Advanced graph theory: Homework 2: CS60047 Autumn 2023

Instructor: Sudebkumar Prasant Pal Teaching assistant: Koustav De

September 01, 2023

Advanced graph theory: Homework 2: CS60047 Autumn 2023

- 1. State Tutte's theorem precisely and illustrate multiple or unique "bad" sets for connected, even, undirected graphs that do not have a perfect matching. One example is $K_{1,3}$. Another example was worked out in class with 20 vertices and a maximum matching of only 9 edges.
- 2. Let $G'(V, E) \neq K_n$ be an *n*-vertex simple connected undirected graph where adding any additional edge e to G' would introduce a perfect matching in G' + e, given that G' has no perfect matching. Naturally G' has an even number of vertices.

Apply Tutte's theorem to answer the following questions.

If $S \subseteq V$ is the "bad" set as per Tutte's theorem whereby o(G'-S) > |S|, then show that

- (i) S induces a complete subgraph in G',
- (ii) each connected component of G' S also introduces a complete subgraph in G', and (iii) the vertices in S are connected to all the vertices in G'.

[Hint: Suppose the edge $\{u,v\}$ is absent in G' where $u,v\in S$. Then adding this edge to G' introduces a perfect matching in $G'+\{u,v\}$ but does not change violated Tutte's condition $o(G'+\{u,v\}-S)=o(G-S)>|S|$, a contradiction. Similar arguments apply for the connected components of G'-S, and also to edges between S and the components of G'-S.]

- 3. Show that 3-regular graphs with no cut edges have a 1-factor.
- 4. Show that every 3-regular graph with at most two cut edges has a 1-factor.
- 5. Show that interval graphs as well as their complements are perfect.

[Given a set of closed intervals on a line, assign a vertex for each interval and an edge for every pair of intersecting intervals to define the interval graph. Sweep a line perpendicular the intervals.]

- 6. Show that the intersection graph of subtrees of a tree is a chordal graph.
- 7. Show that every minimal (by inclusion) vertex separator subset in a chordal graph is a complete graph.
 - [Show that every component of G-S for a minimal separator S, is adjacent to every vertex in S. Also, show that any two vertices u and v in S appear in a 4-cycle.]
- 8. Suppose a graph G has no 1-factor. Show that there must be a vertex x such that each edge adjacent to x is in some maximum matching.

[Let x be an unmatched vertex in a maximum matching F, and let y be a neighbour of x. If y is not covered by an edge e of F then F + xy would be a larger matching. So, see F - e + xy is also a maximum matching.]

9. Let G be a connected 2d-regular graph with an even number of edges. Prove that G has a d-factor.

[Consider an Euler trail L of G. Consider every second edge of the trail. This is possible as the number of edges is even. The selected edges for a d-factor.]

10. Show that the number of triangles in any simple graph of n vertices and m edges is at least $\frac{4m}{3n}(m-\frac{n^2}{4})$.

[For any edge xy there are at least d(x) + d(y) - n vertices adjacent to both x and y. So, this is also a lower bound on the number of triangles sitting on xy. A third of the sum of such estimates over all edges, lower bounds the number of triangles in the graph. This estimate is a third of $\sum (d(x))^2 - mn$, which is at least a third of n times the square of the average of vertex degrees minus mn by the Cauchy-Schwartz inequality.]