

Searching an Array: Linear and Binary Search

Searching

- Check if a given element (**key**) occurs in the array.
- Two methods to be discussed:
 - If the array elements are unsorted.
 - Linear search
 - If the array elements are sorted.
 - Binary search

Linear Search

- **Basic idea:**
 - Start at the beginning of the array.
 - Inspect every element to see if it matches the key.
- **Time complexity:**
 - A measure of how long an algorithm takes to run.
 - If there are **n** elements in the array:
 - **Best case:**
match found in first element (**1 search operation**)
 - **Worst case:**
no match found, or match found in the last element (**n search operations**)
 - **Average:**
($n + 1$) / 2 search operations

Contd.

```
/* If key appears in a[0..size-1], return its location, pos, s.t.  
 a[pos] == key. If key is not found, return -1 */  
int linear_search (int a[], int size, int key)  
{  
    int pos = 0;  
    while ((pos < size) && (a[pos] != key))  
        pos++;  
    if (pos < n)  
        return pos; /* Return the position of match */  
    return -1; /* No match found */  
}
```

Contd.

```
int x[ ]= {12, -3, 78, 67, 6, 50, 19, 10} ;
```

- Trace the following calls :

```
search (x, 8, 6) ;
```

Returns 4

```
search (x, 8, 5) ;
```



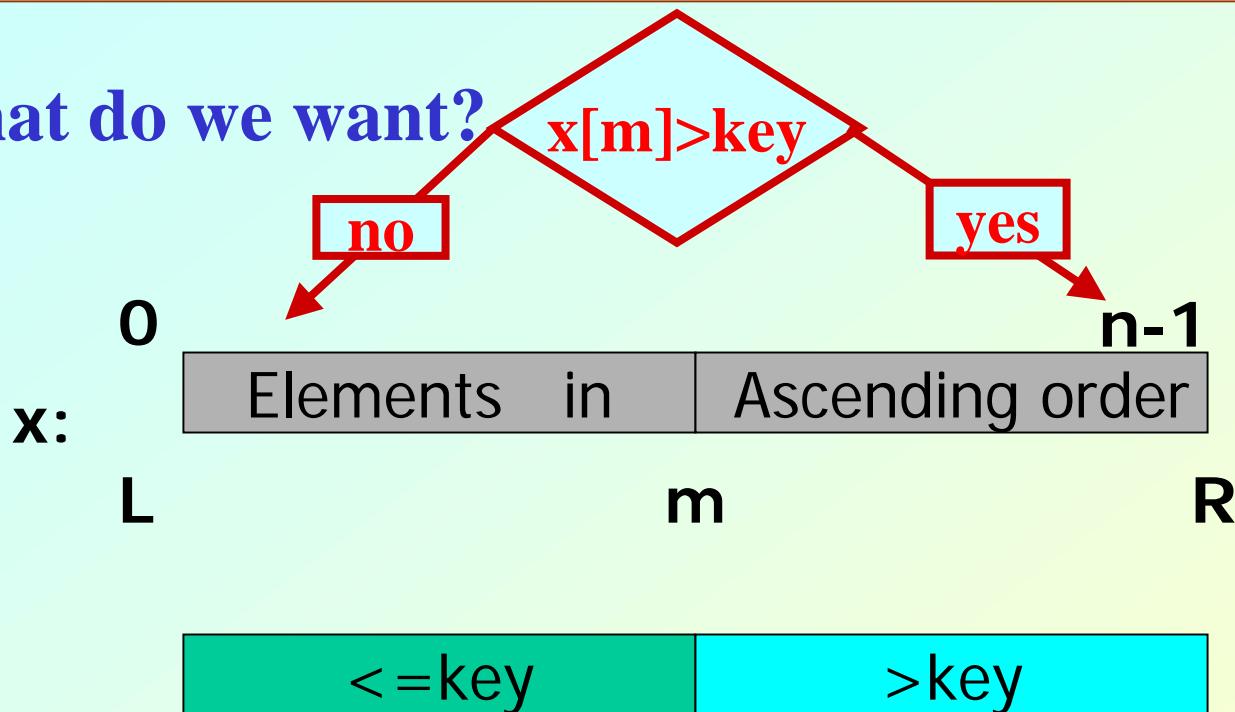
Returns -1

Binary Search

- **Binary search works if the array is sorted.**
 - Look for the target in the middle.
 - If you don't find it, you can ignore half of the array, and repeat the process with the other half.
- **In every step, we reduce the number of elements to search in by half.**

The Basic Strategy

- What do we want?



- Look at $x[(L+R)/2]$. Move L or R to the middle depending on test.
- Repeat search operation in the reduced interval.

Contd.

