

Searching Elements in an Array:

Linear and Binary Search

Searching

- Check if a given element (called *key*) occurs in the array.
 - Example: array of student records; *rollno* can be the key.
- Two methods to be discussed:
 - If the array elements are unsorted.
 - Linear search
 - If the array elements are sorted.
 - Binary search

Linear Search

Basic Concept

- **Basic idea:**
 - Start at the beginning of the array.
 - Inspect elements one by one to see if it matches the key.
- **Time complexity:**
 - A measure of how long an algorithm takes to run.
 - If there are n elements in the array:
 - *Best case:* match found in first element (**1** search operation)
 - *Worst case:* no match found, or match found in the last element (**n** search operations)
 - *Average case:* **$(n + 1) / 2$** search operations

```
#include <stdio.h>

int linear_search (int a[], int size, int key)
{
    for (int i=0; i<size; i++)
        if (a[i] == key) return i;
    return -1;
}

int main()
{
    int x[]={12,-3,78,67,6,50,19,10}, val;
    printf ("\nEnter number to search: ");
    scanf ("%d", &val);
    printf ("\nValue returned: %d \n", linear_search (x,8,val));
}
```

- **What does the function `linear_search` do?**
 - It searches the array for the number to be searched element by element.
 - If a match is found, it returns the array index.
 - If not found, it returns -1.

Contd.

```
int x[] = {12, -3, 78, 67, 6, 50, 19, 10};
```

- Trace the following calls :

```
search (x, 8, 6) ;
```

Returns 4



```
search (x, 8, 5) ;
```

Returns -1



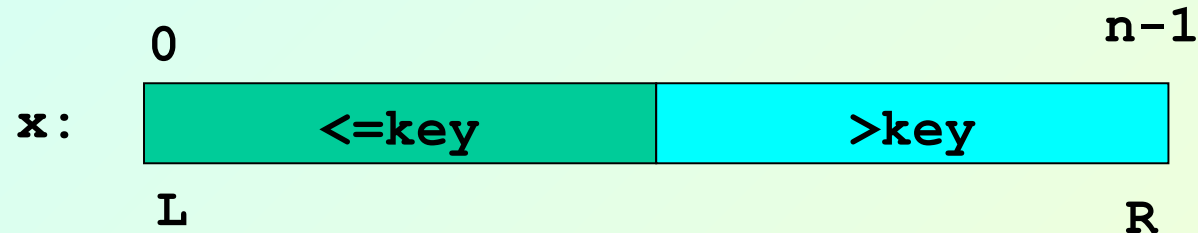
Binary Search

Basic Concept

- Binary search works if the array is *sorted*.
 - Look for the target in the middle.
 - If you don't find it, you can ignore half of the array, and repeat the process with the other half.
- In every step, we reduce the number of elements to search by half.

The Basic Strategy

- What we want?
 - Find split between values larger and smaller than *key*:



- Situation while searching:
 - Initially L and R contains the indices of first and last elements.
- Look at the element at index $[(L+R)/2]$.
 - Move L or R to the middle depending on the outcome of test.

Iterative Version

```
#include <stdio.h>

int bin_search (int a[], int size, int key)
{
    int L, R, mid;
    L = 0; R = size - 1;

    while (L <= R) {
        mid = (L + R) / 2;
        if (a[mid] < key) L = mid + 1;
        else if (a[mid] > key) R = mid - 1;
        else return mid; /* FOUND AT INDEX mid */
    }

    return -1; /* NOT FOUND */
}
```

```
int main()
{
    int x[]={10,20,30,40,50,60,70,80}, val;

    printf ("\nEnter number to search: ");
    scanf ("%d", &val);

    printf ("\nValue returned: %d \n", bin_search (x,8,val));
}
```

Recursive Version

```
#include <stdio.h>

int bin_search (int a[], int L, int R, int key)
{
    int mid;

    if (R < L) return -1;    /* NOT FOUND */
    mid = (L + R) / 2;
    if (a[mid] < key) return (bin_search(a, mid+1, R, key));
    else if (a[mid] > key) return (bin_search(a, L, mid-1, key));
        else return mid;    /* FOUND AT INDEX mid */
}
```

```
int main()
{
    int x[]={10,20,30,40,50,60,70,80}, val;

    printf ("\nEnter number to search: ");
    scanf ("%d", &val);

    printf ("\nValue returned: %d \n", bin_search (x,0,7,val));
}
```

Is it worth the trouble ?

