

CS11001/CS11002

Programming and Data Structures

(PDS) (Theory: 3-0-0)

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Sorting: the basic problem

- Given an array

$x[0], x[1], \dots, x[\text{size}-1]$

reorder entries so that

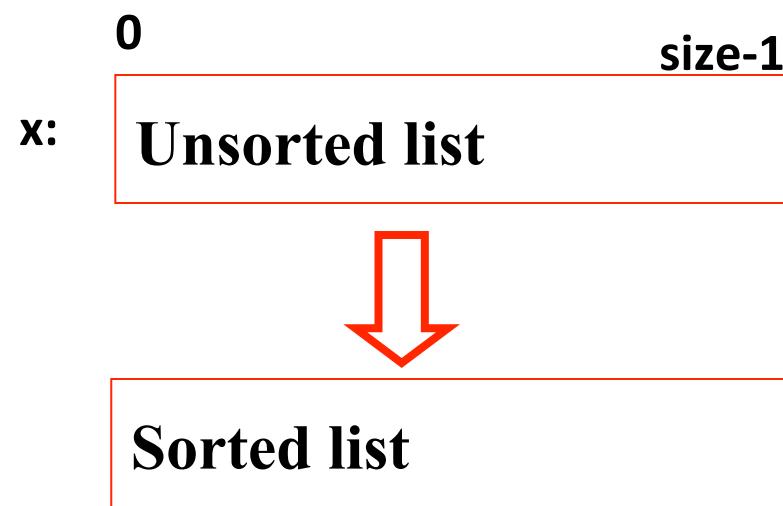
$x[0] \leq x[1] \leq \dots \leq x[\text{size}-1]$

List is in non-decreasing order.

- We can also sort a list of elements in non-increasing order.

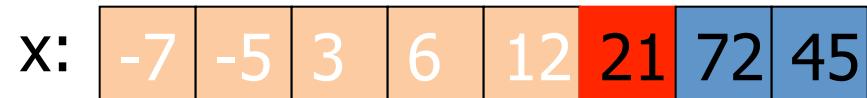
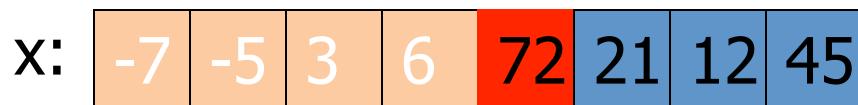
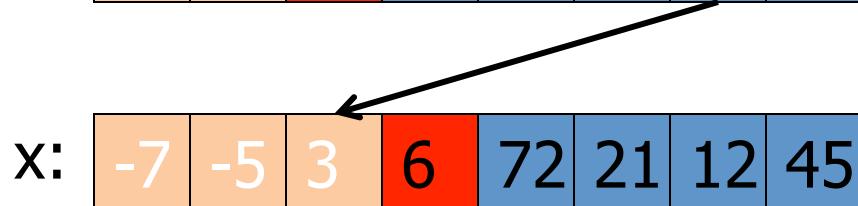
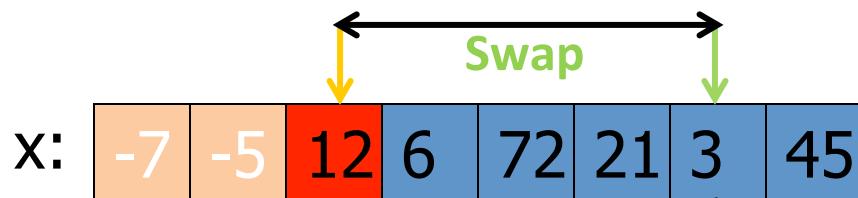
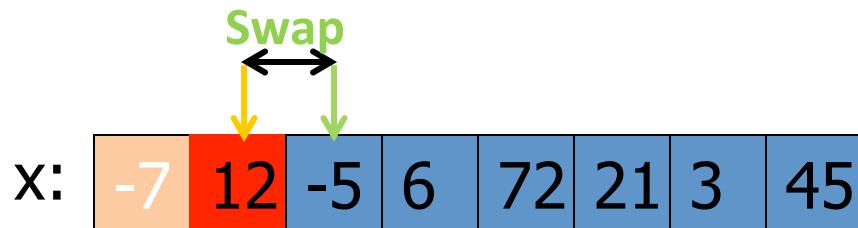
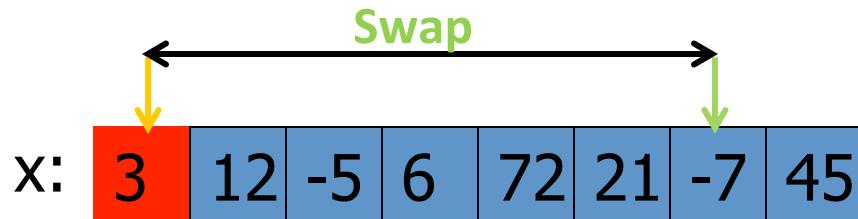
Sorting Problem

- What we want : Data sorted in order
- Input: A list of elements
- Output: A list of elements in sorted (non-increasing/non-decreasing) order



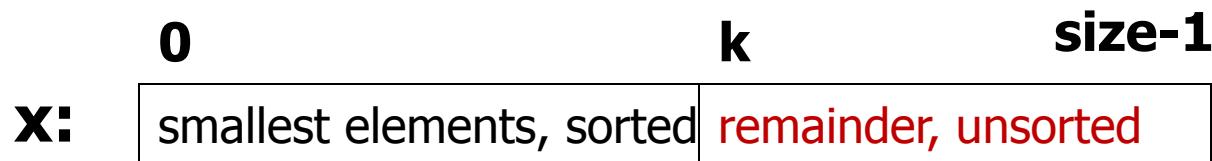
- Original list:
 - 10, 30, 20, 80, 70, 10, 60, 40, 70
- Sorted in non-decreasing order:
 - 10, 10, 20, 30, 40, 60, 70, 70, 80
- Sorted in non-increasing order:
 - 80, 70, 70, 60, 40, 30, 20, 10, 10

Example



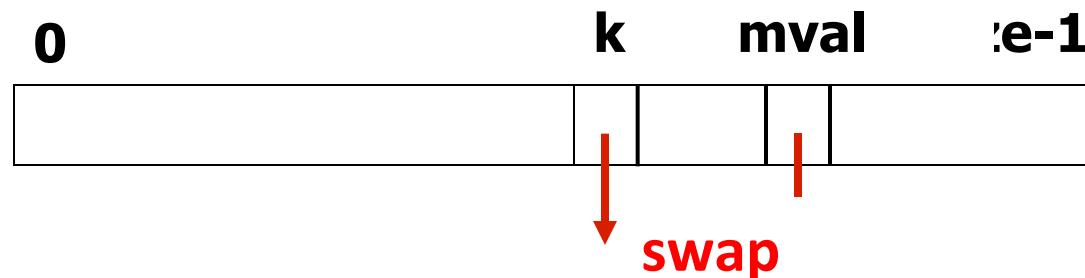
Selection Sort

- General situation :



- Steps :

- Find smallest element, `mval`, in `x[k+1..size-1]`
- Swap smallest element with `x[k]`,
- Increase `k`.



