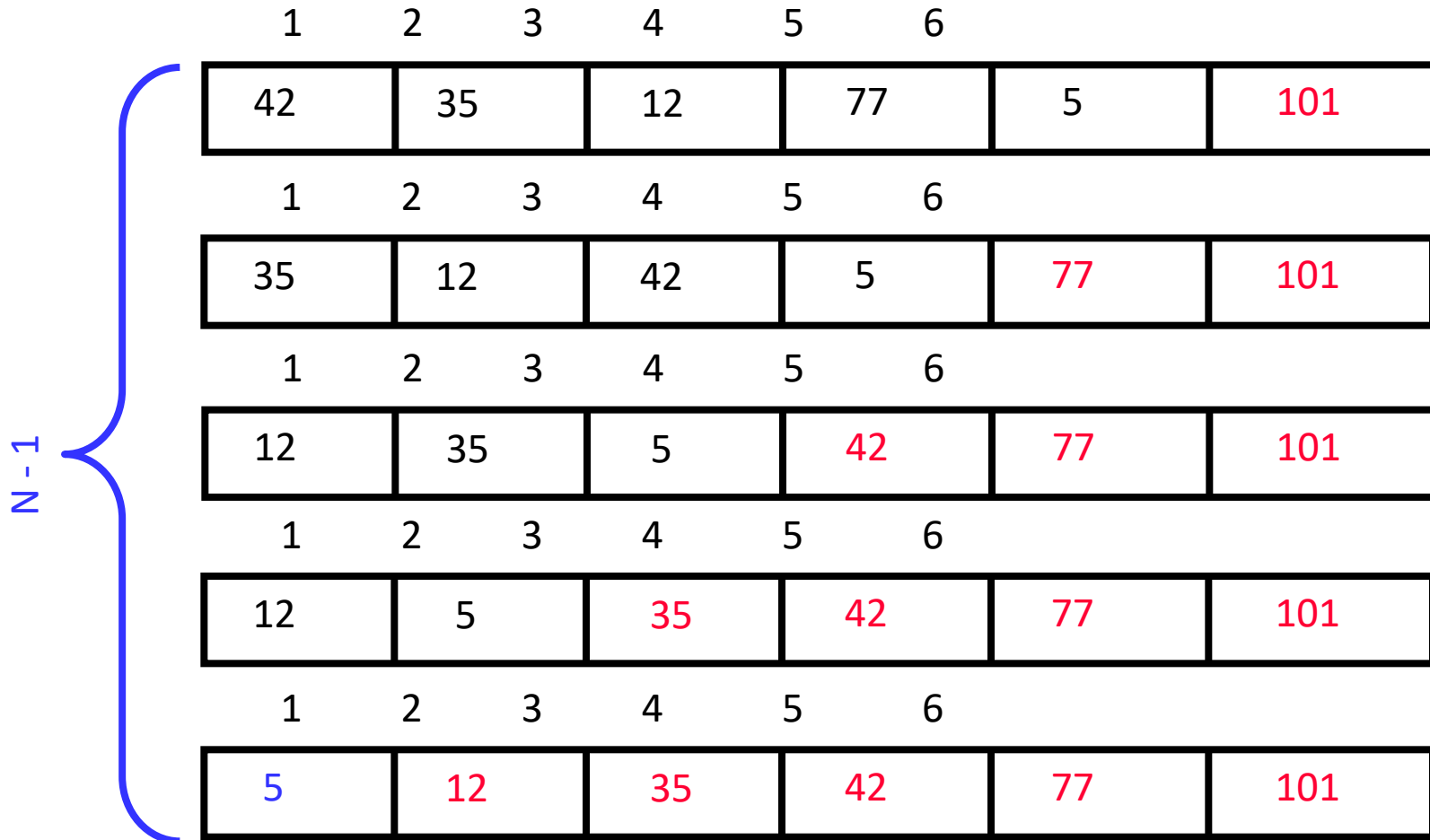
|    |    |    |    |   |     |
|----|----|----|----|---|-----|
| 1  | 2  | 3  | 4  | 5 | 6   |
| 42 | 35 | 12 | 77 | 5 | 101 |