- Suppose that the array x has 1000 elements.
- Ordinary search
 - If key is present in x , it would require 500 comparisons on the average.
- Binary search
 - After 1st compare, left with 500 elements.
 - After 2nd compare, left with 250 elements.
 - After 3rd compare, left with 125 elements.
 - After at most 10 steps, you are done.

Time Complexity

- If there are n elements in the array.
 - Number of searches required in the worst case: $\log_2 n$
- For $n = 64$ (say).
 - Initially, list size = 64.
 - After first compare, list size = 32.
 - After second compare, list size = 16.
 - After third compare, list size = 8.
 -
 - After sixth compare, list size = 1.

$2^k = n$, where k is the number of steps.

$$k = \log_2 n$$

$$\log_2 64 = 6$$

$$\log_2 1024 = 10$$

Sorting

The Basic Problem

- Given an array

`x[0], x[1], ... , x[size-1]`

- reorder entries so that

`x[0] ≤ x[1] ≤ ... ≤ x[size-1]`

- List is in non-decreasing order.
- We can also sort a list of elements in non-increasing order.

Example

- **Original list:**

10, 30, 20, 80, 70, 10, 60, 40, 70

- **Sorted in non-decreasing order:**

10, 10, 20, 30, 40, 60, 70, 70, 80

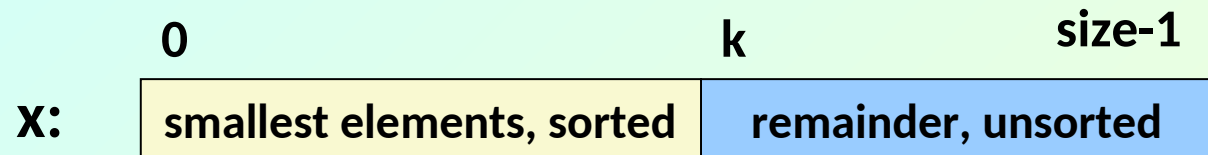
- **Sorted in non-increasing order:**

80, 70, 70, 60, 40, 30, 20, 10, 10

Selection Sort

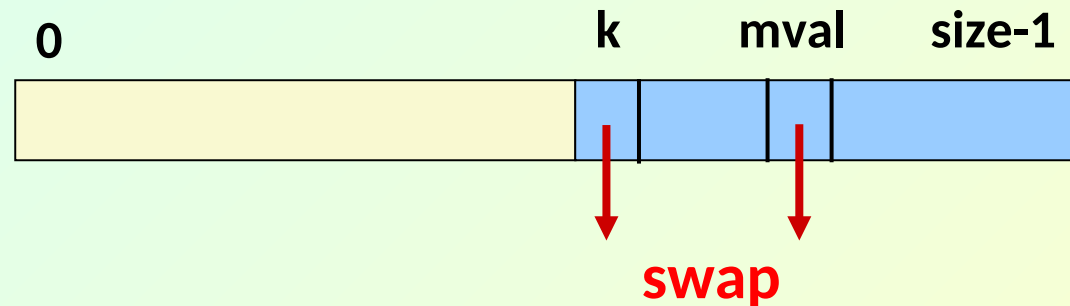
How it works?

- General situation :



- Steps :

- Find smallest element, `mval`, in `x[k..size-1]`
- Swap smallest element with `x[k]`, then increase `k` by 1



An example worked out

PASS 1:

10	5	17	11	-3	12	Find the minimum
10	5	17	11	-3	12	Exchange with 0th
<u>-3</u>	5	17	11	10	12	element

PASS 2:

<u>-3</u>	5	17	11	10	12	Find the minimum
<u>-3</u>	5	17	11	10	12	Exchange with 1st
<u>-3</u>	<u>5</u>	17	11	10	12	element

PASS 3:

<u>-3</u>	<u>5</u>	17	11	10	12	Find the minimum
<u>-3</u>	<u>5</u>	17	11	10	12	Exchange with 2nd
<u>-3</u>	<u>5</u>	<u>10</u>	11	17	12	element

PASS 4:

-3 5 10 **11** 17 12

Find the minimum

-3 5 10 11 17 12

Exchange with 3rd

-3 5 10 11 17 12

element

PASS 5:

-3 5 10 11 17 12

Find the minimum

-3 5 10 11 **17** 12

Exchange with 4th

-3 5 10 11 12 17

element

Subproblem

```
/* Yield index of smallest element in x[k..size-1];*/  
  
int min_loc (int x[], int k, int size)  
{  
    int j, pos;  
  
    pos = k;  
    for (j=k+1; j<size; j++)  
        if (x[j] < x[pos])  
            pos = j;  
    return pos;  
}
```