Subproblem: Find smallest element

```
/* Yield location of smallest element in x[k ..  
size-1] and store in pos*/  
  
int j, pos;  
pos = k;      /* assume first element is the smallest element */  
for (j=k+1; j<size; j++) {  
    if (x[j] < x[pos]) { /* x[pos] is the smallest element as  
of now */  
        pos = j;  
    }  
}  
printf("%d", pos);
```

Subproblem: Swap with smallest element

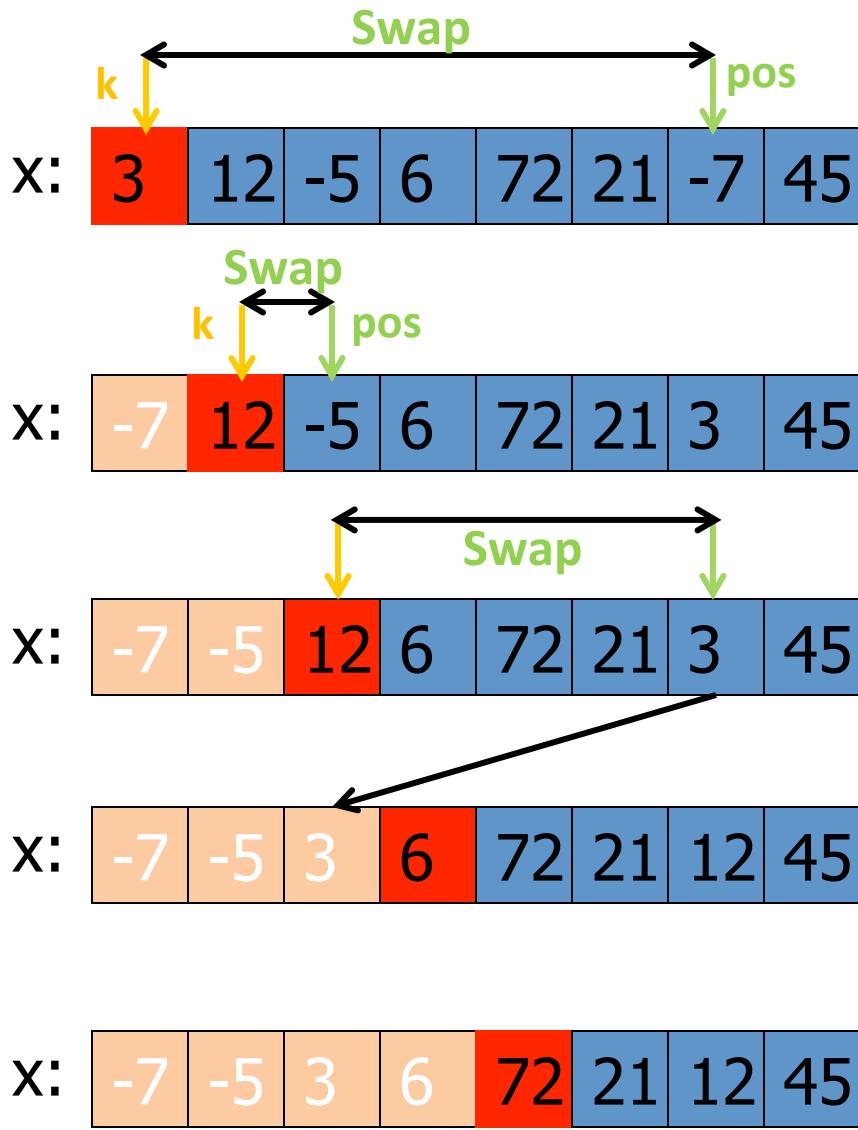
```
/* Yield location of smallest element in x[k .. size-1] and
store in pos*/

int j, pos;
pos = k; /* assume first element is the smallest
element */
for (j=k+1; j<size; j++) {
    if (x[j] < x[pos]) { /* x[pos] is the smallest element
as of now */
        pos = j;
    }
}
printf("%d", pos);
if (x[pos] < x[k]) {
    temp = x[k]; /* swap content of x[k] and x[pos] */
    x[k] = x[pos];
    x[pos] = temp;
}
```

```
#include <stdio.h> /* Sort x[0..size-1] in non-decreasing order */
int main()
{
    int k,j,pos,x[100],size,temp
    printf("Enter the number of elements: ");
    scanf("%d",&size);
    printf("Enter the elements: ");
    for(k=0;k<size;k++) scanf("%d",&x[k]);
    for(k=0; k<size-1; k++) {
        pos = k; /* assume first element is the smallest element */
        for(j=k+1; j<size; j++) {
            if (x[j] < x[pos]) /*x[pos] is the smallest element as of now*/
                pos = j;
        }
        temp = x[k];           /* swap content of x[k] and x[pos] */
        x[k] = x[pos];
        x[pos] = temp;
    }

    for(k=0;k<size;k++) /* print the sorted (non-decreasing) list */
        printf("%d ",x[k]);
    return 0;
}
```

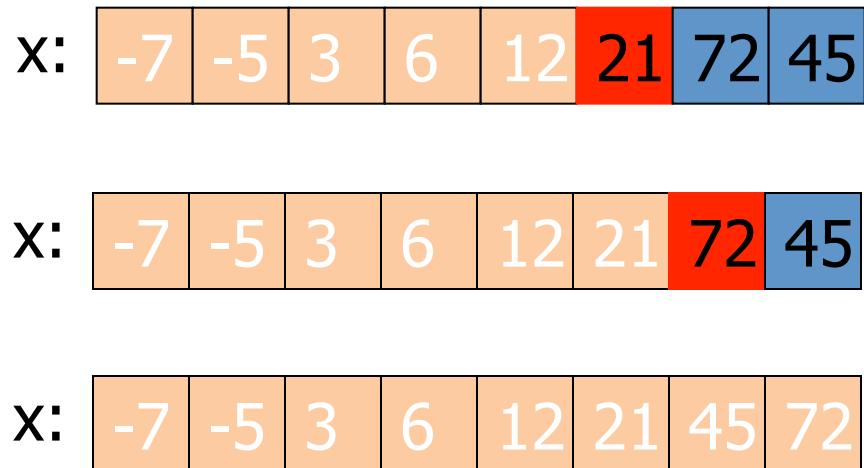
Example



```

for (k=0; k<size-1; k++) {
    pos = k;
    for (j=k+1; j<size; j++) {
        if (x[j] < x[pos])
            pos = j;
    }
    temp = x[k];
    = x[pos];
    x[pos] = temp;
}

```



Analysis

How many steps are required?

Let us assume there are n elements (size=n).
Each statement executes in constant time.

To read

for loop will take of the order of n time

Sorting

When k=0, in worst case

(n-1) comparisons to find minimum.

When k=1, in worst case

(n-2) comparisons to find minimum.

.....

(n-1)+(n-2)+.....+1= n(n-1)/2

To print

for loop will take of the order of n time

```
for (k=0; k<size; k++)  
  
    scanf ("%d", &x[k]);  
for (k=0; k<size-1; k++) {  
    pos = k;  
    for (j=k+1; j<size; j++) {  
        if (x[j] < x[pos])  
            pos = j;  
    }  
    temp = x[k];  
    x[k] = x[pos];  
    x[pos] = temp;  
}  
for (k=0; k<size; k++)  
    printf ("%d ", x[k]);
```

Total time= $2 \times \text{order of } n + \text{order of } n^2 = \text{order of } n^2$

Analysis

How many steps are required?