Largest value correctly placed

# Repeat “Bubble Up” How Many Times?

- If we have  $N$  elements...
- And if each time we bubble an element, we place it in its correct location...
- Then we **repeat the “bubble up” process  $N - 1$  times.**
- **This guarantees we’ll correctly place all  $N$  elements.**

# “Bubbling” All the Elements



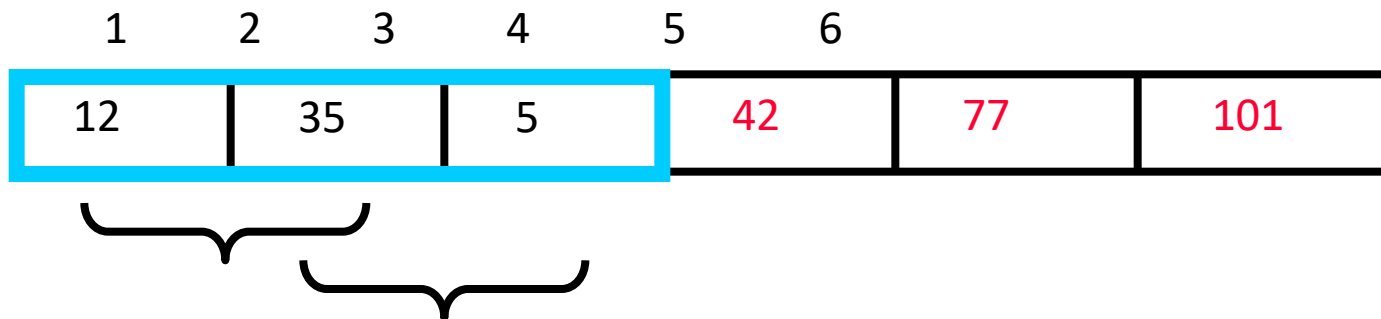
# Reducing the Number of Comparisons

| 1  | 2  | 3  | 4  | 5   | 6   |
|----|----|----|----|-----|-----|
| 77 | 42 | 35 | 12 | 101 | 5   |
| 1  | 2  | 3  | 4  | 5   | 6   |
| 42 | 35 | 12 | 77 | 5   | 101 |
| 1  | 2  | 3  | 4  | 5   | 6   |
| 35 | 12 | 42 | 5  | 77  | 101 |
| 1  | 2  | 3  | 4  | 5   | 6   |
| 12 | 35 | 5  | 42 | 77  | 101 |
| 1  | 2  | 3  | 4  | 5   | 6   |
| 12 | 5  | 35 | 42 | 77  | 101 |



# Reducing the Number of Comparisons

- On the  $N^{\text{th}}$  “bubble up”, we only need to do **MAX-N comparisons**.
- For example:
  - This is the 4<sup>th</sup> “bubble up”
  - MAX is 6
  - Thus we have **2 comparisons** to do



Putting It All Together

# Bubble Sort

```
procedure Bubblesort(A isoftype in/out Arr_Type)
  to_do, index isoftype Num
  to_do <- N - 1

  loop ←
    exitif(to_do = 0)
    index <- 1
    loop ←
      exitif(index > to_do)
      if(A[index] > A[index + 1]) then
        Swap(A[index], A[index + 1])
      endif
      index <- index + 1
    endloop
    to_do <- to_do - 1
  endloop
endprocedure // Bubblesort
```

Inner loop

Outer loop

# Already Sorted Collections?

- What if the collection was already sorted?
- What if only a few elements were out of place and after a couple of “bubble ups,” the collection was sorted?
- We want to be able to **detect this** and “stop early”!

|   |    |    |    |    |     |
|---|----|----|----|----|-----|
| 1 | 2  | 3  | 4  | 5  | 6   |
| 5 | 12 | 35 | 42 | 77 | 101 |

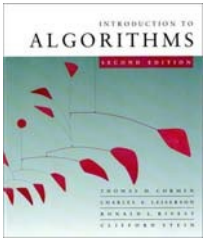
# Using a Boolean “Flag”

- We can use a boolean variable to determine if any swapping occurred during the “bubble up.”
- If no swapping occurred, then we know that the collection is already sorted!
- This boolean “flag” needs to be reset after each “bubble up.”

```
did_swap isoftype Boolean
did_swap <- true
```

```
loop
  exitif ((to_do = 0) OR NOT(did_swap))
  index <- 1
  did_swap <- false
  loop
    exitif(index > to_do)
    if(A[index] > A[index + 1]) then
      Swap(A[index], A[index + 1])
      did_swap <- true
    endif
    index <- index + 1
  endloop
  to_do <- to_do - 1
endloop
```