/* If **key** appears in x[0..size-1], return its location, pos s.t.
x[pos]==key. If not found, return -1 */

```
int bin_search (int x[], int size, int key)
{
    int L, R, mid;
    _____;
    while ( _____ )
    {
        _____;
    }
    _____;
}
```

The basic search iteration

```
/* If key appears in x[0..size-1], return its location, pos s.t. x[pos]==key. If  
not found, return -1 */
```

```
int bin_search (int x[], int size, int key)  
{  
    int L, R, mid;  
    _____;  
    while ( _____ )  
    {  
        mid = (L + R) / 2;  
        if (x[mid] > key)  
            R = mid;  
        else L = mid;  
    }  
    _____;  
}
```

Loop termination

```
/* If key appears in x[0..size-1], return its location, pos s.t. x[pos]==key. If  
not found, return -1 */
```

```
int bin_search (int x[], int size, int key)  
{  
    int L, R, mid;  
    _____;  
    while ( L+1 != R )  
    {  
        mid = (L + R) / 2;  
        if (x[mid] <= key)  
            L = mid;  
        else R = mid;  
    }  
    _____;  
}
```

Return result

```
/* If key appears in x[0..size-1], return its location, pos s.t. x[pos]==key. If  
not found, return -1 */
```

```
int bin_search (int x[], int size, int key)
```

```
{
```

```
    int L, R, mid;
```

```
    _____;
```

```
    while ( L+1 != R )
```

```
{
```

```
        mid = (L + R) / 2;
```

```
        if (x[mid] <= key)
```

```
            L = mid;
```

```
        else R = mid;
```

```
}
```

```
    if (L >= 0 && x[L] == key) return L;
```

```
    else return -1;
```

```
}
```

Initialization

```
/* If key appears in x[0..size-1], return its location, pos s.t. x[pos]==key. If  
not found, return -1 */
```

```
int bin_search (int x[], int size, int key)
```

```
{
```

```
    int L, R, mid;
```

```
    L = -1;  R = size;
```

```
    while ( L+1 != R )
```

```
{
```

```
        mid = (L + R) / 2;
```

```
        if (x[mid] <= key)
```

```
            L = mid;
```

```
        else R = mid;
```

```
}
```

```
    if (L >= 0 && x[L] == key)  return L;
```

```
    else return -1;
```

```
}
```

Binary Search Examples

Sorted array

-17 -5 3 6 12 21 45 63 50

Trace :

binsearch (x, 9, 3); →

binsearch (x, 9, 145);

binsearch (x, 9, 45);

L= -1; R= 9; x[4]=12;

L= -1; R=4; x[1]=-5;

L= 1; R=4; x[2]=3;

L=2; R=4; x[3]=6;

L=2; R=3; return L;

We may modify the algorithm by checking equality with x[mid].

Is it worth the trouble ?

- Suppose there are **1000 elements.**
- **Ordinary search**
 - If **key** is a member of **x**, it would require **500 comparisons on the average.**
- **Binary search**
 - after **1st compare, left with 500 elements.**
 - after **2nd compare, left with 250 elements.**
 - **After at most 10 steps, you are done.**

Time Complexity

- If there are n elements in the array.

- Number of searches required:

$$\log_2 n \leftarrow$$

- For $n = 64$ (say).

- Initially, list size = 64.
 - After first compare, list size = 32.
 - After second compare, list size = 16.
 - After third compare, list size = 8.
 -
 - After sixth compare, list size = 1.

$$2^k = n,$$

Where k is the number of steps.

$$\log_2 64 = 6$$

Sorting: the basic problem

- Given an array

$x[0], x[1], \dots, x[\text{size}-1]$

reorder entries so that

$x[0] \leq x[1] \leq \dots \leq x[\text{size}-1]$

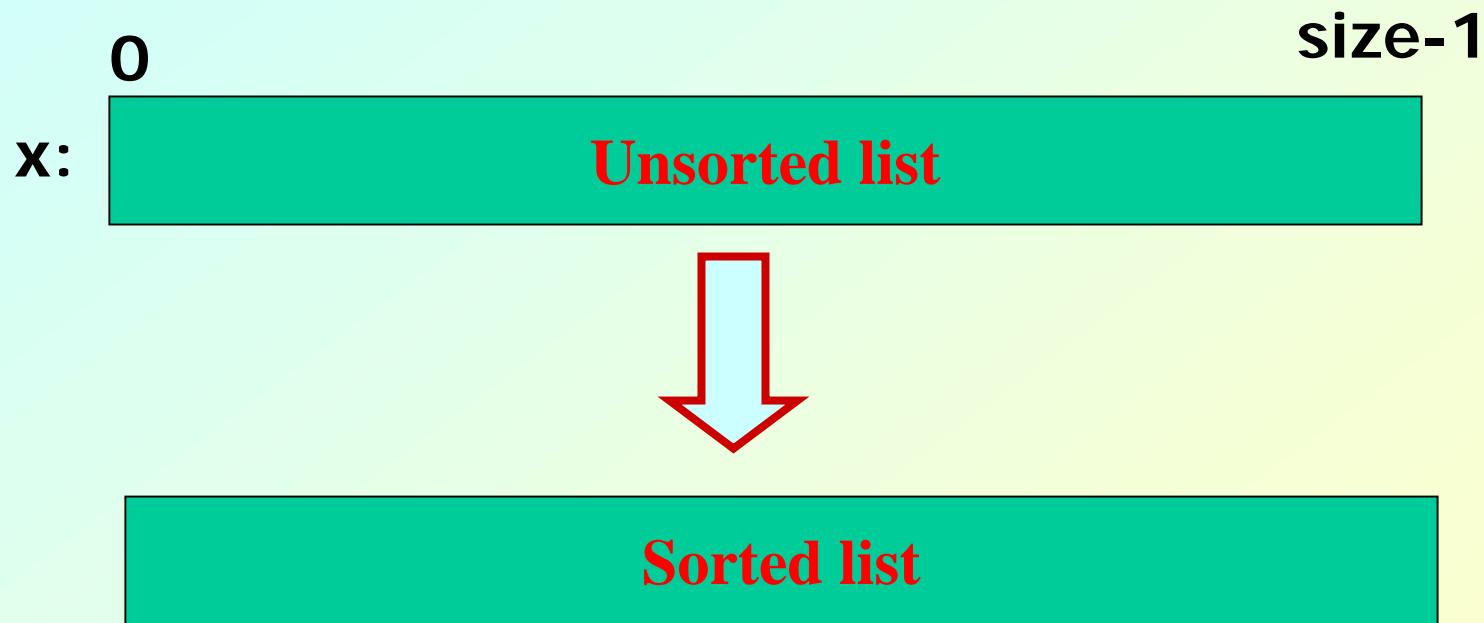
- List is in non-decreasing order.
- We can also sort a list of elements in non-increasing order.