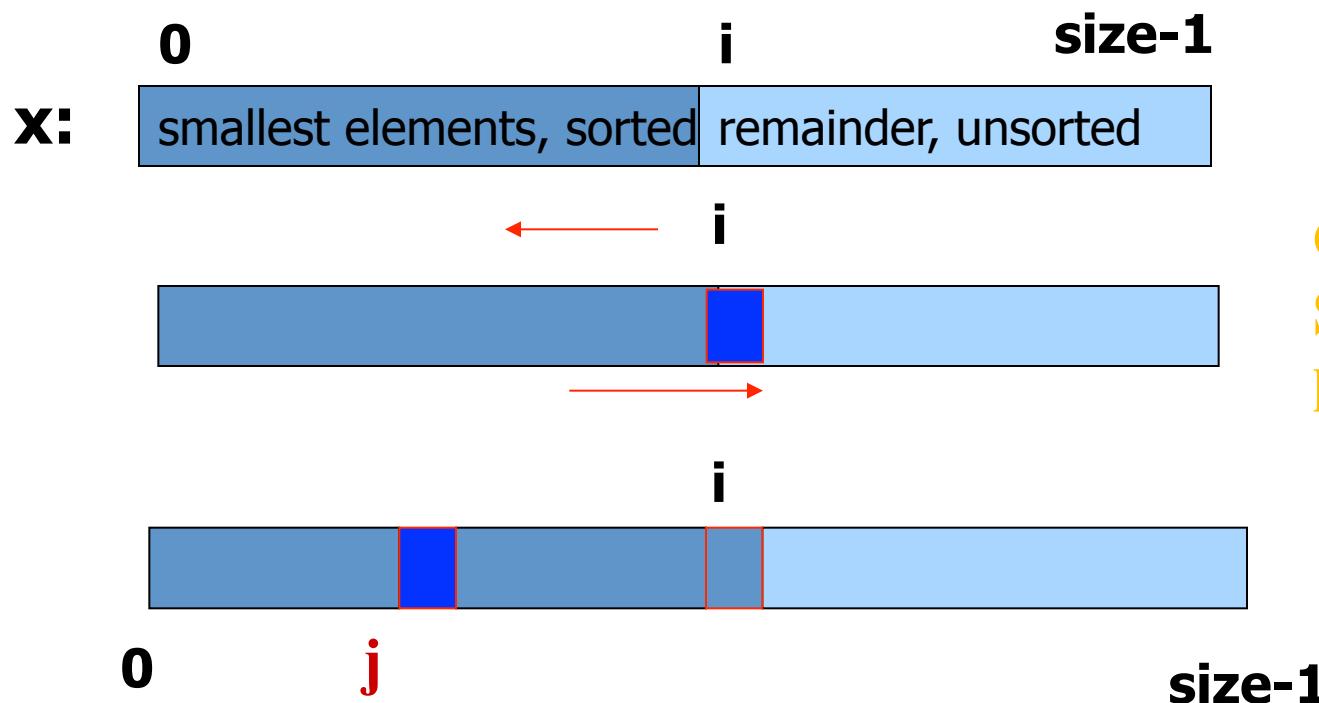
Let us assume there are n elements ($\text{size}=n$).
Each statement executes in constant time.

Best Case?
Worst Case?
Average Case?

```
for (k=0; k<size; k++)  
  
    scanf ("%d", &x[k]);  
for (k=0; k<size-1; k++) {  
    pos = k;  
    for (j=k+1; j<size; j++) {  
        if (x[j] < x[pos])  
            pos = j;  
    }  
    temp = x[k];  
    x[k] = x[pos];  
    x[pos] = temp;  
}  
for (k=0; k<size; k++)  
    printf ("%d ", x[k]);
```

Insertion Sort

- General situation :



Insertion Sorting

```
int list[100], size;  
_____  
;  
  
for (i=1; i<size; i++) {  
    item = list[i] ;  
    for (j=i-1; (j>=0) && (list[j] > item); j--)  
        list[j+1] = list[j];  
    list[j+1] = item ;  
}  
_____;
```

Complete Insertion Sort

```
#include <stdio.h> /* Sort x[0..size-1] in non-decreasing order */
#define SIZE 100
int main()
{
    int i,j,x[SIZE],size,temp;
    printf("Enter the number of elements: ");
    scanf("%d",&size);
    printf("Enter the elements: ");
    for(i=0;i<size;i++)
        scanf("%d",&x[i]);
    for (i=1; i<size; i++) {
        temp = x[i] ;
        for (j=i-1; (j>=0) && (x[j] > temp); j--)
            x[j+1] = x[j];
        x[j+1] = temp ;
    }
    for(i=0;i<size;i++) /* print the sorted (non-decreasing) list */
        printf("%d ",x[i]);
    return 0;
}
```

Insertion Sort Example

```
for (i=1; i<size; i++) {  
    temp = x[i] ;  
    for (j=i-1; (j>=0)&& (x[j] > temp); j--)  
        x[j+1] = x[j];  
    x[j+1] = temp ;  
}
```

Input:

3 12 -5 6 72 21 -7 45

3 12 -5 6 72 21 -7 45

-5 3 12 6 72 21 -7 45

-5 3 6 12 72 21 -7 45

-5 3 6 12 72 21 -7 45

-5 3 6 12 21 72 -7 45

-7 -5 3 6 12 21 72 45

-7 -5 3 6 12 21 45 72

Output:

-5 3 6 12 21 72 -7 45
-7

-5 3 6 12 21 72 _____ 45

-5 3 6 12 21 _____ 72 45

-5 3 6 12 _____ 21 72 45

-5 3 _____ 12 21 72 45

-5 _____ 6 12 21 72 45

_____ -5 3 6 12 21 72 45

_____ -7 -5 3 6 12 21 72 45



Time Complexity

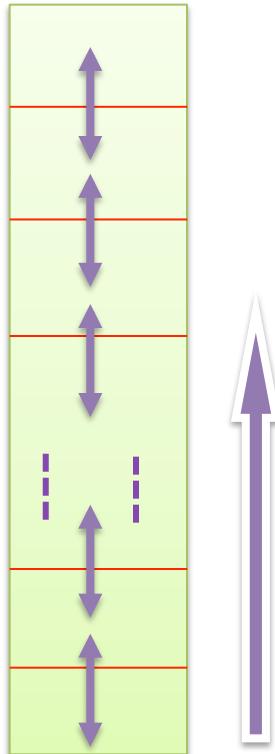
- Number of comparisons and shifting:
 - Worst Case?

$$1+2+3+ \dots + (n-1) = n(n-1)/2$$

- Best Case?