# **INSERTION SORT**

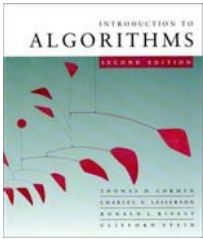


# Insertion sort

“pseudocode”

```
INSERTION-SORT ( $A, n$ )     $\triangleleft A[1 \dots n]$   
  for  $j \leftarrow 2$  to  $n$   
    do  $key \leftarrow A[j]$   
       $i \leftarrow j - 1$   
      while  $i > 0$  and  $A[i] > key$   
        do  $A[i+1] \leftarrow A[i]$   
           $i \leftarrow i - 1$   
       $A[i+1] = key$ 
```

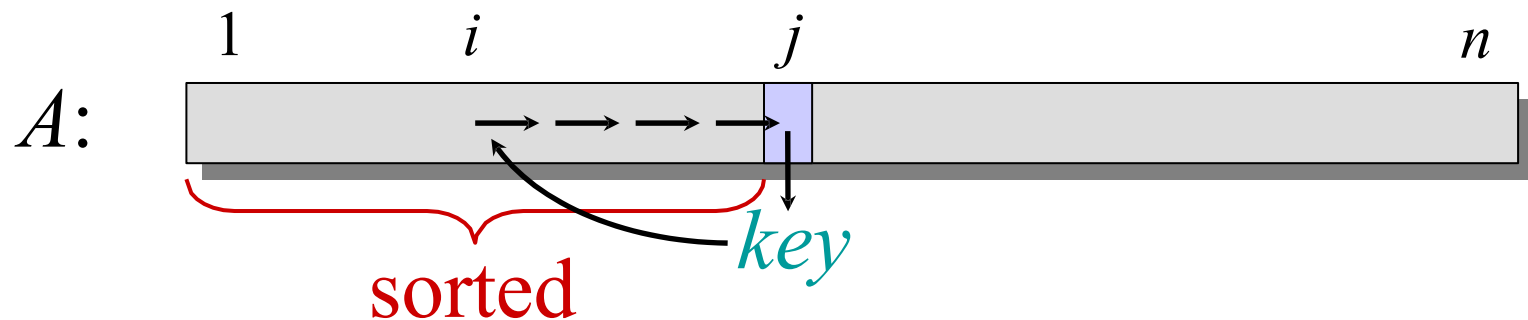


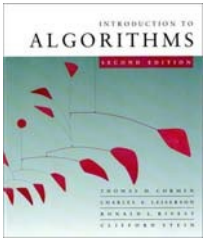


# Insertion sort

“pseudocode”