Example

- **Original list:**
 - 10, 30, 20, 80, 70, 10, 60, 40, 70
- **Sorted in non-decreasing order:**
 - 10, 10, 20, 30, 40, 60, 70, 70, 80
- **Sorted in non-increasing order:**
 - 80, 70, 70, 60, 40, 30, 20, 10, 10

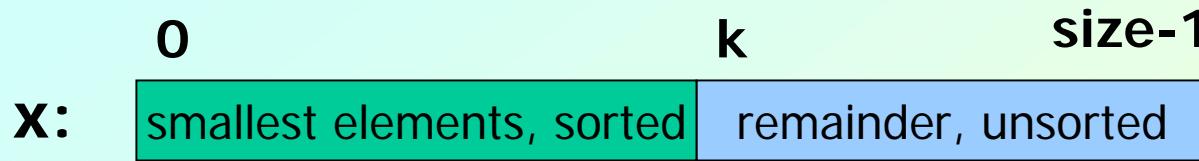
Sorting Problem

- What we want : Data sorted in order



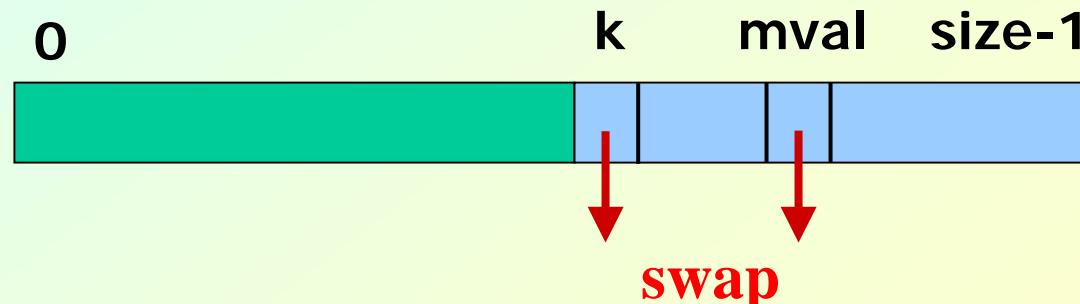
Selection Sort

- General situation :



- Step :

- Find smallest element, **mval**, in $x[k..size-1]$
- Swap smallest element with $x[k]$, then increase k.



Subproblem

```
/* Yield location of smallest element in x[k .. size-1];*/
```

```
int min_loc (int x[ ], int k, int size)
```

```
{
```

```
    int j, pos; /* x[pos] is the smallest element found so far */
```

```
    pos = k;
```

```
    for (j=k+1; j<size; j++)
```

```
        if (x[j] < x[pos])
```

```
            pos = j;
```

```
    return pos;
```

```
}
```

The main sorting function

```
/* Sort x[0..size-1] in non-decreasing order */
```

```
int selsort (int x[], int size)
{
    int k, m;
    for (k=0; k<size-1; k++)
    {
        m = min_loc(x, k, size);
        temp = a[k];
        a[k] = a[m];
        a[m] = temp;
    }
}
```

Example

x: 

x: 

x: 

x: 

x: 

x: 

x: 

x: 

Analysis

- How many steps are needed to sort n things ?
 - Total number of steps proportional to n^2
 - No. of comparisons?