$$1+1+\dots \text{ (n-1) times} = (n-1)$$

Bubble Sort

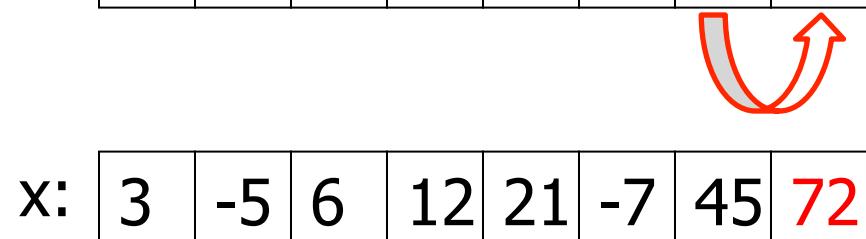
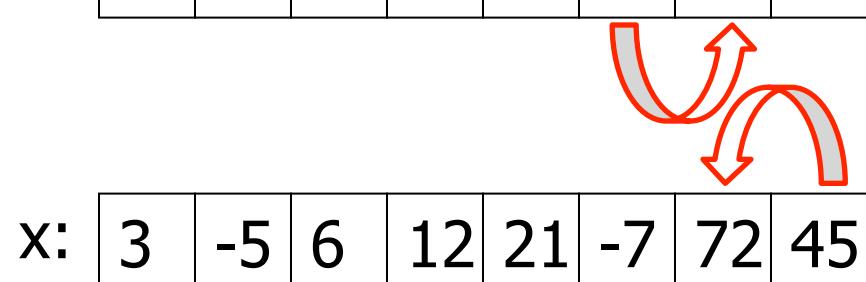
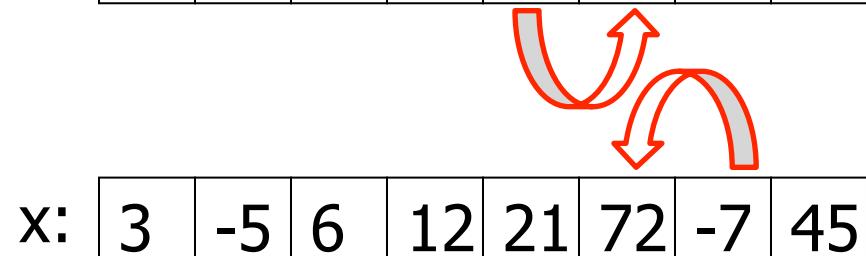
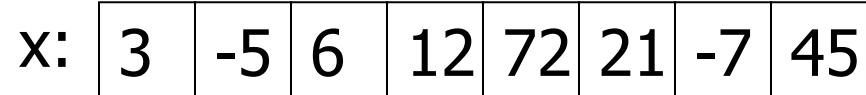
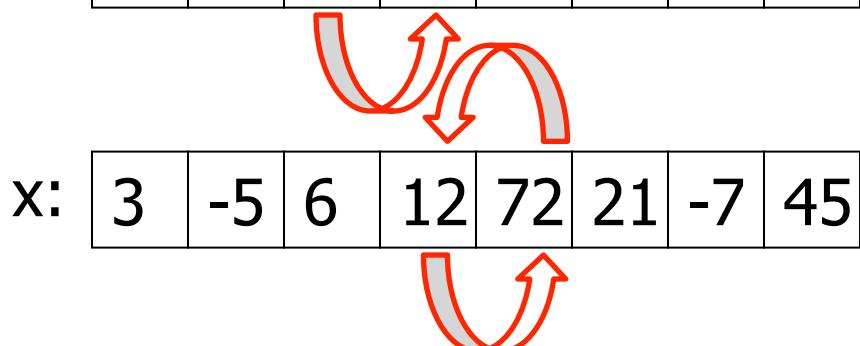
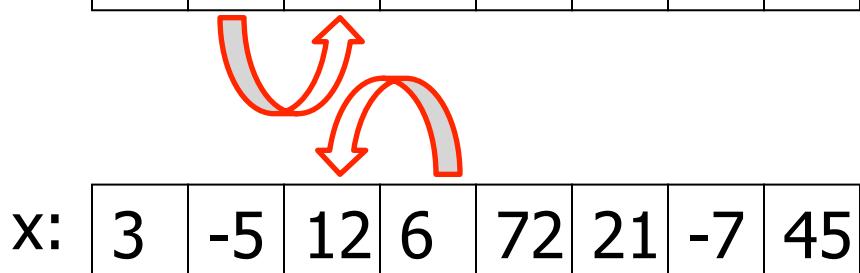
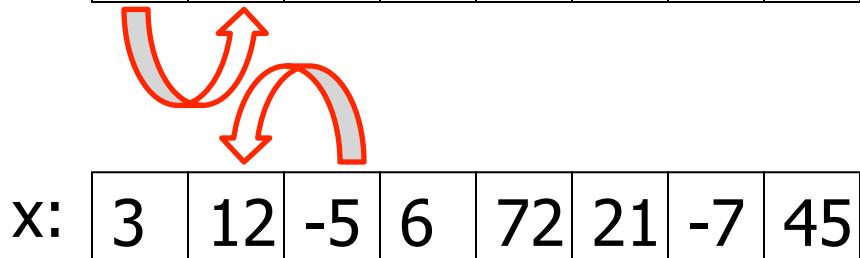
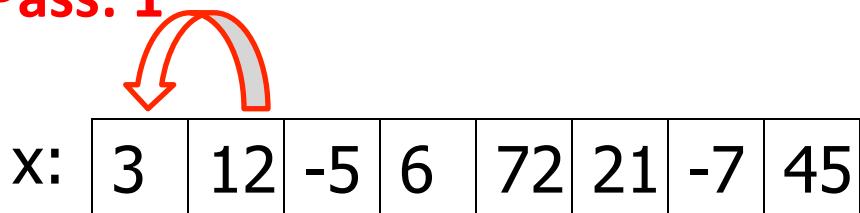


In every iteration heaviest element drops at the bottom.

The bottom moves upward.

Example

Pass: 1

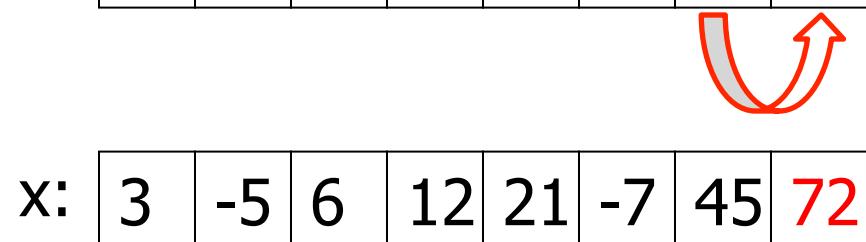
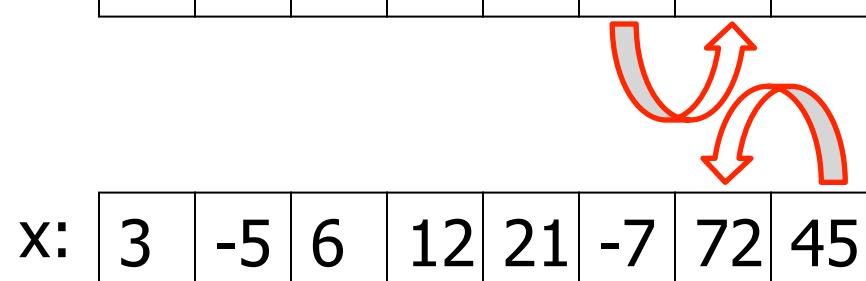
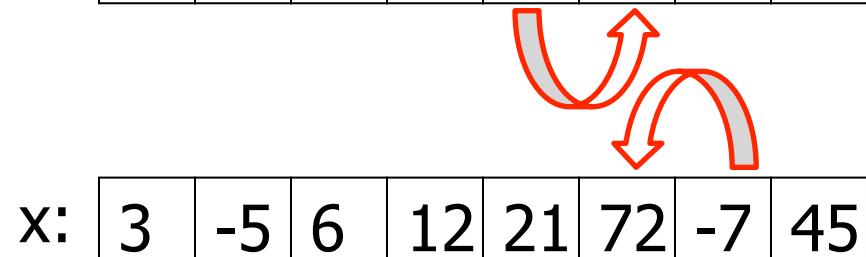
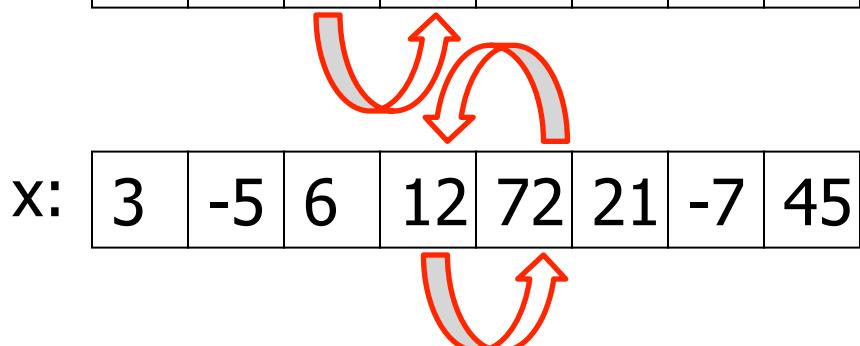
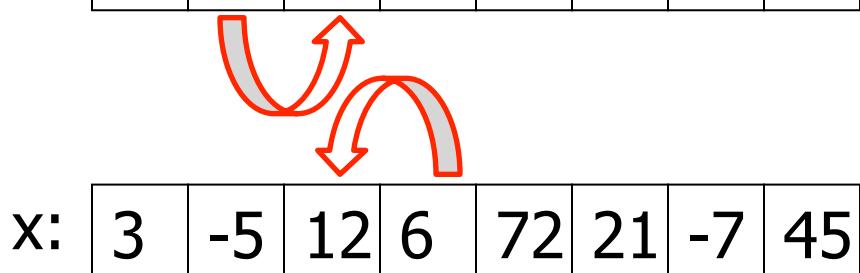
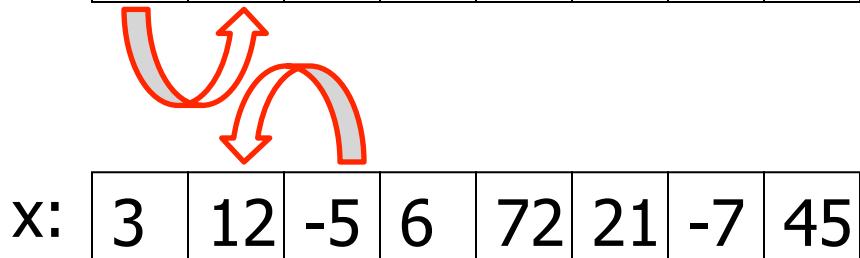
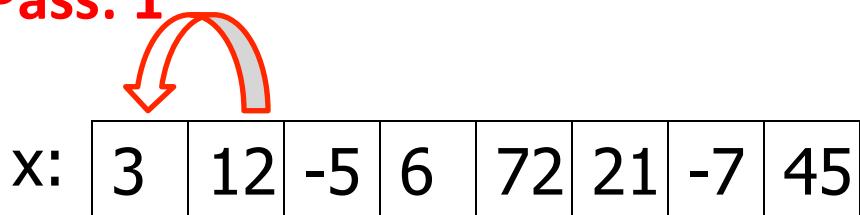