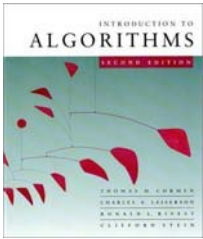
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  for  $j \leftarrow 2$  to  $n$   
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        $i \leftarrow j - 1$   
       while  $i > 0$  and  $A[i] > key$   
         do  $A[i+1] \leftarrow A[i]$   
             $i \leftarrow i - 1$   
        $A[i+1] = key$ 
```





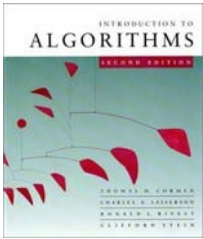
# Example of insertion sort

8 2 4 9 3 6



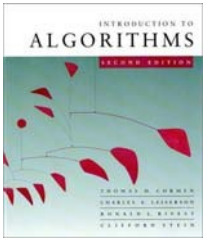
# Example of insertion sort



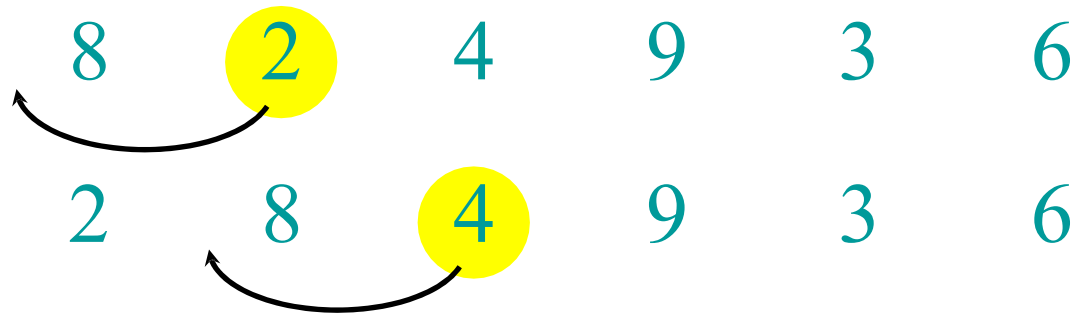


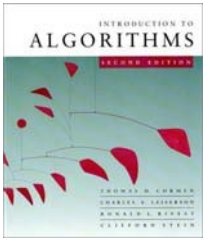
# Example of insertion sort



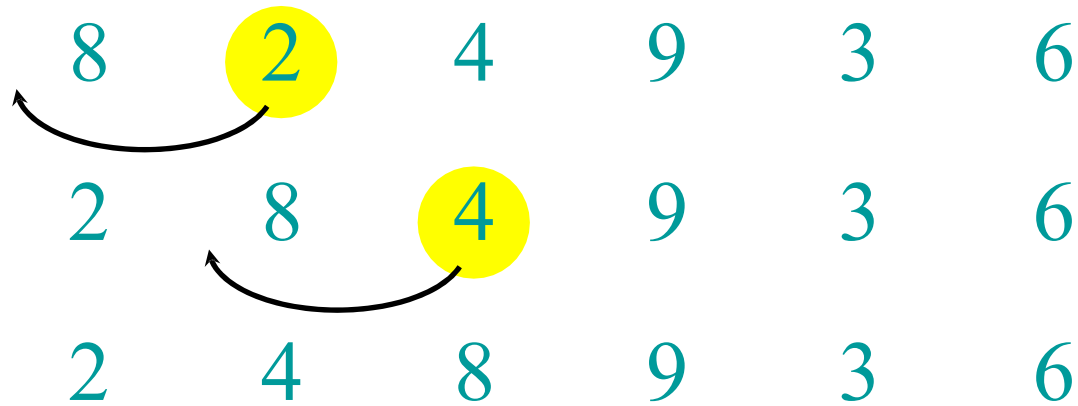