The main sorting function

```
/* Sort x[0..size-1] in non-decreasing order */  
  
int sel_sort (int x[], int size)  
{ int k, m;  
  
  for (k=0; k<size-1; k++)  
  {  
    m = min_loc (x, k, size);  
    temp = a[k];  
    a[k] = a[m];  
    a[m] = temp;  
  }  
}
```

SWAP

```
int main()
{
    int x[ ]={-45,89,-65,87,0,3,-23,19,56,21,76,-50};
    int i;
    for(i=0;i<12;i++)
        printf("%d ",x[i]);
    printf("\n");
    sel_sort(x,12);
    for(i=0;i<12;i++)
        printf("%d ",x[i]);
    printf("\n");
}
```

```
-45 89 -65 87 0 3 -23 19 56 21 76 -50
-65 -50 -45 -23 0 3 19 21 56 76 87 89
```

Analysis

- How many steps are needed to sort n items ?
 - Total number of steps proportional to n^2 .
 - No. of comparisons?

$$(n-1) + (n-2) + \dots + 1 = n(n-1)/2$$

Of the order of n^2

- Worst Case? Best Case? Average Case?

Insertion Sort

Basic Idea

- Insert elements one at a time, and create a partial sorted list.
 - Sorted list of 2 elements, 3 elements, 4 elements, and so on.
- In general, in every iteration an element is compared with all the elements before it.
- After finding the position of insertion, space is created for it by shifting the other elements and the desired element is then inserted at the suitable position.
- This procedure is repeated for all the elements in the list.

An example worked out

PASS 1:

<u>10</u>	5	17	11	-3	12	<code>item = 5</code>
<u>?</u>	<u>10</u>	17	11	-3	12	<code>compare 5 and 10</code>
<u>5</u>	<u>10</u>	17	11	-3	12	<code>insert 5</code>

PASS 2:

<u>5</u>	<u>10</u>	17	11	-3	12	<code>item = 17</code>
<u>5</u>	<u>10</u>	<u>17</u>	11	-3	12	<code>compare 17 and 10</code>

PASS 3:

<u>5</u>	<u>10</u>	<u>17</u>	11	-3	12	<code>item = 11</code>
<u>5</u>	<u>10</u>	<u>?</u>	<u>17</u>	-3	12	<code>compare 11 and 17</code>
<u>5</u>	<u>10</u>	<u>?</u>	<u>17</u>	-3	12	<code>compare 11 and 10</code>
<u>5</u>	<u>10</u>	<u>11</u>	<u>17</u>	-3	12	<code>insert 11</code>

PASS 4:

5 10 11 17 -3 12
5 10 11 ? 17 12
5 10 ? 11 17 12
5 ? 10 11 17 12
? 5 10 11 17 12
-3 5 10 11 17 12

item = -3

compare -3 and 17

compare -3 and 11

compare -3 and 10

compare -3 and 5

insert -3

PASS 5:

-3 5 10 11 17 12
-3 5 10 11 ? 17
-3 5 10 11 ? 17
-3 5 10 11 12 17

item = 12

compare 12 and 17

compare 12 and 11

insert 12

Insertion Sort

```
void insert_sort (int list[], int size)
{
    int i,j,item;

    for (i=1; i<size; i++)
    {
        item = list[i] ;
        j = i - 1;
        while ((item < list[j]) && (j >= 0))
        {
            list[j+1] = list[j];
            j--;
        }
        list[j+1] = item ;
    }
}
```

```

int main()
{
    int x[ ]={-45,89,-65,87,0,3,-23,19,56,21,76,-50};
    int i;
    for(i=0;i<12;i++)
        printf("%d ",x[i]);
    printf("\n");
    insert_sort (x,12);
    for(i=0;i<12;i++)
        printf("%d ",x[i]);
    printf("\n");
}

```

```

-45 89 -65 87 0 3 -23 19 56 21 76 -50
-65 -50 -45 -23 0 3 19 21 56 76 87 89

```

Time Complexity

- Number of comparisons and shifting:

- Worst case?

$$1 + 2 + 3 + \dots + (n-1) = n(n-1)/2$$

- Best case?

$$1 + 1 + \dots \text{ to } (n-1) \text{ terms} = (n-1)$$

Bubble Sort

How it works?

- The sorting process proceeds in several passes.
 - In every pass we go on comparing neighboring pairs, and swap them if out of order.
 - In every pass, the largest of the elements under considering will *bubble* to the top (i.e., the right).