Bubble Sort

```
int pass,j,a[100],temp;  
  
/* read number of elements as n and elements as a[] */  
  
for(pass=0; pass<n; pass++) {  
    /* sub pass */  
    for(j=0; j<n-1; j++) {  
        if(a[j]>a[j+1]) {  
            temp = a[j+1];  
            a[j+1] = a[j];  
            a[j] = temp;  
        }  
    }  
}  
  
/* print sorted list of elements */
```

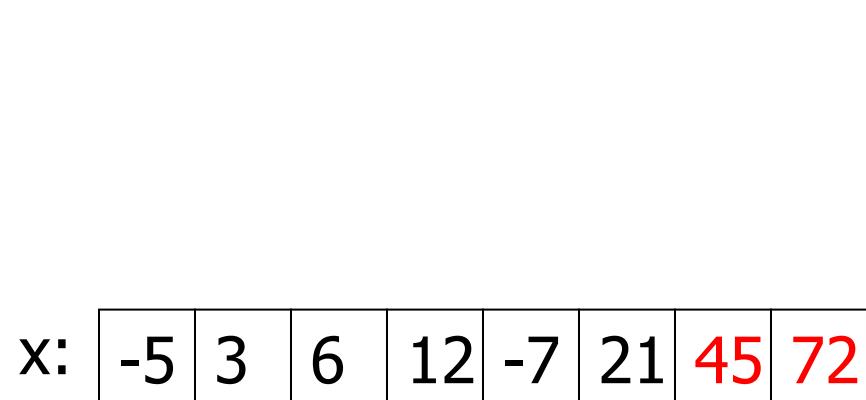
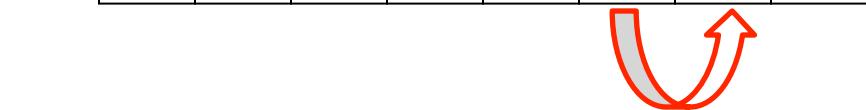
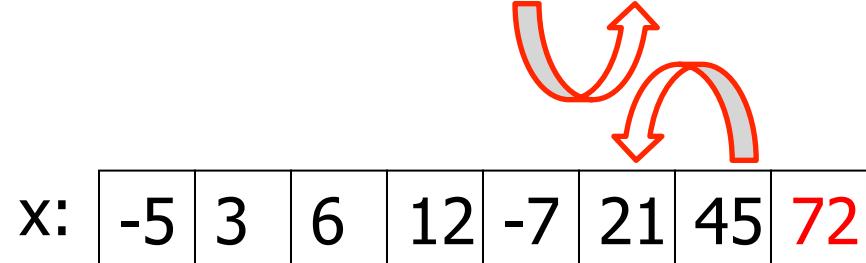
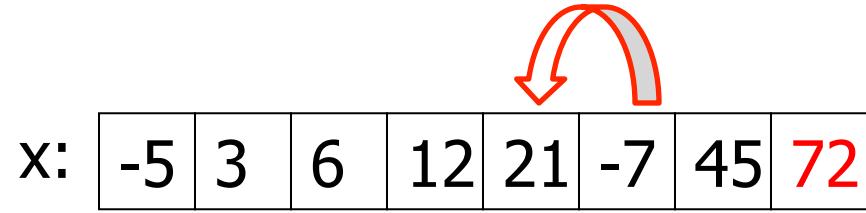
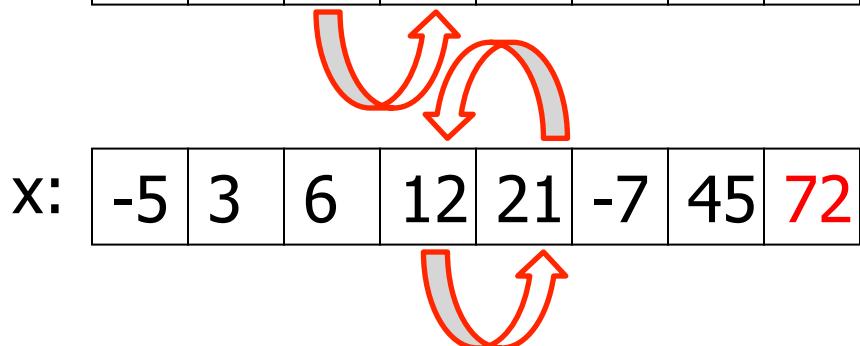
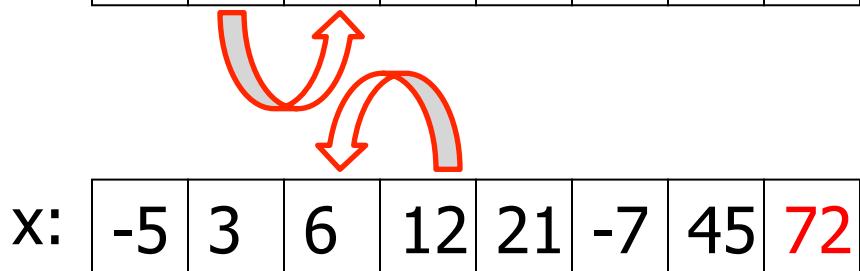
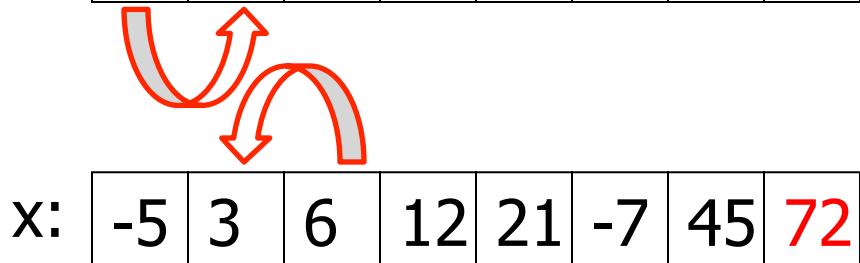
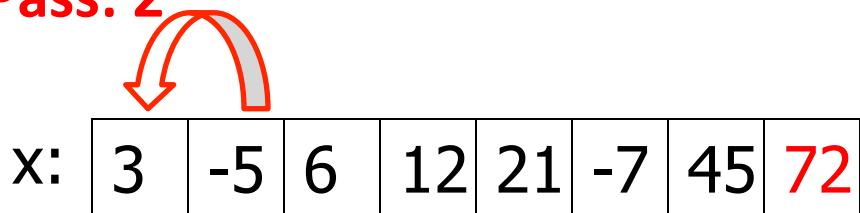
Example

Pass: 1



Example

Pass: 2



Time Complexity

- Number of comparisons :

- Worst Case?

$$1+2+3+ \dots + (n-1) = n(n-1)/2$$

- Best Case?