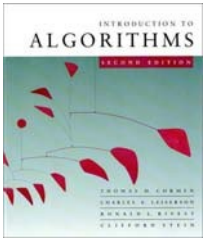
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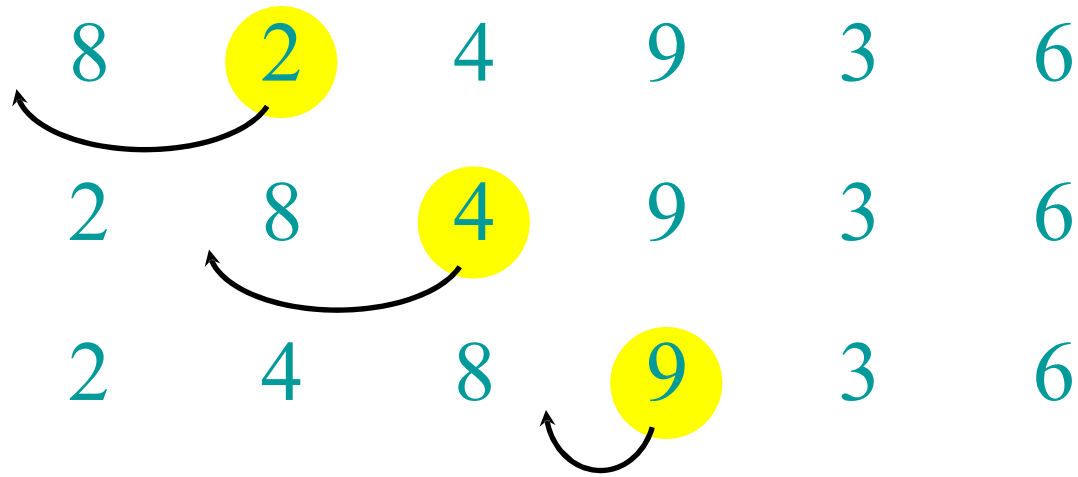


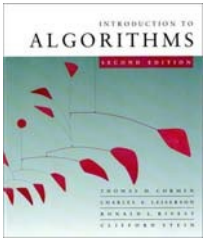
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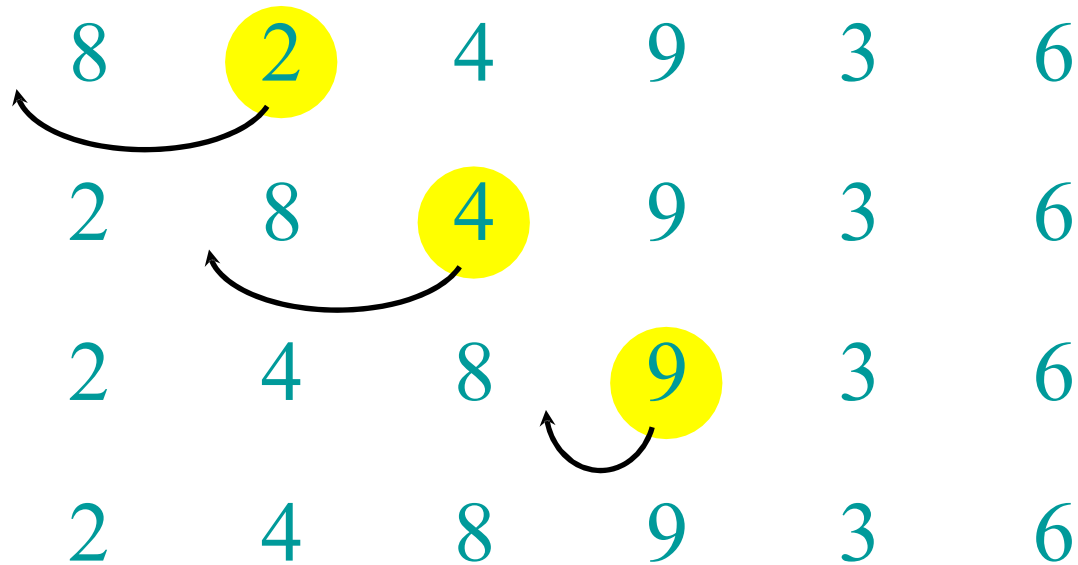


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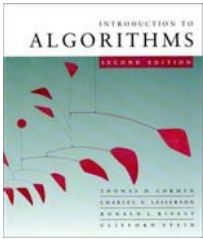




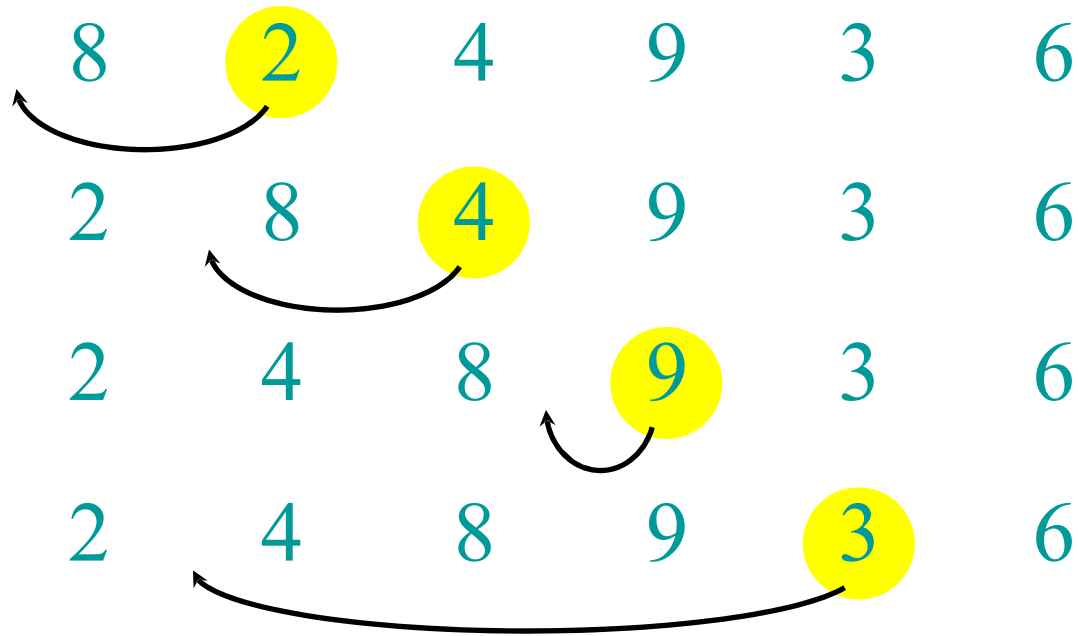
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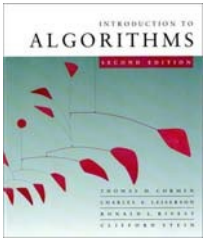




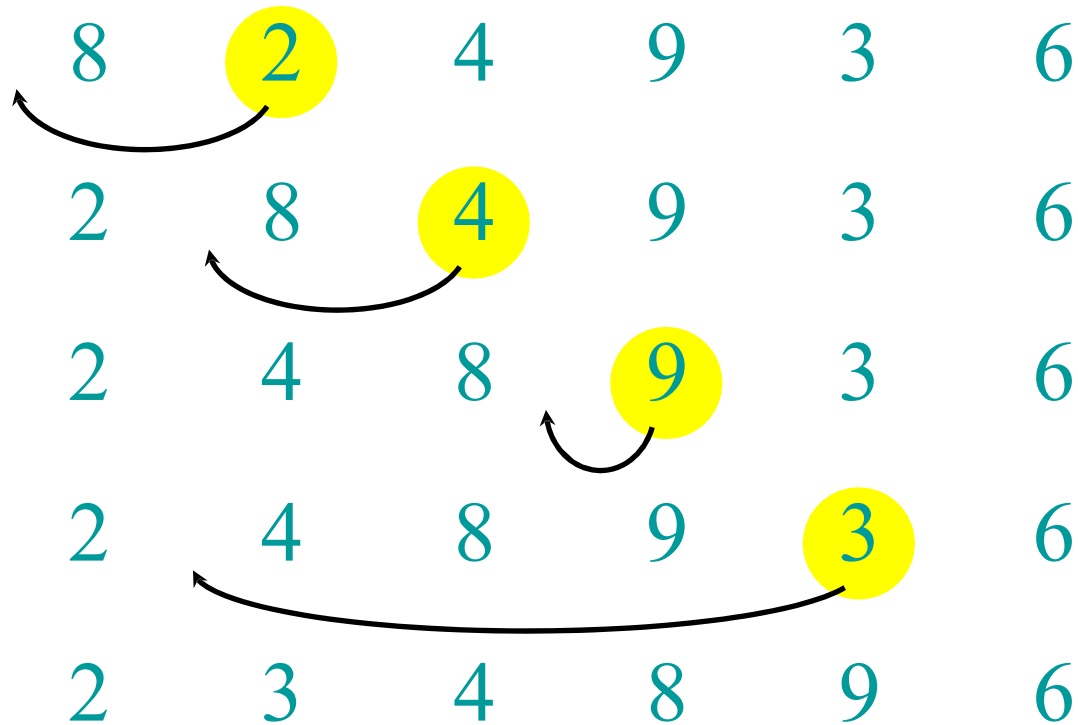


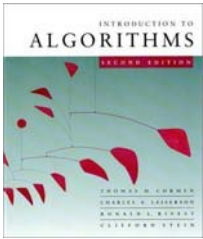
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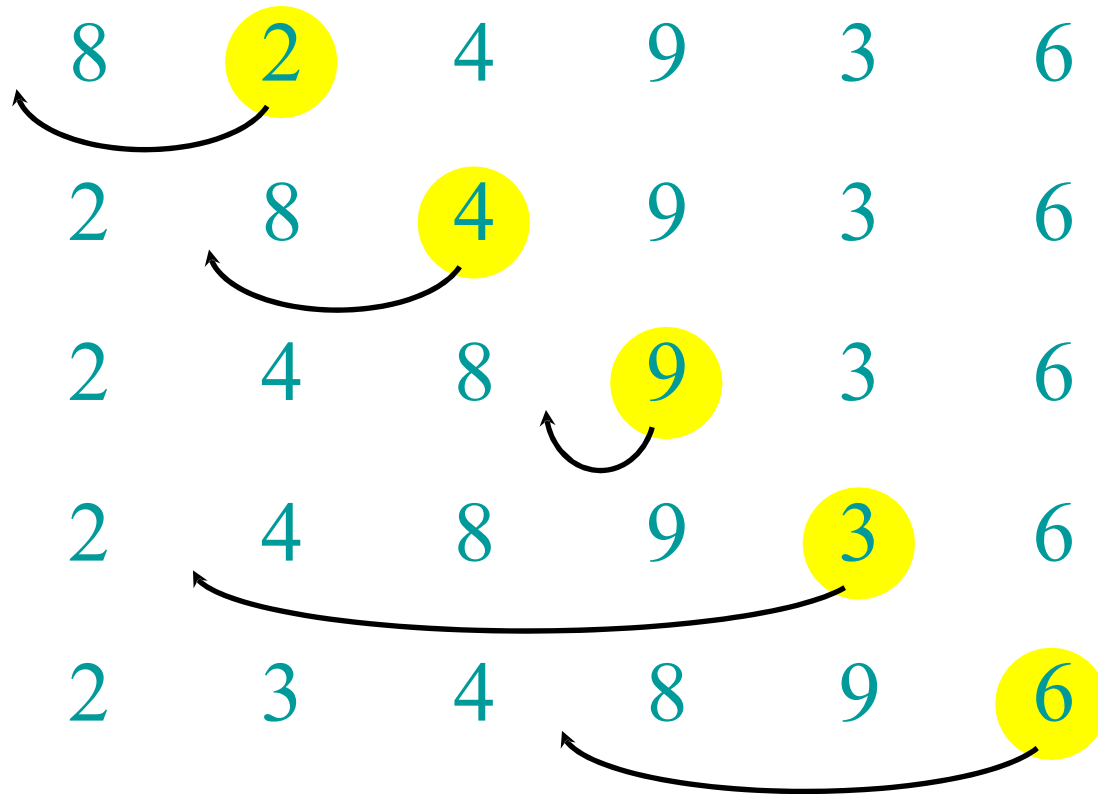


