# CS11001/CS11002 Programming and Data Structures (PDS) (Theory: 3-0-0) 

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## An Example: Random Number Generation



## Passing Arrays to a Function

- An array name can be used as an argument to a function.
- Permits the entire array to be passed to the function.
- Array name is passed as the parameter, which is effectively the address of the first element.
- Rules:
- The array name must appear by itself as argument, without brackets or subscripts.
- The corresponding formal argument is written in the same manner.
- Declared by writing the array name with a pair of empty brackets.
- Dimension or required number of elements to be passed as a separate parameter.


## Example 1: Minimum of a set of numbers



## Parameter Passing mechanism

- When an array is passed to a function, the values of the array elements are not passed to the function.
- The array name is interpreted as the address of the first array element.
- The formal argument therefore becomes a pointer to the first array element.
- When an array element is accessed inside the function, the address is calculated using the formula stated before.
- Changes made inside the function are thus also reflected in the calling program.


## Parameter Passing mechanism

- Passing parameters in this way is called call-by-reference.
- Normally parameters are passed in C using call-by-value.
- Basically what it means?
- If a function changes the values of array elements, then these changes will be made to the original array that is passed to the function.
- This does not apply when an individual element is passed on as argument.


## Example: Average of numbers



## Call by Value and Call by Reference

- Call by value
- Copy of argument passed to function
- Changes in function do not effect original
- Use when function does not need to modify argument
- Avoids accidental changes
- Call by reference
- Passes original argument
- Changes in function effect original
- Only used with trusted functions


## Example: Max Min function



## Scope of a variable



## Storage Class of Variables

## What is Storage Class?

- It refers to the permanence of a variable, and its scope within a program.
- Four storage class specifications in C:
- Automatic: auto
- External: extern
- Static: static
- Register: register


## Automatic Variables

- These are always declared within a function and are local to the function in which they are declared.
- Scope is confined to that function.
- This is the default storage class specification.
- All variables are considered as auto unless explicitly specified otherwise.
- The keyword auto is optional.
- An automatic variable does not retain its value once control is transferred out of its defining function.


## auto: Example



## Static Variables

- Static variables are defined within individual functions and have the same scope as automatic variables.
- Unlike automatic variables, static variables retain their values throughout the life of the program.
- If a function is exited and re-entered at a later time, the static variables defined within that function will retain their previous values.
- Initial values can be included in the static variable declaration.
- Will be initialized only once.
- An example of using static variable:
- Count number of times a function is called.


## static: Example



## External Variables

- They are not confined to single functions.
- Their scope extends from the point of definition through the remainder of the program.
- They may span more than one functions.
- Also called global variables.
- Alternate way of declaring global variables.
- Declare them outside the function, at the beginning.


## global: Example

\#include <stdio.h>
int count=0;
void print()
\{
printf("Hello World!! ");
count++;
\}
int main()
\{
int $\mathrm{i}=0$;
while(i<10) \{ print();
i++; printf("is printing \%d times.\n",count);
\}
return 0 ;
\}

## static vs global

```
#include <stdio.h>
void print()
{
    static int count=0;
    printf("Hello World!! ");
    count++;
    printf("is printing %d times.\n",count);
}
int main()
{
    int i=0;
    while(i<10) {
        print();
        i++;
    }
    return 0;
}
```

\#include <stdio.h>
int count=0;

## Register Variables

- These variables are stored in high-speed registers within the CPU.
- Commonly used variables like loop variables/counters may be declared as register variables.
- Results in increase in execution speed.
- User can suggest, but the
\#include<stdio.h>
int main()
\{
int sum;
register int count;
for(count=0;count<20;count++)
sum=sum+count;
printf("\nSum of Numbers:\%d", sum);
return(0);
\} allocation is done by the compiler.


## \#include: Revisited

- Preprocessor statement in the following form \#include "filename"
- Filename could be specified with complete path. \#include "/home/pralay/C-header/myfile.h"
- The content of the corresponding file will be included in the present file before compilation and the compiler will compile thereafter considering the content as it is.