Same

How do you make best case with
(n-1) comparisons only?

Bubble Sort

```
int pass,j,a[100],temp,swapflag;
/* read number of elements as n and elements as a[] */

for(pass=0; pass<n; pass++)  {

    swapflag=0;
    for(j=0; j<n-1; j++)  {
        if(a[j]>a[j+1])  {
            temp = a[j+1];
            a[j+1] = a[j];
            a[j] = temp;
            swapflag=1;
        }
    }
    if(swapflag==0)
        break;
}
/* print sorted list of elements */
```

Can we improve the sorting time?

Basis of efficient sorting algorithms

- Two of the most popular sorting algorithms are based on divide-and-conquer approach.
 - Quick sort
 - Merge sort
- Basic concept:

```
sort (list)
{
    if the list has length greater than 1
    {
        Partition the list into lowlist and highlist;
        sort (lowlist);
        sort (highlist);
        combine (lowlist, highlist);
    }
}
```

Merge Sort

Example

x: [3 | 12 | -5 | 6 | 72 | 21 | -7 | 45]

[3 | 12 | -5 | 6]

Splitting arrays

[72 | 21 | -7 | 45]

[3 | 12]

[-5 | 6]

[72 | 21]

[-7 | 45]

[3 | 12]

[-5 | 6]

[72 | 21]

[-7 | 45]

[3 | 12]

[-5 | 6]

[21 | 72]

[-7 | 45]

[-5 | 3 | 6 | 12]

Merging two sorted arrays

[-7 | 21 | 45 | 72]



[-7 | -5 | 3 | 6 | 12 | 21 | 45 | 72]

Merge Sort C program

```
#include<stdio.h>
void mergesort(int a[],int i,int j);
void merge(int a[],int i1,int j1,int i2,int j2);

int main()
{
    int a[30],n,i;
    printf("Enter no of elements:");
    scanf("%d",&n);
    printf("Enter array elements:");

    for(i=0;i<n;i++)
        scanf("%d",&a[i]);

    mergesort(a,0,n-1);

    printf("\nSorted array is :");
    for(i=0;i<n;i++)
        printf("%d ",a[i]);

    return 0;
}
```

Merge Sort C program

```
void mergesort(int a[],int i,int j)
{
    int mid;

    if(i<j)      {
        mid=(i+j)/2;
        /* left recursion */
        mergesort(a,i,mid);
        /* right recursion */
        mergesort(a,mid+1,j);
        /* merging of two sorted sub-arrays */
        merge(a,i,mid,mid+1,j);
    }
}
```

Merge Sort C program

```
void merge(int a[],int i1,int i2,int j1,int j2)
{
    int temp[50];      //array used for merging
    int i=i1,j=j1,k=0;

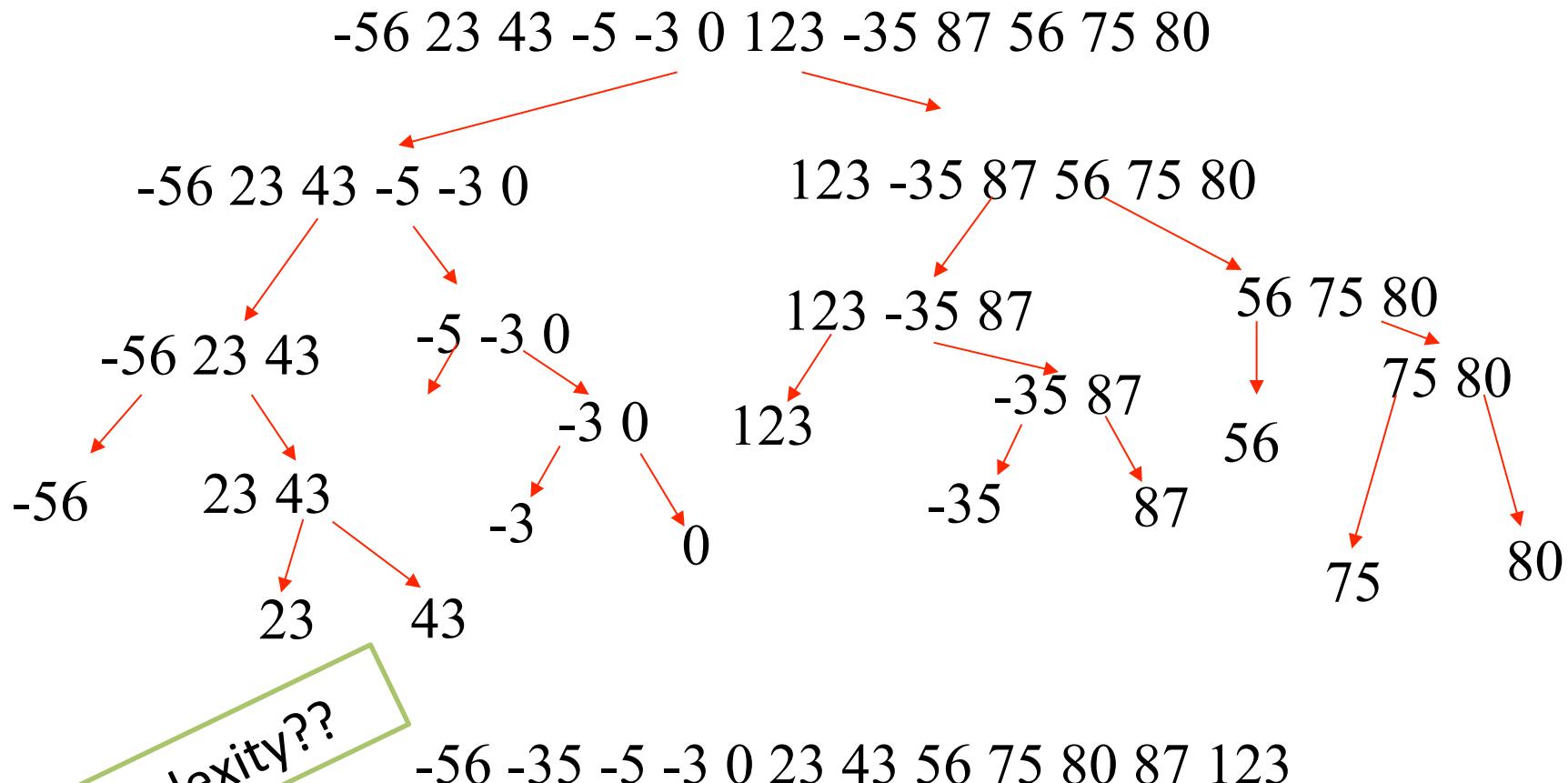
    while(i<=i2 && j<=j2)      //while elements in both lists
    {
        if(a[i]<a[j])
            temp[k++]=a[i++];
        else
            temp[k++]=a[j++];
    }

    while(i<=i2)      //copy remaining elements of the first list
        temp[k++]=a[i++];

    while(j<=j2)      //copy remaining elements of the second list
        temp[k++]=a[j++];

    for(i=i1,j=0;i<=j2;i++,j++)
        a[i]=temp[j];      //Transfer elements from temp[] back to a[]
}
```