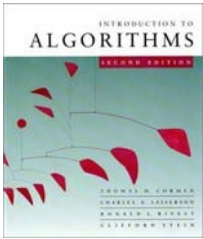
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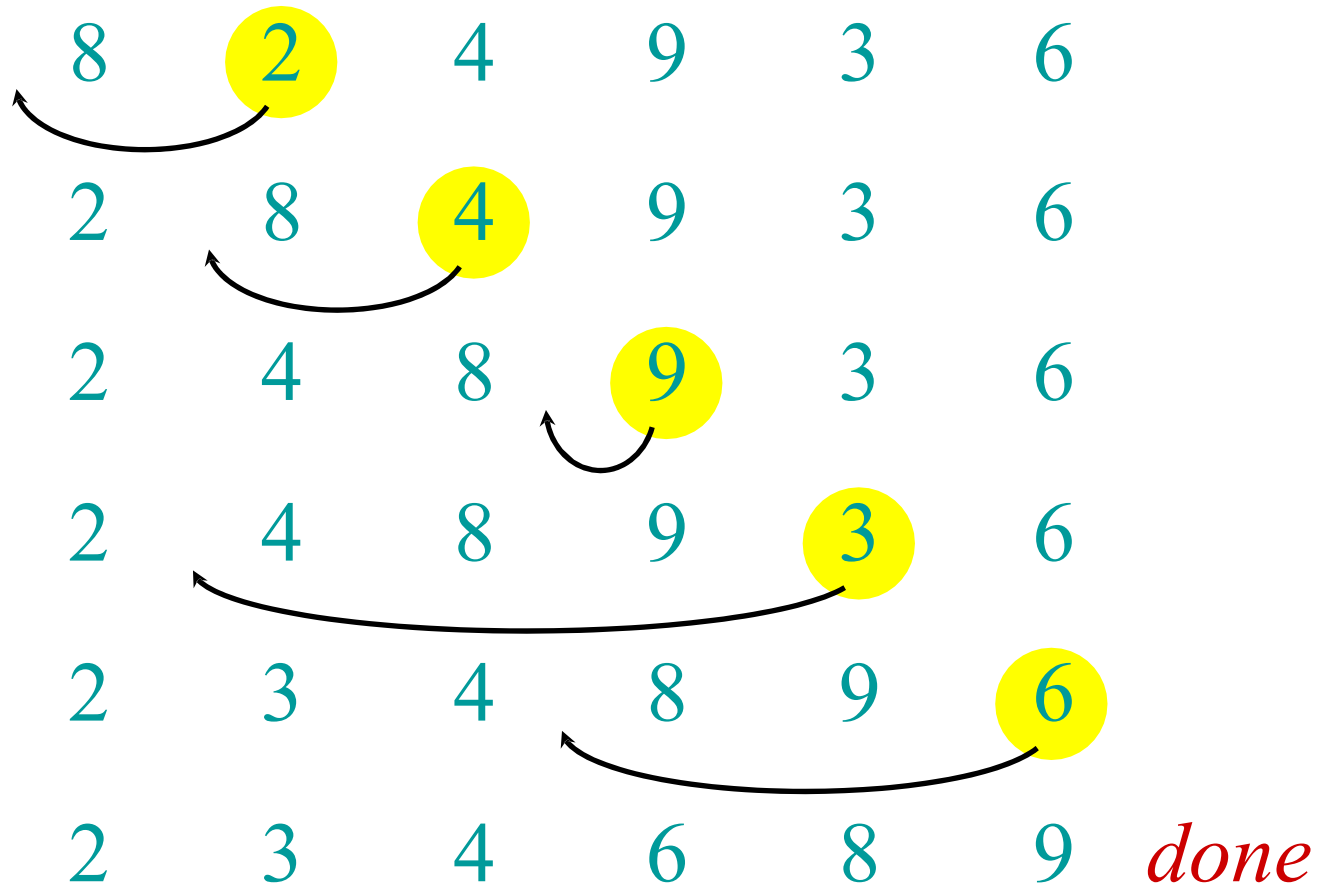


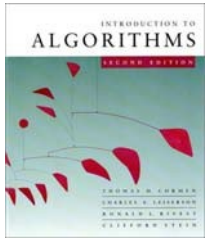
# Example of insertion sort





# Example of insertion sort





# Insertion sort analysis

***Worst case:*** Input reverse sorted.

$$T(n) = \sum_{j=2}^n \Theta(j) = \Theta(n^2) \quad [\text{arithmetic series}]$$

***Average case:*** All permutations equally likely.

$$T(n) = \sum_{j=2}^n \Theta(j/2) = \Theta(n^2)$$

*Is insertion sort a fast sorting algorithm?*

- Moderately so, for small  $n$ .
- Not at all, for large  $n$ .