## \#include: Revisited



# Variable number of arguments 

- General syntax:
scanf (control string, arg1, arg2, ..., argn); printf (control string, arg1, arg2, ..., argn);


## How is it possible?

## Example: GCD calculation

> int gcd (int A, int B)
> \{
> int temp;
> while (( $B \% A)!=0)$ \{

```
/* Compute the GCD of four numbers */
#include <stdio.h>
int gcd(int A, int B);
void main()
{
}
    int n1, n2, n3, n4, result;
    scanf ("%d %d %d %d", &n1, &n2, &n3, &n4);
    result = gcd (gcd (n1, n2), gcd (n3, n4) );
    printf ("The GCD of %d, %d, %d and %d is %d \n",
    n1, n2, n3, n4, result);
}
```


## Example: GCD calculation



## Recursion

- A process by which a function calls itself repeatedly.
- Either directly.
- X calls X.
- Or cyclically in a chain.
- $X$ calls $Y$, and $Y$ calls $X$.
- Used for repetitive computations in which each action is stated in terms of a previous result.
- fact( $n$ ) $=\mathrm{n}^{*}$ fact ( $\mathrm{n}-1$ )


## Recursion

- For a problem to be written in recursive form, two conditions are to be satisfied:
- It should be possible to express the problem in recursive form.
- The problem statement must include a stopping condition

$$
\begin{aligned}
\text { fact(n) } & =1, & \text { if } n=0 \\
& =n * \text { fact }(n-1), & \text { if } n>0
\end{aligned}
$$

## Recursion

- Examples:
- Factorial:
fact(0) $=1$
fact( $n$ ) $=n *$ fact $(n-1)$, if $n>0$
- GCD:
$\operatorname{gcd}(m, m)=m$
$\operatorname{gcd}(m, n)=\operatorname{gcd}(m-n, n)$, if $m>n$
$\operatorname{gcd}(m, n)=\operatorname{gcd}(n, n-m)$, if $m<n$
- Fibonacci series (1,1,2,3,5,8,13,21,....)
fib (0) = 1
fib (1) $=1$
fib $(\mathrm{n})=$ fib $(\mathrm{n}-1)+$ fib $(\mathrm{n}-2)$, if $\mathrm{n}>1$


## Example 1 :: Factorial



## Mechanism of Execution

- When a recursive program is executed, the recursive function calls are not executed immediately.
- They are kept aside (on a stack) until the stopping condition is encountered.
- The function calls are then executed in reverse order.


# Advantage of Recursion :: Calculating fact(5) 

- First, the function calls will be processed:
fact(5) $=5$ * fact(4)
fact(4) $=4$ * fact(3)
fact(3) $=3 *$ fact(2)
fact(2) $=2$ * fact(1)
fact(1) $=1$ * fact(0)
- The actual values return in the reverse order:
fact( 0 ) $=1$
fact(1) $=1$ * $1=1$
fact(2) $=2 * 1=2$
fact(3) $=3$ * $2=6$
fact(4) $=4 * 6=24$
Fact(5) $=5 * 24=120$


## Facts on fact

$-5!=5 * 4 * 3 * 2 * 1$

- Notice that
- $5!=5$ * 4 !
- 4 ! = 4 * 3 ! ...
- Can compute factorials recursively
-Solve base case (1! = $0!=1$ ) then plug in
- 2! = 2 * 1! = 2 * $1=2$;
- $3!=3$ * $2!=3 * 2=6$;


## Example 2 :: Fibonacci series



## Execution of Fibonacci number

- Fibonacci number fib(n) can be defined as:
- The successive Fibonacci numbers are:

$$
0,1,1,2,3,5,8,13,21, \ldots . .
$$



## Set of recursive calls for fibonacci (3)



## Inefficiency of Recursion

- How many times the function is called when evaluating $f(4)$ ?
- Same thing is
 computed several times.


## Performance Tip

- Avoid Fibonacci-style recursive programs which result in an exponential "explosion" of calls.


## Example 3: Towers of Hanoi Problem



- The problem statement:
- Initially all the disks are stacked on the LEFT pole.
- Required to transfer all the disks to the RIGHT pole.
- Only one disk can be moved at a time.
- A larger disk cannot be placed on a smaller disk.