Splitting Trace

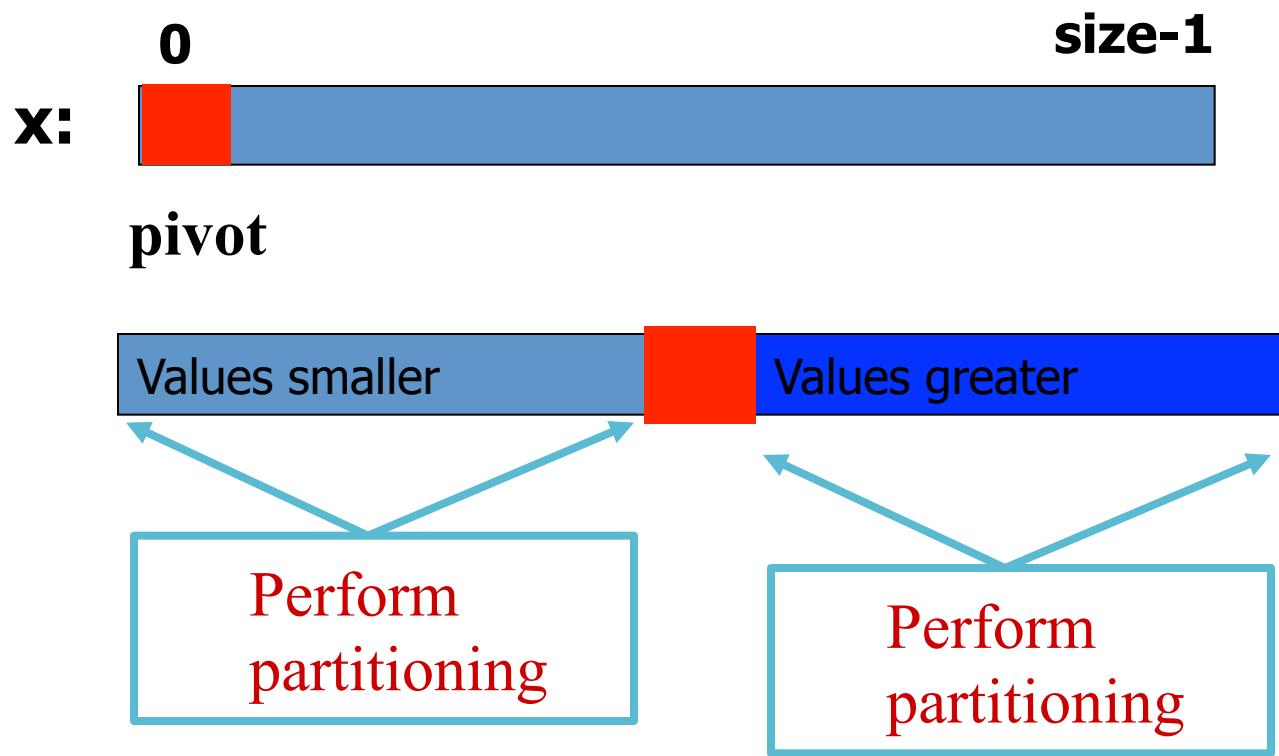


Worst Case: $O(n \cdot \log(n))$

Quicksort

- At every step, we select a *pivot element* in the list (usually the first element).
 - We put the pivot element in the final position of the sorted list.
 - All the elements less than or equal to the pivot element are to the left.
 - All the elements greater than the pivot element are to the right.

Partitioning



Quick Sort program

```
#include <stdio.h>
void quickSort( int[], int, int);
int partition( int[], int, int);

void main()
{
    int i,a[] = { 7, 12, 1, -2, 0, 15, 4, 11, 9};
    printf("\n\nUnsorted array is: ");
    for(i = 0; i < 9; ++i)
        printf(" %d ", a[i]);
    quickSort( a, 0, 8);
    printf("\n\nSorted array is: ");
    for(i = 0; i < 9; ++i)
        printf(" %d ", a[i]);
}

void quickSort( int a[], int l, int r)
{
    int j;
    if( l < r ) { // divide and conquer
        j = partition( a, l, r);
        quickSort( a, l, j-1);
        quickSort( a, j+1, r);
    }
}
```

Quick Sort program

```
int partition( int a[], int l, int r)
{
    int pivot, i, j, t;
    pivot = a[l];
    i = l;
    j = r+1;
    while( 1) {
        do {
            ++i;
        } while(a[i]<=pivot && i<=r);

        do {
            --j;
        } while( a[j] > pivot );
        if( i >= j ) break;
        t = a[i];
        a[i] = a[j];
        a[j] = t;
    }
    t = a[l];
    a[l] = a[j];
    a[j] = t;
    return j;
}
```

Trace of Partitioning

Input: 45 -56 78 90 -3 -6 123 0 -3 45 69 68

45 -56 78 90 -3 -6 123 0 -3 45 69 68

-6	-56	-3	0	-3	45	123	90	78	45	69	68
-56	-6	-3	0	-3	68	90	78	45	69	123	
-3	0	-3			45	68	78	90	69		
-3	0				69		78			90	

Output: -56 -6 -3 -3 0 45 45 68 69 78 90 123

Time Complexity

- Partitioning with n elements.

- No. of comparisons:

$$n-1$$

Choice of pivot element affects the time complexity.

- Worst Case Performance:

$$(n-1)+(n-2)+(n-3)+\dots\dots\dots+1 = n(n-1)/2$$

- Best Case performance:

$$(n-1)+2((n-1)/2-1)+4(((n-1)/2-1)-1)/2-1) \dots k \text{ steps}$$

$$= O(n \cdot \log(n))$$

$$2^k=n$$

Searching an Array: Linear and Binary Search

Searching

- Check if a given element (**key**) occurs in the array.

Linear Search

- Basic idea:
 - Start at the beginning of the array.
 - Inspect every element to see if it matches the key.

```
/* If key appears in a[0..size-1], print its location pos, where a[pos] == key. Else print unsuccessful search */
#include <stdio.h>
int main()
{
    int size,a[100],key,i,pos;

    printf("Enter the number of elements: ");
    scanf("%d",&size);
    printf("Enter the elements: ");
    for(i=0;i<size;i++)
        scanf("%d",&a[i]);
    printf("Enter the key element: ");
    scanf("%d",&key);
    for(pos=-1,i=0;i<size;i++) { /* initializing pos as unsuccessful search*/
        if(a[i]==key) {
            pos=i;
            break;
        }
    }
    if(pos==-1)
        printf("Unsuccessful search\n");
    else
        printf("The element is present at %d position\n",pos+1);
    return 0;
}
```

```

/* If key appears in a[0..size-1], print its location pos, where a[pos] == key. Else print unsuccessful search */
#include <stdio.h>
#include <stdlib.h>      /* for exit() function */
#define SIZE 100

void main()
{
    int size,a[SIZE],key,i,pos;

    printf("Enter the number of elements: ");
    scanf("%d",&size);
    if(size>SIZE) { /* size is a variable, SIZE is not!! */
        printf("Array Size error!!! I am exiting .... \n");
        exit(0);
    }
    printf("Enter the elements: ");
    for(i=0;i<size;i++)a
        scanf("%d",&a[i]);
    printf("Enter the key element: ");
    scanf("%d",&key);
    for(pos=-1,i=0;i<size;i++) { /* initializing pos as unsuccessful search*/
        if(a[i]==key) {
            pos=i;
            break;
        }
    }
    (pos== -1)? printf("Unsuccessful search\n"):printf("The element is present at %d position\n",pos+1);
}

```

Linear Search

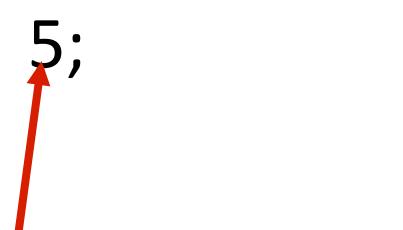
1. If function type is **void** then nothing is to be returned from the function.
2. It is always good to assign a return data type (including void) with each function.
3. **exit()** will exit from the program without executing remaining statements of the program. This needs **stdlib.h** as header file.