# Analysis

|  | <i>cost</i> | <i>times</i>             |
|--|-------------|--------------------------|
| INSERTION-SORT( <i>A</i> )   |             |                          |
| 1 <b>for</b> $j = 2$ <b>to</b> $A.length$                            | $c_1$       | $n$                      |
| 2 $key = A[j]$   | $c_2$       | $n - 1$                  |
| 3     // Insert $A[j]$ into the sorted<br>sequence $A[1 .. j - 1]$ . | 0           | $n - 1$                  |
| 4 $i = j - 1$  | $c_4$       | $n - 1$                  |
| 5 <b>while</b> $i > 0$ and $A[i] > key$                              | $c_5$       | $\sum_{j=2}^n t_j$       |
| 6 $A[i + 1] = A[i]$  | $c_6$       | $\sum_{j=2}^n (t_j - 1)$ |
| 7 $i = i - 1$  | $c_7$       | $\sum_{j=2}^n (t_j - 1)$ |
| 8 $A[i + 1] = key$   | $c_8$       | $n - 1$                  |

# Comparison

| Sorting Algorithm | Time Complexity |              |            | Space Complexity |
|-------------------|-----------------|--------------|------------|------------------|
|                   | Best Case       | Average Case | Worst Case | Worst Case       |
| Bubble Sort       | $O(N)$          | $O(N^2)$     | $O(N^2)$   | $O(1)$           |
| Selection Sort    | $O(N^2)$        | $O(N^2)$     | $O(N^2)$   | $O(1)$           |
| Insertion Sort    | $O(N)$          | $O(N^2)$     | $O(N^2)$   | $O(1)$           |