$$(n-1)+(n-2)+\dots+1 = n(n-1)/2$$

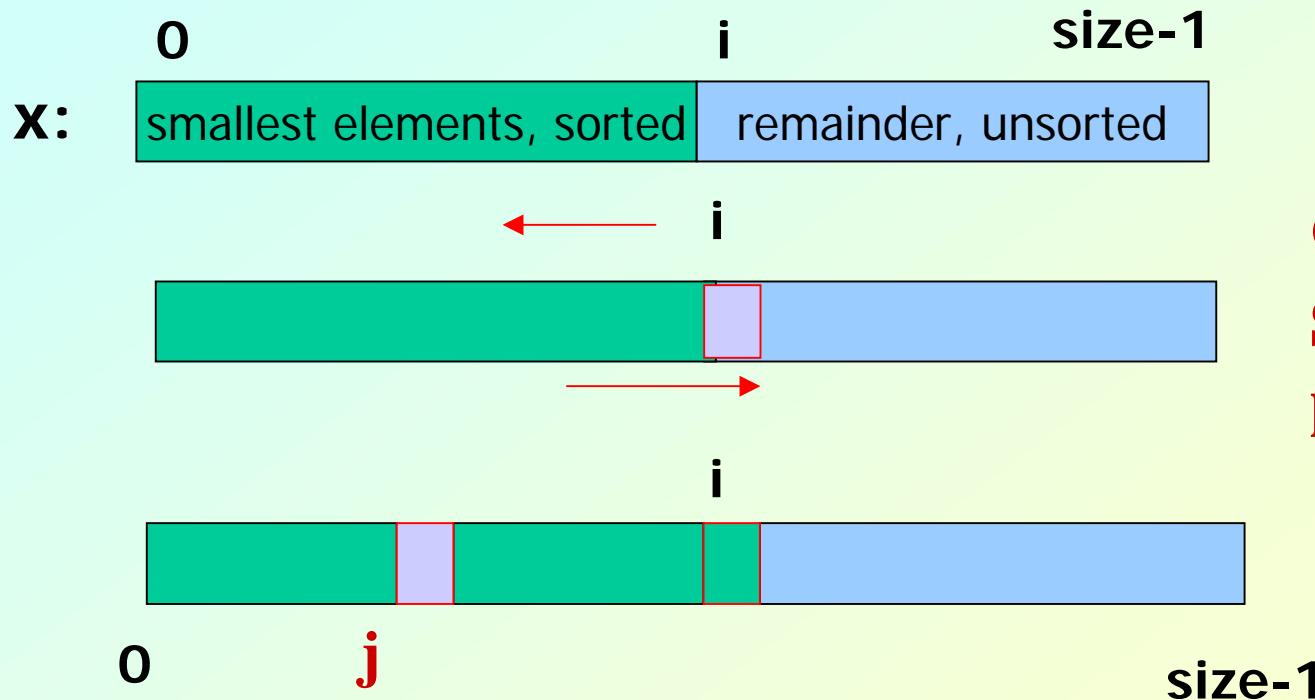


Of the order of n^2

- Worst Case? Best Case? Average Case?

Insertion Sort

- General situation :



Compare and Shift till $x[i]$ is larger.

Insertion Sorting

```
void InsertSort (int list[], int size)
{
    for (i=1; i<size; i++)
    {
        item = list[i] ;
        for (j=i-1; (j>=0)&& (list[j] > i); j--)
            list[j+1] = list[j];
        list[j+1] = item ;
    }
}
```

Insertion Sort

```
#define MAXN 100
void InsertSort (int list[MAXN], int size) ;
main () {
    int index, size;
    int numbers[MAXN];
    /* Get Input */
    size = readarray (numbers) ;
    printarray (numbers, size) ;
    InsertSort (numbers, size) ;
    printarray (numbers, size) ;
}
```

Time Complexity

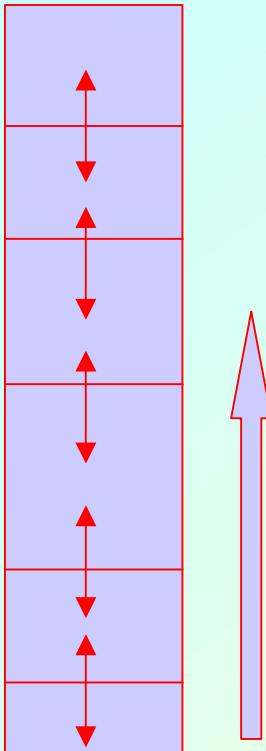
- Number of comparisons and shifting:
 - Worst Case?

$$1+2+3+\dots+(n-1) = n(n-1)/2$$

- Best Case?

$$1+1+\dots \text{ (n-1) times} = (n-1)$$

Bubble Sort



In every iteration heaviest element
drops at the bottom.

The bottom moves upward.