Linear Search

```
int x[ ]= {12,-3,78,67,6,50,19,10} ;
```

- Trace the following calls :

search 6;  Returns 5

search 5;  Unsuccessful search

Linear Search

- Basic idea:
 - Start at the beginning of the array.
 - Inspect every element to see if it matches the key.
- Time complexity:
 - A measure of how long an algorithm takes to run.
 - If there are n elements in the array:
 - Best case:
match found in first element (1 search operation)
 - Worst case:
no match found, or match found in the last element (n search operations)
 - Average:
 $(n + 1) / 2$ search operations

Search on Sorted List

~~int x[]= {12,-3,78,67,6,50,19,10,11};~~



int x[]= {-3,6,10,11,12,19,50,67,78} ;

- Trace the following calls :

search 6;

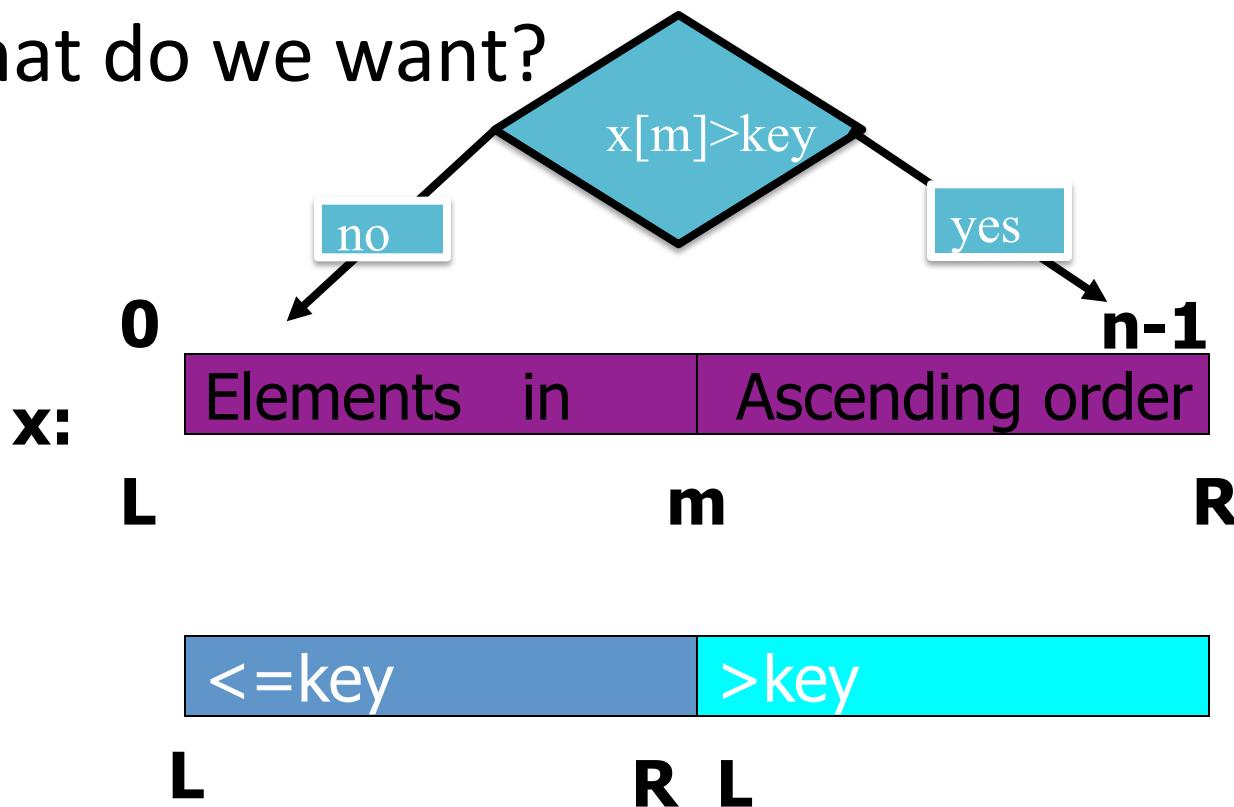
search 5;

Binary Search

- Binary search works if the array is sorted.
 - Look for the target in the middle.
 - If you don't find it, you can ignore half of the array, and repeat the process with the other half.
- In every step, we reduce the number of elements to search in by half.

The Basic Strategy

- What do we want?



- Look at $[(L+R)/2]$. Move L or R to the middle depending on test.
- Repeat search operation in the reduced interval.

Binary Search

```
/* If key appears in x[0..size-1], prints its location pos where  
x[pos]==key. If not found, print -1 */
```

```
int main ()  
{  
    int x[100],size,key;  
    int L, R, mid;  
    _____;  
  
    while ( _____ )  
    {  
        _____;  
    }  
    _____;  
}
```

The basic search iteration

```
/* If key appears in x[0..size-1], prints its location pos where x[pos]==key. If  
not found, print -1 */  
  
int main ()  
{  
    int x[100],size,key;  
    int L, R, mid;  
    _____;  
    while ( _____ )  
    {  
        mid = (L + R) / 2;  
        if (x[mid] > key)  
            R = mid;  
        else L = mid;  
    }  
    _____;  
}
```

Loop termination

```
/* If key appears in x[0..size-1], prints its location pos where x[pos]==key. If  
not found, print -1 */
```

```
int main ()  
{  
    int x[100],size,key;  
    int L, R, mid;  
    _____;  
while ( L+1 != R )  
{  
    mid = (L + R) / 2;  
    if (x[mid] <= key)  
        L = mid;  
    else R = mid;  
}  
_____;  
}
```

Print result

```
/* If key appears in x[0..size-1], prints its location pos where x[pos]==key. If
not found, print -1 */

int main ()
{
    int x[100],size,key;
    int L, R, mid;
    _____;
    while ( L+1 != R )
    {
        mid = (L + R) / 2;
        if (x[mid] <= key)
            L = mid;
        else R = mid;
    }
    if (L >= 0 && x[L] == key) printf("%d",L);
    else printf("-1");
}
```

Initialization

```
/* If key appears in x[0..size-1], prints its location pos where x[pos]==key. If not found, print -1 */
```

```
int main ()
{
    int x[100],size,key;
    int L, R, mid;
    _____;
    L = -1;  R = size;
    while ( L+1 != R )
    {
        mid = (L + R) / 2;
        if (x[mid] <= key)
            L = mid;
        else R = mid;
    }
    if (L >= 0 && x[L] == key) printf("%d",L);
    else printf("-1");
}
```

Complete C Program

```
/* If key appears in x[0..size-1], prints its location pos where x[pos]==key. If not found, print -1 */
```

```
void main ()  
{  
    int x[100],size,key;  
    int L, R, mid;  
    printf("Enter the number of elements: ");  
    scanf("%d",&size);  
    printf("Enter the elements: ");  
    for(i=0;i<size;i++)  
        scanf("%d",&a[i]);  
    printf("Enter the key element: ");  
    scanf("%d",&key);  
  
    L = -1;  R = size;  
    while ( L+1 != R )  {  
        mid = (L + R) / 2;  
        if (x[mid] <= key)  
            L = mid;  
        else R = mid;  
    }  
    if  (L >= 0  &&  x[L] == key)  printf("%d",L);  
    else printf("-1");  
}
```

Binary Search Examples

Sorted array

```
-17 -5  3   6   12  21  45  63  50
```

Trace :

binsearch 3;

binsearch 145;

binsearch 45;

L= -1; R= 9; x[4]=12;

L= -1; R=4; x[1]=-5;

→ L= 1; R=4; x[2]=3;

L=2; R=4; x[3]=6;

L=2; R=3; return L;

Is it worth the trouble ?

- Suppose there are 1000 elements.
- Ordinary search
 - If key is a member of x , it would require 500 comparisons on the average.
- Binary search
 - after 1st compare, left with 500 elements.
 - after 2nd compare, left with 250 elements.
 - After at most 10 steps, you are done.

What is best case?
What is worst case?

Time Complexity

- If there are n elements in the array.
 - Number of searches required:
 $\log_2 n$
 - For $n = 64$ (say)
 - Initially, list size = 64.
 - After first compare, list size = 32.
 - After second compare, list size = 16.
 - After third compare, list size = 8.
 -
 - After sixth compare, list size = 1.
- $2^k = n,$
Where k is the number of steps.
- $$\log_2 64 = 6$$

Homework

**Modify the algorithm by
checking equality with $x[\text{mid}]$.**