## Recursion is implicit

## - General problem of $n$ disks.

- Step 1:
- Move the top (n-1) disks from LEFT to CENTER.
- Step 2:
- Move the largest disk from LEFT to RIGHT.
- Step 3:
- Move the (n-1) disks from CENTER to RIGHT.


## Recursive C code: Towers of Hanoi

```
#include <stdio.h>
void transfer (int n, char from, char to, char temp);
int main()
{
    int n; /* Number of disks */
    scanf ("%d", &n);
    transfer (n, 'L', 'R', 'C');
    return 0;
}
void transfer (int n, char from, char to, char temp)
{
    if (n>0) {
        transfer (n-1, from, temp,to);
        printf ("Move disk %d from %c to %c \n", n, from, to);
        transfer (n-1, temp, to, from);
    }
    return;
}
```


## Towers of Hanoi: Example Run


4
Move disk 1 from $L$ to $C$
Move disk 2 from $L$ to $R$
Move disk 1 from $C$ to $R$
Move disk 3 from $L$ to $C$
Move disk 1 from $R$ to $L$
Move disk 2 from $R$ to $C$
Move disk 1 from $L$ to $C$
Move disk 4 from $L$ to $R$
Move disk 1 from $C$ to $R$
Move disk 2 from $C$ to $L$
Move disk 1 from $R$ to $L$
Move disk 3 from $C$ to $R$
Move disk 1 from $L$ to $C$
Move disk 2 from $L$ to $R$
Move disk 1 from $C$ to $R$


## Recursion vs. Iteration

- Repetition
- Iteration: explicit loop
- Recursion: repeated function calls
- Termination
- Iteration: loop condition fails
- Recursion: base case recognized
- Both can have infinite loops
- Balance
- Choice between performance (iteration) and good software engineering (recursion)


## Performance Tip

- Avoid using recursion in performance situations. Recursive calls take time and consume additional memory.


## How are function calls implemented?

- In general, during program execution
- The system maintains a stack in memory.
- Stack is a last-in first-out structure.
- Two operations on stack, push and pop.
- Whenever there is a function call, the activation record gets pushed into the stack.
- Activation record consists of the return address in the calling program, the return value from the function, and the local variables inside the function.
- At the end of function call, the corresponding activation record gets popped out of the stack.


## At the system



## Example:: main() calls fact(3)



TRACE OF THE STACK DURING EXECUTION


## Homework

## Trace of Execution for Fibonacci Series

## Sorting: the basic problem

- Given an array

$$
x[0], x[1], \ldots, x[\text { size }-1]
$$

reorder entries so that

$$
x[0]<=x[1]<=\ldots<=x[\text { size }-1]
$$

- List is in non-decreasing order.
- We can also sort a list of elements in nonincreasing order.


## Sorting Problem

- What we want : Data sorted in order
- Input: A list of elements
- Output: A list of elements in sorted (non-increasing/nondecreasing) order

- Original list:
$-10,30,20,80,70,10,60,40,70$
- Sorted in non-decreasing order:
- $10,10,20,30,40,60,70,70,80$
- Sorted in non-increasing order:
- 80, 70, 70, 60, 40, 30, 20, 10, 10


## Example



## Selection Sort

- General situation :

|  | 0 | k |  |
| :---: | :---: | :---: | :---: |
| X | smallest elements, sorted remainder, unsorted |  |  |

- Steps :
- Find smallest element, mval, in $x[k+1$..size-1]
- Swap smallest element with $\mathrm{x}[\mathrm{k}-1]$,
- Increase k.

swap


## Subproblem: Find smallest element

```
/* Yield location of smallest element in x[k .. size-1] and store in pos*/
int j, pos;
pos = k; /* assume first element is the smallest element */
for (j=k+1; j<size; j++) {
    if (x[j] < x[pos]) { /* x[pos] is the smallest element as of now */
        pos = j;
    }
}
printf("%d",pos);
```