Bubble Sort

```
#include <stdio.h>
```

```
void swap(int *x,int *y)
```

```
{
```

```
    int tmp=*x;
```

```
    *x=*y;
```

```
    *y=tmp;
```

```
}
```

```
void bubble_sort(int x[],int n)
{
    int i,j;
    for(i=n-1;i>0;i--)
        for(j=0;j<i;j++)
            if(x[j]>x[j+1]) swap(&x[j],&x[j+1]);
}
```

Contd.

```
int main(int argc, char *argv[])
{
    int x[ ]={-45,89,-65,87,0,3,-23,19,56,21,76,-50};
    int i;
    for(i=0;i<12;i++) printf("%d ",x[i]);
    printf("\n");
    bubble_sort(x,12);
    for(i=0;i<12;i++) printf("%d ",x[i]);
    printf("\n");
}
```

-45 89 -65 87 0 3 -23 19 56 21 76 -50
-65 -50 -45 -23 0 3 19 21 56 76 87 89

Time Complexity

- Number of comparisons :

- Worst Case?

$$1+2+3+ \dots + (n-1) = n(n-1)/2$$

- Best Case?

- same.**

How do you make best case with $(n-1)$ comparisons only?

Some Efficient Sorting Algorithms

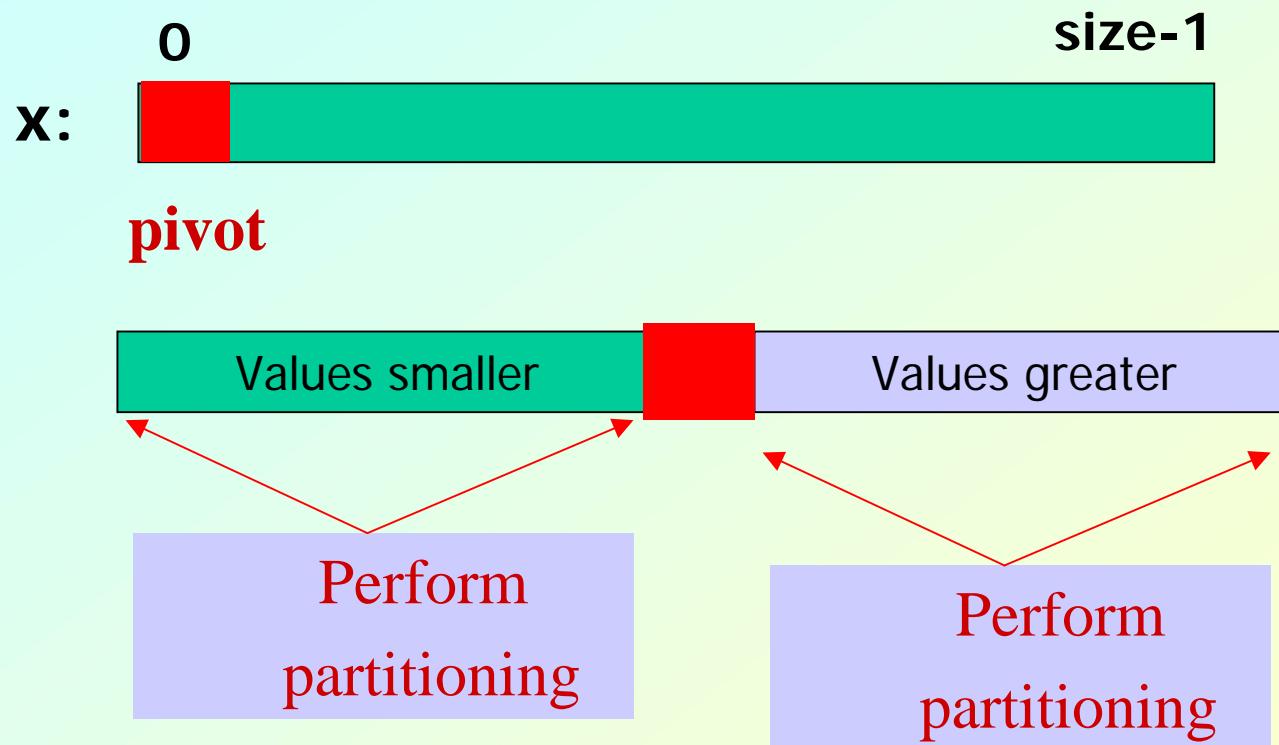
- Two of the most popular sorting algorithms are based on **divide-and-conquer** approach.
 - Quick sort
 - Merge sort
- Basic concept:

```
sort (list)
{
    if the list has length greater than 1
    {
        Partition the list into lowlist and highlist;
        sort (lowlist);
        sort (highlist);
        combine (lowlist, highlist);
    }
}
```

Quicksort

- At every step, we select a pivot element in the list (usually the first element).
 - We put the pivot element in the final position of the sorted list.
 - All the elements less than or equal to the pivot element are to the left.
 - All the elements greater than the pivot element are to the right.