10 5 17 11 -3 12

5 10 17 11 -3 12

5 10 17 11 -3 12

5 10 11 17 -3 12

5 10 11 -3 17 12

5 10 11 -3 12 17

← Largest

An example worked out

PASS 1:

10	5	17	11	-3	12
5	10	17	11	-3	12
5	10	17	11	-3	12
5	10	11	17	-3	12
5	10	11	-3	17	12
5	10	11	-3	12	<u>17</u>

PASS 2:

5	10	11	-3	12	<u>17</u>
5	10	11	-3	12	<u>17</u>
5	10	11	-3	12	<u>17</u>
5	10	-3	11	12	<u>17</u>
5	10	-3	11	<u>12</u>	<u>17</u>

PASS 3:

5	10	-3	11	<u>12</u>	<u>17</u>
5	10	-3	11	<u>12</u>	<u>17</u>
5	-3	10	11	<u>12</u>	<u>17</u>
5	-3	10	<u>11</u>	<u>12</u>	<u>17</u>

PASS 4:

5	-3	10	<u>11</u>	<u>12</u>	<u>17</u>
-3	5	10	<u>11</u>	<u>12</u>	<u>17</u>
-3	5	<u>10</u>	<u>11</u>	<u>12</u>	<u>17</u>

PASS 5:

-3	5	<u>10</u>	<u>11</u>	<u>12</u>	<u>17</u>
<u>-3</u>	<u>5</u>	<u>10</u>	<u>11</u>	<u>12</u>	<u>17</u>

Sorted ←

- **How the passes proceed?**
 - In pass 1, we consider index 0 to $n-1$.
 - In pass 2, we consider index 0 to $n-2$.
 - In pass 3, we consider index 0 to $n-3$.
 -
 -
 - In pass $n-1$, we consider index 0 to 1.

Bubble Sort

```
void swap(int *x, int *y)
{
    int tmp = *x;
    *x = *y;
    *y = tmp;
}
```

```
void bubble_sort
    (int x[], int n)
{
    int i, j;

    for (i=n-1; i>0; i--)
        for (j=0; j<i; j++)
            if (x[j] > x[j+1])
                swap(&x[j], &x[j+1]);
}
```

```

int main()
{
    int x[ ]={-45,89,-65,87,0,3,-23,19,56,21,76,-50};
    int i;
    for(i=0;i<12;i++)
        printf("%d ",x[i]);
    printf("\n");
    bubble_sort (x,12);
    for(i=0;i<12;i++)
        printf("%d ",x[i]);
    printf("\n");
}

```

```

-45 89 -65 87 0 3 -23 19 56 21 76 -50
-65 -50 -45 -23 0 3 19 21 56 76 87 89

```

Time Complexity

- Number of comparisons :

- Worst case?

$$1 + 2 + 3 + \dots + (n-1) = n(n-1)/2$$

- Best case?

- Same

- How do you make best case with $(n-1)$ comparisons only?
 - By maintaining a variable **flag**, to check if there has been any swaps in a given pass.
 - If no swaps during a pass, the array is already sorted.

```
void bubble_sort (int x[], int n)
{
    int i, j, flag;

    for (i=n-1; i>0; i=flag)
    {
        flag = 0;
        for (j=0; j<i; j++)
            if (x[j] > x[j+1])
            {
                swap(&x[j], &x[j+1]);
                flag = j;
            }
    }
}
```

Some Efficient Sorting Algorithms

- Two of the most popular sorting algorithms are based on **divide-and-conquer** approach.
 - Quick sort
 - Merge sort
- Basic concept (divide-and-conquer method):