## Subproblem: Swap with smallest element

```
/* Yield location of smallest element in x[k .. size-1] and store in pos*/
int j, pos;
pos = k; /* assume first element is the smallest element */
for (j=k+1; j<size; j++) {
        if (x[j] < x[pos]) { /* x[pos] is the smallest element as of now */
                        pos = j;
        }
}
printf("%d",pos);
if (x[pos] < x[k]) {
    temp = x[k]; /* swap content of }\textrm{x}[\textrm{k}]\mathrm{ and }\textrm{x}[\textrm{pos}] *
    x[k] = x[pos];
    x[pos] = temp;
}
```

```
#include <stdio.h>
    /* Sort x[0..size-1] in non-decreasing order */
int main()
{
    int k,j,pos,x[100],size,temp;
    printf("Enter the number of elements: ");
    scanf("%d",&size);
    printf("Enter the elements: ");
    for(k=0;k<size;k++)
        scanf("%d",&x[k]);
    for (k=0; k<size-1; k++) {
        pos = k; /* assume first element is the smallest element */
        for (j=k+1; j<size; j++) {
            if (x[j]<x[pos]) /* x[pos] is the smallest element as of now */
                pos = j;
        }
        if (x[pos] < x[k]) {
            temp = x[k];
            x[k] = x[pos];
            x[pos] = temp;
        }
    }
    for(k=0;k<size;k++) /* print the sorted (non-decreasing) list */
        printf("%d ",x[k]);
    return 0;
```


## Example



## Analysis

How many steps are required?
Let us assume there are $n$ elements (size=n). Each statement executes in constant time.

To read
for loop will take of the order of $n$ time
Sorting
When $\mathrm{k}=0$, in worst case
( $\mathrm{n}-1$ ) comparisons to find minimum.
When $\mathrm{k}=1$, in worst case
( $n-2$ ) comparisons to find minimum.

$$
(n-1)+(n-2)+\ldots \ldots+1=n(n-1) / 2
$$

To print

```
for(k=0;k<size;k++)
    scanf("%d",&x[k]);
for (k=0; k<size-1; k++) {
    pos = k;
    for (j=k+1; j<size; j++) {
        if (x[j] < x[pos])
                pos = j;
    }
    temp = x[k];
    x[k] = x[pos];
    x[pos] = temp;
}
for(k=0;k<size;k++)
    printf("%d ",x[k]);
```

            for loop will take of the order of \(n\) time
    Total time $=\mathbf{2} \times$ order of $\mathbf{n}+$ order of $\mathbf{n}^{\mathbf{2}}=$ order of $\mathbf{n}^{\mathbf{2}}$

## Analysis

How many steps are required?

Let us assume there are $n$ elements (size=n). Each statement executes in constant time.

```
for(k=0;k<size;k++)
    scanf("%d",&x[k]);
for (k=0; k<size-1; k++) {
    pos = k;
    for (j=k+1; j<size; j++) {
        if (x[j] < x[pos])
                                    pos = j;
    }
    temp = x[k];
    x[k] = x[pos];
    x[pos] = temp;
}
for(k=0;k<size;k++)
    printf("%d ",x[k]);
```


## Insertion Sort

## - General situation :



Compare and Shift till $x[i]$ is larger.
i


0
size-1

## Insertion Sorting

```
int list[100], size;
```

$\qquad$

``` _;
for (i=1; i<size; i++) {
        item = list[i] ;
    for (j=i-1; (j>=0)&& (list[j] > item); j--)
        list[j+1] = list[j];
    list[j+1] = item ;
}
```

$\qquad$

``` ;
```


## Complete Insertion Sort

```
#include <stdio.h>
    /* Sort x[0..size-1] in non-decreasing order */
#define SIZE 100
int main()
{
    int i,j,x[SIZE],size,temp;
    printf("Enter the number of elements: ");
    scanf("%d",&size);
    printf("Enter the elements: ");
    for(i=0;i<size;i++)
        scanf("%d",&x[i]);
    for (i=1; i<size; i++) {
        temp = x[i];
        for (j=i-1; (j>=0)&& (x[j] > temp); j--)
                x[j+1] = x[j];
            x[j+1] = temp ;
    }
    for(i=0;i<size;i++) /* print the sorted (non-decreasing) list */
        printf("%d ",x[i]);
    return 0;
```


## Insertion Sort Example



## Time Complexity

- Number of comparisons and shifting:

Worst Case?

$$
1+2+3+\ldots \ldots+(n-1)=n(n-1) / 2
$$

- Best Case?
$1+1+\ldots \ldots(n-1)$ times $=(n-1)$