Partitioning



Quick Sort: Example

```
#include <stdio.h>
#include <stdlib.h>

void print(int x[],int low,int high)
{
    int i;

    for(i=low;i<=high;i++)
        printf(" %d", x[i]);
    printf("\n");
}

void swap(int *a,int *b)
{
    int tmp=*a;
    *a=*b; *b=tmp;
}
```

Contd.

```
void partition(int x[],int low,int high)
{
    int i=low+1,j=high;
    int pivot=x[low];
    if(low>=high) return;
    while(i<j)
    {
        while ((x[i]<pivot) && (i<high)) i++;
        while ((x[j]>=pivot) && (j>low)) j--;
        if(i<j) swap(&x[i],&x[j]);
    }
    if (j==high) {
        swap(&x[j],&x[low]);
        partition(x,low,high-1);
    }
    else
        if (i==low+1)
            partition(x,low+1,high);
        else {
            swap(&x[j],&x[low]);
            partition(x,low,j-1);
            partition(x,j+1,high);
        }
    }
}
```

Contd:

```
int main(int argc,char *argv[])
{
    int *x;
    int i=0;
    int num;

    num=argc-1;
    x=(int *) malloc(num * sizeof(int));

    for(i=0;i<num; i++)
        x[i]=atoi(argv[i+1]);
    printf("Input: ");
    for(i=0; i<num; i++)
        printf(" %d", x[i]);
    printf("\n");
    partition(x,0,num-1);

    printf("Output: ");
    for(i=0; i<num; i++)
        printf(" %d", x[i]);
    printf("\n");
}
```

Trace of Partitioning

```
./a.out 45 -56 78 90 -3 -6 123 0 -3 45 69 68
```

Input: 45 -56 78 90 -3 -6 123 0 -3 45 69 68

45 -56 78 90 -3 -6 123 0 -3 45 69 68

-6	-56	-3	0	-3	45	123	90	78	45	69	68
-56	-6	-3	0	-3	68	90	78	45	69	123	
-3	0	-3			45	68	78	90	69		
-3	0				69	78			90		

Output: -56 -6 -3 -3 0 45 45 68 69 78 90 123

Time Complexity

- Partitioning with n elements.

- No. of comparisons:

$n-1$

Choice of pivot element affects the time complexity.

- Worst Case Performance:

$$(n-1)+(n-2)+(n-3)+\dots\dots\dots +1 = n(n-1)/2$$

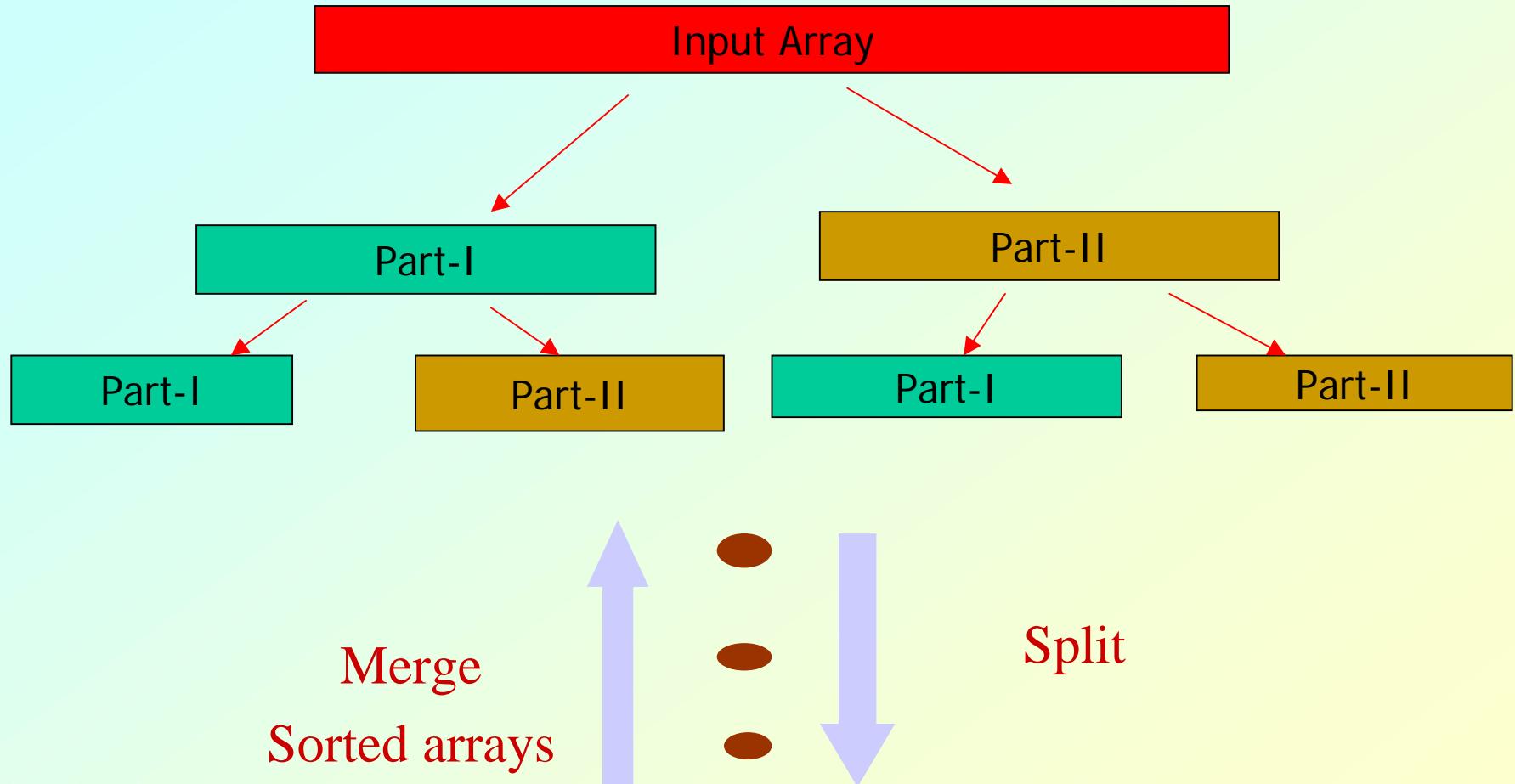
- Best Case performance:

$$(n-1)+2((n-1)/2-1)+4(((n-1)/2-1)-1)/2-1) \dots k \text{ steps}$$

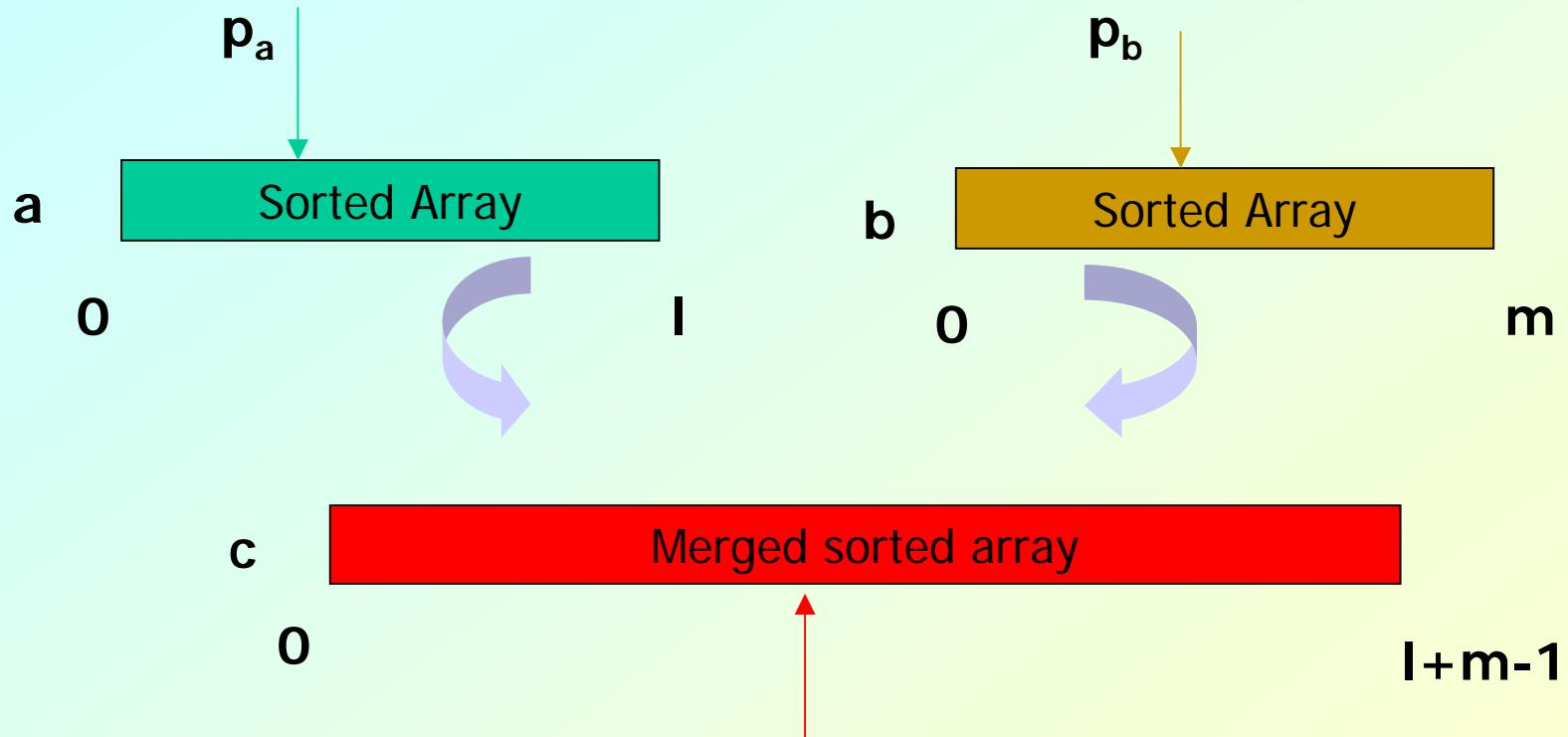
$$= O(n \cdot \log(n))$$

$$2^k = n$$

Merge Sort



Merging two sorted arrays



Move and copy elements pointed by p_a if its value is smaller than the element pointed by p_b in $(l+m-1)$ operations and otherwise.