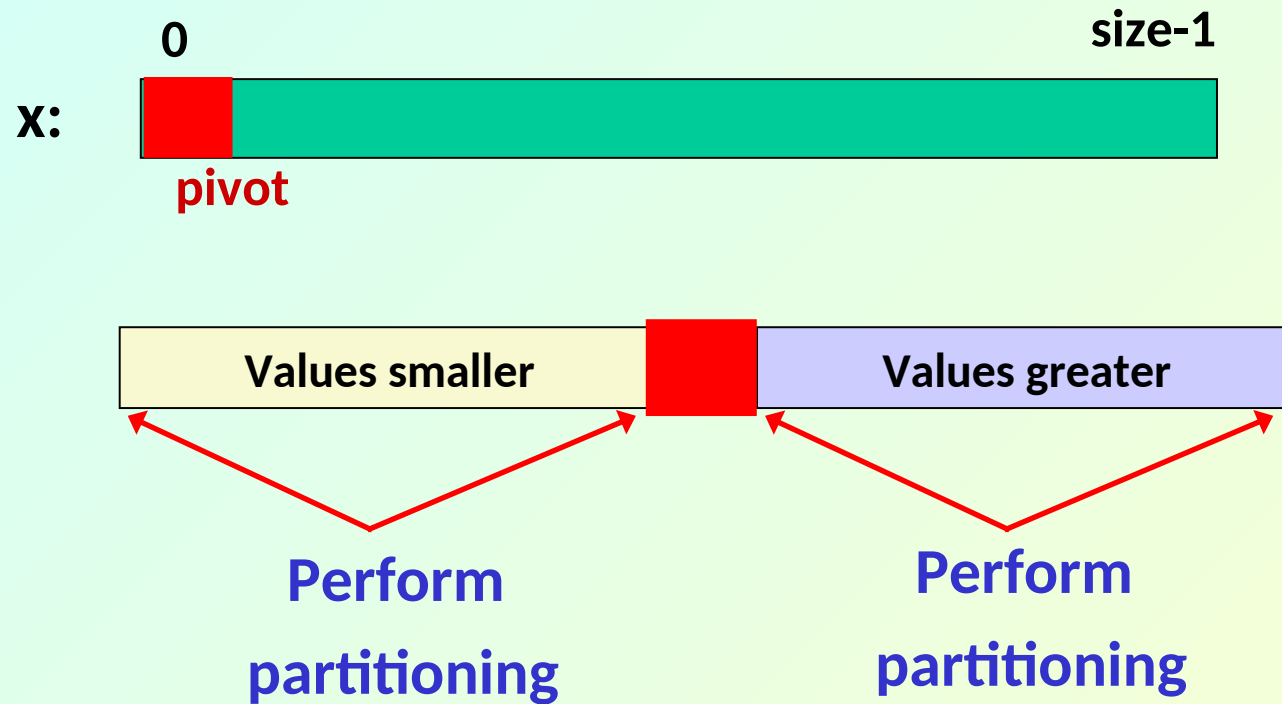
```
sort (list)
{
  if the list has greater than 1 elements
  {
    Partition the list into lowlist and highlist;
    sort (lowlist);
    sort (highlist);
    combine (lowlist, highlist);
  }
}
```

Quick Sort

How it works?

- At every step, we select a *pivot element* in the list (usually the *first* element).
 - We put the pivot element in the final position of the sorted list.
 - All the elements less than or equal to the pivot element are to the left.
 - All the elements greater than the pivot element are to the right.

Partitioning



```
void print (int x[], int low, int high)
{
    int i;

    for(i=low; i<=high; i++)
        printf(" %d", x[i]);
    printf("\n");
}
```

```
void swap (int *a, int *b)
{
    int tmp=*a;
    *a=*b;
    *b=tmp;
}
```

```
void partition (int x[], int low, int high)
{
    int i = low+1, j = high;
    int pivot = x[low];
    if (low >= high) return;
    while (i<j) {
        while ((x[i]<pivot) && (i<high)) i++;
        while ((x[j]>=pivot) && (j>low)) j--;
        if (i<j) swap (&x[i], &x[j]);
    }
    if (j==high) {
        swap (&x[j], &x[low]);
        partition (x, low, high-1);
    }
    else
        if (i==low+1)
            partition (x, low+1, high);
        else {
            swap (&x[j], &x[low]);
            partition (x, low, j-1);
            partition (x, j+1, high);
        }
}
```

```
int main (int argc, char *argv[])
{
    int x[] = {-56,23,43,-5,-3,0,123,-35,87,56,75,80};
    int i=0;
    int num;

    num = 12;    /* Number of elements */

    partition(x,0,num-1);

    printf("Sorted list: ");
    print (x,0,num-1);
}
```

Trace of Partitioning: an example

45 -56 78 90 -3 -6 123 0 -3 45 69 68

45 -56 78 90 -3 -6 123 0 -3 45 69 68

-6 -56 -3 0 -3 45 123 90 78 45 69 68

-56 -6 -3 0 -3 68 90 78 45 69 123

-3 0 -3 45 68 78 90 69

-3 0 69 78 90

Sorted list: -56 -6 -3 -3 0 45 45 68 69 78 90 123

Time Complexity

- Worst case:

n^2 ==> list is already sorted

- Average case:

$n \log_2 n$

- Statistically, quick sort has been found to be one of the fastest algorithms.