Merge Sort

```
#include <stdio.h>
#include <stdlib.h>
main()
{
    int i, num;
    int a[ ]={-56,23,43,-5,-3,0,123,-35,87,56,75,80};
    for(i=0;i<12;i++) printf("%d ",a[i]); printf("\n");
    merge_sort(a, 12);
    printf("\n The sorted sequence follows \n");
    for(i=0;i<12;i++) printf("%d ",a[i]); printf("\n");
}
```

/* Recursive function for dividing an array a[0..n-1] into two halves and sort them and merge them subsequently. */

```
void merge_sort(int *a,int n)
{
int i,j,l,m;
int *b, *c;

if(n>1){
l=n/2; m=n-l;
b=(int *) calloc(l,sizeof(int));
c=(int *) calloc(m,sizeof(int));
```

```
for(i=0;i<l;i++) b[i]=a[i];
for(j=l;j<n;j++) c[j-l]=a[j];

merge_sort(b,l);
merge_sort(c,m);
merge(b,c,a,l,m);
free(b); free(c);
}
```

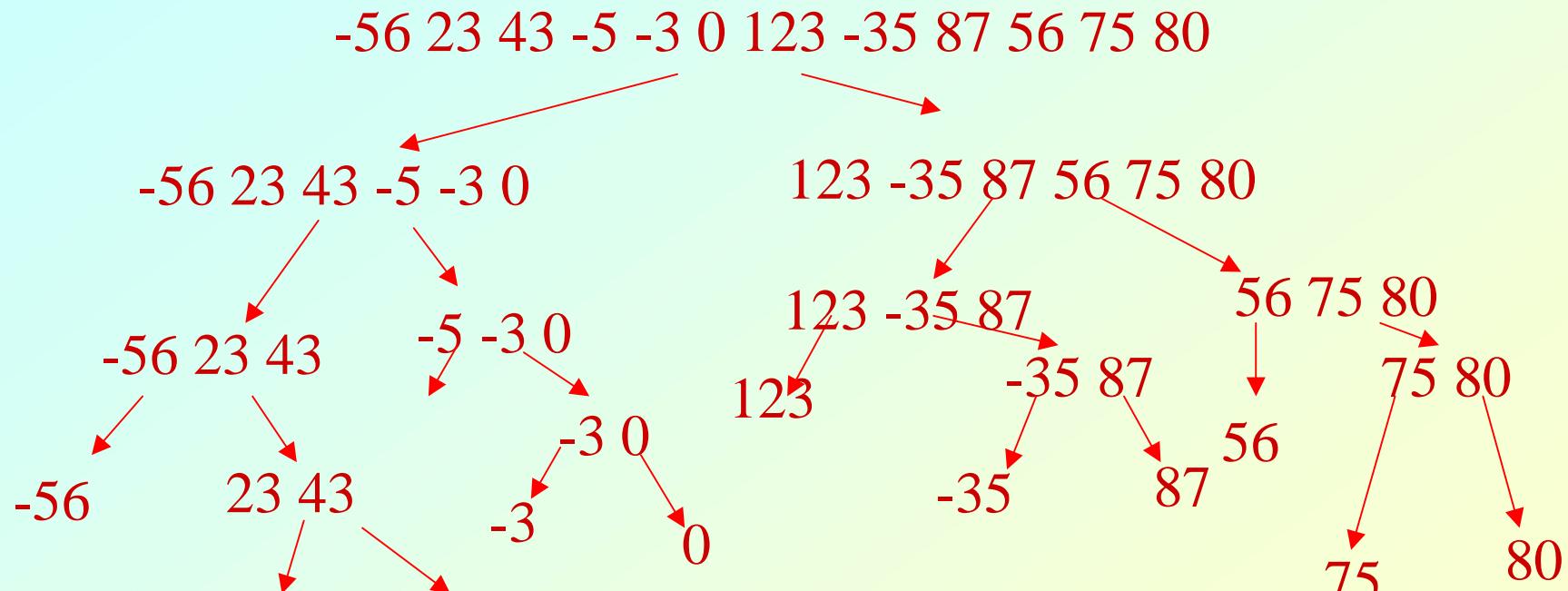
/* Merging of two sorted arrays a[0..m-1] and b[0..n-1] into a single array c[0..m+n-1] */

```
void merge(int *a,int *b,int *c,  
          int m,int n)  
{int i,j,k,l;  
  
i=j=k=0;  
  
do{  
if(a[i]<b[j]){ c[k]=a[i]; i=i+1; }  
else{ c[k]=b[j]; j=j+1;}    ←  
k++;  
} while((i<m)&&(j<n));
```

```
if(i==m){  
for(l=j;l<n;l++){ c[k]=b[l]; k++;}  
}  
else{  
for(l=i;l<m;l++){c[k]=a[l]; k++;}  
}
```

Pointer movement and copy operations.

Splitting Trace



Worst Case: $O(n \log(n))$

-56 -35 -5 -3 0 23 43 56 75 80 87 123